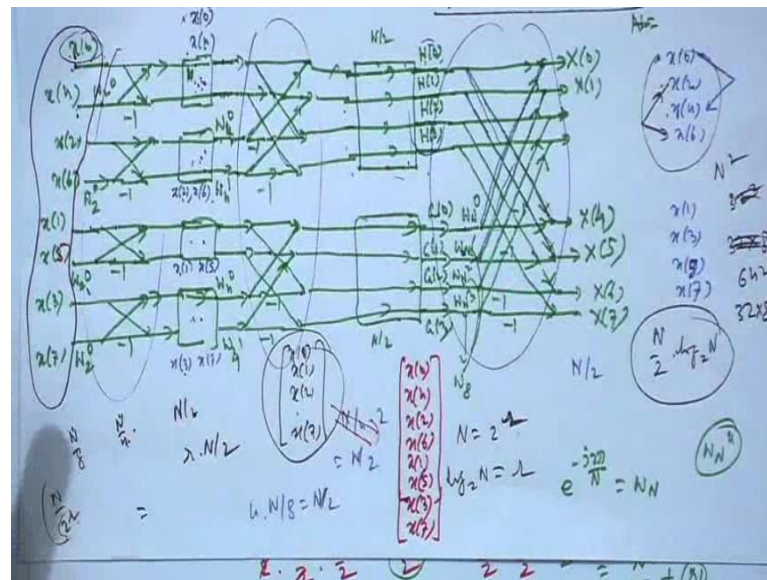


Discrete Time Signal Processing
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Lecture – 39
Complexity Analysis of FFT

So, we continue from where we stopped. Now, consider an example.

(Refer Slide Time: 00:29)



Suppose, N equal to 8. So, I have got x_0, x_1, \dots, x_7 . So, suppose I have got the H ; H_0, H_1, H_2, H_3 suppose they are available because N is 8; so N by 2 is 4. If you take the even points X_0, X_2, X_4, X_6 , put them, and then, takes their DFT. Now, they are coming out H ; and here, suppose we have got this for the odd part, ok. Then, from this formula, you start; you have to find out some output, final output x at 0. So, that will be H_0 plus this factor times G_0 . This factor by the way, e to the power, excuse me, we adopted notation, very widely used notation. This factor is called a twiddle factor; we call it W_N ; and so, this is W_N to the power k . So, it was W_N to the power k here, this term.

So, start with k equal to 0, and find out capital X_0 , X_1 , x , upper half. So, suppose, I have to find out X_0 ; that will be H_0 plus W^N whole to the power k , k is 0, W into the power 0 into G_0 . So, suppose I multiply it by, I substitute W to the power 0 is 1, but I am still writing. And then, I add these one, two arrows meet at a point, assume there is an adder here; because, if we start drawing adder everywhere, it will be very clumsy and complicated. So, when I write some constant by the side of an H , that means, you are multiplying the, the signal G_0 by this, and they up, and is getting added. So, what you have is H_0 , plus W^N to the power 0 into G_0 ; that is your X_0 .

On the other hand, if k jumps from 0 to 0 by N by 2, N by 2 here is 4; so, from 0 to 4, a or there was 0, I have found out. Now, 0 plus N by 2, that is, 0 plus 4; I will have same H_0 , because H_k is periodic over N by 2. So, N by 2 is 4. So, H_0 and H_4 , they are same; G_0 , G_4 , they are same; but only this part will be minus of what it was. So, therefore, what it will be, same H_0 , and same G_0 ; but G_0 now will be multiplied by the same constant, but after that, minus sign; it will be subtracted. So, G_0 multiplied by the same constant; G_0 into this, G_0 into this; but after that, it will be multiplied by minus 1, or, and then, added with the same H_0 . You can either multiply by minus 1, but in practice, we will not use a multiplier; we will just use a subtractor; adder we know how to do addition; in binary domain, we can also design a subtractor, if you are using 2's complement number system, you take 2's complement, and add, you get the same thing. But, this will give me X_4 .

Similarly, by the same way, suppose here, you want to find out capital X_1 . So, it will be H_1 , plus G_1 into W^N to the power 1. So, that means, G_1 into W^N to the power 1. So, this will be added, to your X_1 . And now, instead of capital X_1 , if k jumps from 1 to 1 plus N by 2; N by 2 is 4. So, 1 to 5, you will have the same H , same G ; there is H_1 , H_5 same; G_1 , G_5 same; this factor will be same; but this will undergo sign change; it will become minus. So, once again, this will be your X_5 ; same manner, k equal to 2; k equal to 2 means, H_2 plus W^N to the power 2 G_2 . And, instead of 2, if k jumps from 2 to 2 plus N by 2, that is 4; that is 2 plus 4 is 6; we will have the same thing, same gain, but there will be undergo a minus sign. It will come from here, and by the same token, lastly, also G_3 , we multiplied by W^N to the power 3; then, added in this X_1 minus; in the same logic, and then added; that will give you X_7 .

But, what is capital H_0, H_1, H_2, H_3 is a DFT of X sequence x_0, x_2, x_3, x_4 . So, this is that, it was computing that N by 2 point DFT, N by 2 point DFT; on what sequence, this was working on this sequence, x_0 ; then, next event, then, next event, then next event. And, this was working on first odd, next odd, next odd, next odd; this is 5, on this. But, this is a sequence; on this also I can apply the same procedure. I can take out its even component, these 2, the 0th term, second term, and what components these 2, either we subsequence x_0, x_4 , and, another subsequence x_2, x_6 . By the same token this DFT can be obtained. Now, this box has no meaning; this will be another box. Now, this was 4 point; this is the 2 point; this will give you.

So, I already have used up in H and G , I will not use the name here; it will take this x_0, x_4 here, as a 2 point sequence, and it will give 2 outputs. And, this will take x_3 , sorry; this will take x_0, x_4 , and this will take x_2, x_6 ; 0, 0th point, first point, second point, third point; first and third, first and third are the odd ones, H_2, H_6 ; and 0th and second, there is x_0, x_4 , (Refer Time: 08:36) x_0, x_4 , I calculated 2 point DFTs. They come out here, and here; here also they come out. There, what I have to do? This will get multiplied by now W ; earlier, it was W_8 ; N was 8; you can, instead of W_N , here, you can write as W_8 , in all of them. I wrote in a general way; all of them, it will W_8 everywhere.

Because, you have to, the total length was 8. Similarly, here the total length is 4; that is, the length of the sequence is 4. I am trying to find out its DFT. So, it is 4, and I am dividing this into 2 halves; due to two halves, but W formula, the W_N , N was coming from this length, final length that is 4; that is why, now this will be W_4 to the power 0, W_4 to the power 1. After that, they will be added, and then, this direction will undergo the minus 1. This direction will undergo the minus 1; this is that we have already seen. So, then, this will be here; this will be here; it is just replica, I mean, same thing, and this will go like this. Same way, x_1 and x_5 , they will come under this box; x_1, x_5 ; and, x_3, x_7 here; N is 4 again. So, W_4 ; here, W_4 to the power 0, W_4 to the power 1; this will get added. Here, this will not add; mind you, I am telling. Now, these two get added; and in this direction, this minus 1, minus 1, it gets added; these get added. So, it is here; it is here; it is here; all right. So, now, again, this is having 2 point sequence now, x_0 and x_4 . This was having x_2, x_6 . I am writing horizontally, x_2, x_6 . This was having x_1, x_5 ; and, this is having x_3, x_7 , right. So, 2 point.

Again, I break it into its even part, odd part; very simple now. Even part has x_0 ; odd part has x_4 . So, every, I mean, even part is a 1 point sequence; odd part is a 1 point sequence. So, 1 point sequence have only 1 point DFT, and that DFT is same as itself; summation $\sum_{n=0}^{N-1} x_n e^{-j2\pi kn/N}$ by, in this formula the (Refer Time: 11:49) DFT formula, the (Refer Time: 11:52) DFT formula, if capital N is 1, it is $n=0$ to 0 , x_0 into e^{j0} , so, x_0 , capital X k, that is, k equal to 0 only one possibility. So, that is same as the, capital X 0 is same as small x_0 . So, DFT is the value of the sequence itself; X, the sample itself, it is a DFT. 1 point sequence, 1 point DFT; sample value and DFT are same. So, I will have one point here, and one point here. So, W, here, it was 4, now W^2 ; this will be what, x_0 , x_4 ; and, this will get added, and this will get minus 1, and then, added. So, boxes do not do anything now; you can remove the boxes; it was just for explaining.

Similarly, here x_2 , x_6 , 1 point sequence, 1 point sequence. So, its DFT itself, x_6 is DFT itself; H_k , G_k , like that. So, this we combine, this I multiply by again by W^2 , to the power 0, because W^N , you go from 0, 1 up to N, N by 2 minus one. So, N here is 2. So, 2 by 2 minus 1 is 0. So, you can go from the 0 to 0 only. And then, you add, and here minus 1, like this. Then again, by the same philosophy, sequence is x_1 , x_3 . Now, W^2 to the power 0 is 1, I agree, but still, I will be writing like that; last one is x_3 , x_7 . So, this is the entire diagram, you see. Now, let us see what computational advantage we get. We suppose have, here N is 8, but suppose, a general N; for general N, I have got this box. This is, suppose generality, this is by (Refer Time: 14:27) I have got H_k , for k equal to 0 1 up to N by 2 minus 1; and, G_k also. So, for this butterfly, these are butterflies.

How many are there? For this structure, this multiplication is common; for the next butterfly also, this multiplication is common. This butterfly, this is common, this is common. How many butterflies? N by 2 butterflies; like this one, and this one, they together form one. Then, this one, and this one, they form one, totally, N by 2 butterflies; and for per butterfly, only one multiplication, only one multiplication. So, you have got N by 2 complex multiplications. Of course, I can say that, W into the power 0 is 1. So, no need to multiply. So, it will be N by 2 minus 1; all those things I am not doing; very rough calculation, N by 2 multiplications, right.

Then again, I assume that, these are somehow generated. Now, how to generate them? This is again generated by this. That suppose, these 2 DFT s are available, then, you have got how many? This was N by 2; N by 2. So, now, N by 4, N by 4, N by 4, N by 4 here, N by 4, N by 4; how many butterflies? N by 4 butterflies here, it was N by 2, N by 2; how many butterflies? N by 2; N by 2, because, one from here, and here, and another from here, or here; one butterfly means, you take one output, and one output, you have combine this way after multiplying, and then, again, going in this direction, multiplying by minus 1, and in this direction, there is 1 butterfly. So, obviously, if we use N by 2, here, N by 2, here I have totally N by 2 butterfly. Similarly, if it is N by 4, N by 4, you have N by 4 butterflies. Like, N was 8 in the previous example. So, 8 by 4 is 2.

So, you have got one butterfly; this, another butterfly this. 8 by 4 butterflies; or again, what butterfly, how many multiplier? 2 multipliers, N by 4; how many butterflies I have got, N by 4; per butterfly one multiplier. So, N by 4 into 1; but, I have got N by 4 here, another N by 4 here. So, again, N by 4 into 2, so N by 2; next time, N by 4, if N by 4 will give rise to N by 2, sorry, N by 8, N by 8, total N by 8 butterflies; per butterfly 1 multiplier, N by 8 multiplier, but, how many N by 8s? Because, this is giving rise to N by 8, N by 8; this is giving rise to N by 8, N by 8; this is giving rise to N by eight N by 8; this is giving rise to N by 8 N by 8. So, how many? 4. So, it will be again N by 2; I repeat, here I, what I had? These 2 blocks, like, these 2 blocks, give rise to, I mean, because, I was computing this DFT. So, these 2 blocks, they give rise 2 butterflies. In general, it was N by 2, N by 4, N by 4. So, this block, these 2 blocks, give rise to N by 4 butterflies; N by 4 into one multiplication, because, per butterfly 1, and 1 more case here, 2 into N by 4.

Now, this was N by 2, N by 4. Now, this N by 4, this is giving rise to N by 8, N by 8; N by 8 butterflies; per butterfly, one multiplier; N by 8 into 1. Again here also N by 8 into 1, here also, N by 8 into 1, N by 8 into 1; so 4 into N by 8; again, N by 2, so every stage, this is one stage; this is one stage; this is one stage. Every stage, your length of the DFT reduces, but number of such parallel butterflies, you know, parallel, this increases. So, total multiplier, again same. So, if you have got, say k number of stages, 1 number of stages, 1 into N by 2; that will be the total multiplications. Now, how many stages you have? N by 2; divisible by N by 2.

The next time, N by 4; next time, N by 8; finally, it will be 1 point. So, that will happen only if it is N by N . So, how many stages I have? It is N by 2; then, N by 4; N by what? So, Actually, your (Refer Time: 19:20) by N was 2 to the power r . So, basically, you have r stages, r into N by 2. You have r stages; like, N was 8 here. So, 2 to the power 3, one stage, one stage, one stage; per stage I have got? N by 2 computation; N is 8; you can put 8 here, or write N by 2, 3 into N by 2. Why r , because, N by 2; then N by 4 then, N by 8. N by 2; if the basic block was, I am dividing into N by 4, then, N by 8. So, 1 stages, another stage, another stage. So, how many cases? It was 2 to the power 1, 2 to the power 2, 2 to the power 3, up to 2 to the power r ; then only, N by N , 1. So, I will get 1 point. Here, I am taking N point sequence, and breaking into N by 4, N by 4; into N by 2, N by 2; total was length N , even point here, odd point here; that is, corresponding to these DFTs, N by 2, N by 2.

Then again, each N by 2 was broken down into further half. So, N by 4, N by 4; N by 4, N by 4; they are corresponding to this sequence, $x_0, x_4, x_2, x_6; x_1, x_3, x_5, x_7$. Then again, this was broken into two; N by 8, N by 8, N by 8, N by 8, N by 8, N by 8, N by 8, N by 8; till the point I get 1 point sequence, x_0, x_4 ; it become $x_0, 1$ point, that is the even point of this 2, length 2 sequence; because, length 2 x_0, x_4 , even point or odd point. So, I got x_0 is the only point, even point, x_4 the only point, odd point. So, it is a 1 point sequence. So, then, then only I, can stop.

So, when do I reach 1 point? I started with N by 2, stage 1; then, N by 4, 2 to the power 2, stage 2; N by 8, stage 3; three stages; N by 2 to the power r , to the power is N ; then only, N by N is equal to 1. So, after r stages, sequences then break into sub sequence, then break into sub sequence, I get 1 point sequences. There is here, here, there is 4 point, 4 point; then, 4 points were divided into 2 point, 2 point, 2 point, 2 point; then, 2 point was divided into 1 point, 1 point, 1 point, 1 point. So, sequence length becomes 1 point; that happens only when I have got r stages. I repeat again; start with N by 2, N by 2 point; N by 2 point; then, you go for another stage, everybody is N by 4 point; not still 1 point; then N by 8 point. So, when will I have 1 point sequence, like x_0 is a 1 point sequence; x_4 is a 1 point sequence; when? When? In the denominator, I have N ; N is 2 to the power r . So, I start it with 2 to the power 1, one stage; if I worked N by 4, 2 stages, 2 to the power 2; if I go for N by 8, 3 stages, 2 to the power 3. So, if I go up 2 to the

power, I mean, 2^r , r stages, then only, I get 1 point sequence; I stop this business. So, r stages; per stage, total computation is N by 2 complex multiplications. So, total will be r into N by 2, which also you write as r was, what was r , this. So, $\log_2 N$, that is 2^r is your r , right. So, there is your famous formula comes, N by $2 \log_2 N$.

So, it is a huge saving. Originally, it was N^2 . Suppose, N is 64, 64^2 and it is 32 into 64. If N is 64, it is $8 \log_2 64$, 64 is 8, 8 into 32. So, you have 32 square on one hand, and then, is just 32 into 8. See, the kind of Not even 32; 6 into 8; sorry, 32. If you have N equal to 64, sorry, if is N equal to 64, originally you had a 64 square; or here, 64 by 2, that is 32, and $\log_2 64$, that is 6. See the huge difference. This is 64 square. So, 32 into 2 into 64; that is, 32 into 128; and here it is only 32 into 8, huge difference; 16 times less. And, if the capital N is large, I mean, your gain will be more and more. Lastly, you see, our original sequence was in this order, x_0, x_1 ; like, if you write a vector, it will be, if you it write in a vector form, it will be, you know, this kind of form x_0, x_1, x_2 , dot, dot, dot; x_7 . Write the vector form.

This is the sequence, but look at this order. This is, I am writing here only, this is different. This is x_0 , then x_4 , then x_2 , then x_6 , then x_1 , then x_5 , it cannot be 3, it is 5; x_5 ; then x_3 , x_7 . So, the whole order has got disrupted. So, question is what is this new order? If I give a vector, and that vector is long, there is no point in going back calculating like this, and you know, forming the new vector. There is a way to get this new vector.

(Refer Slide Time: 25:24)

Handwritten notes on a whiteboard:

- $N = 2^r$
- $n = 0, 1, \dots, N-1$
- $n = a_{r-1} \dots a_1 a_0$
- $n = a_0 + a_1 2 + a_2 2^2 + \dots + a_{r-1} 2^{r-1}$
- $$\begin{aligned} n: \text{even} &\Rightarrow a_0 = 0 \\ n: \text{odd} &\Rightarrow a_0 = 1 \end{aligned}$$
- Small box: © IIT KGP

How to get this new vector? Now, suppose your N is 2 to the power r . So, I have got number 0, 1, ..., $N-1$, so, index n , x_n ; n will be 0, 1, dot, dot, dot, up to N minus 1. Suppose, there is an example N is 64; that is 2 to the power 8. So, it will be 0, 1, up to 63. Or, say N equal to 16. So, it will be 0, 1 up to 15. Now, all this decimal numbers, integers, 0, 1, 2, 3, up to 15, if I want to represent them by binary words, how many bits will you require? 4 bits here; 2 to the power 4; 16, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0, 0, 2, like that; in general, if it is 2 to the power r , I will require r bits. So, every index value n , 0, 1, 2, all the index values, up to capital N minus 1, like 15 here, or, if capital N is 64, it is 63 here, that will require r number of bits; if N is 64, 2 to the power 6. So, we will go up to 0, 1, up to 63, 0, 1, up to 63.

These, all these indices can require if you want to represent all of them, you require 6 bits. If N is 16, there is 0, 1, up to 15, there is index set, the small n goes about this. Then, if you want to represent each value of small n , time index, I will require 4 bits; 2 to the power 4 is 16. So, in general, I have got, suppose a 0, 1, ..., a_{r-1} , how many, r minus 1. So, how many bits? 0th bit, first bit, second bit, up to r minus 1 bit, total r . This is how my n is, and what is the binary value? a 0, because, the decimal value a 0, plus a 1 into 2, a 2 into 2 to the power and, these are all positive. I am not

representing a negative number here; negative is impossible; this is up to range. So, a_0 to the power 2, to the power 2, dot, dot, dot, 2 to the power r minus 1.

Now, suppose n is of even, n is even, in that case, if n is even, then, it is divisible by 2; that means, a_0 cannot be 1; a_0 has to be 0, because if $a_0 = 1$, this remaining part, you see, they are all 2, 2 to the power 2, 2 to the power 3, up to 2 to the power r minus 1. So, they are the, that is divisible by 2; but, if it is 1, and if you divide the entire thing by 2, there will be remainder 1. So, it is not be divisible by 2. So, n is even means, a_0 must be 0, and, n odd means, a_0 must be 1. Now, I am dividing this sequence into even component and odd component. So, what I will do, I will basically 0, 2, 4, 8, they will come under 1 umbrella; and, 1, 3, 5, that will come under 1 sub sequence; that is you remember h_n here, and g_n , 2 length N by 2 sequences.

But here, indices will be 0 to, in terms of x , x_0 , x_2 , that is, h_n is x^{2n} ; g_n is x^{2n+1} . So, it is, in terms of x , x_0 , x_2 , x_4 , x_6 , like that. So, n will be taking even numbers; that means I will take the N ; it is their binary representation of this form for various N s; whenever a_0 is 0, and it is even. So, I will put them in one half, and whenever a_0 is 1, it is odd, I will put them in another half; while a_0 is 0, I will consider 0, 0, 0, 0, 0, then, 0, 0, 0, 0, 1, 0; that is decimal value 2; 0, 0, 0, 1, 0, 0, that is decimal value 4. Then, 0, 0, 0, 0, 1, 1, 0, decimal value 6, like them. So, in that order, I will put them; same for the odd part. I think time is off for this class, I will continue from here, in the next.

Thank you very much.