

**Advanced Engineering Mathematics**  
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**Lecture - 41**  
**Discrete Fourier Transforms - II**

Hello friends. Welcome to my second lecture on Discrete Fourier Transform. We will first show that the vectors  $U_k$  given by  $e^{j2\pi kn/N}$  where  $n$  varies from 0 up to  $N-1$ , it takes value 0,1,2,3 and so on  $N-1$  they are orthogonal that is they form an orthogonal basis over the set of  $N$  dimensional complex vectors. So in order to prove that this vector  $U_k$  form an orthogonal basis we have to show that

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Orthogonality property

The vectors  $u_k = (e^{j2\pi kn/N} : n = 0, 1, 2, \dots, N-1)^T, k = 0, 1, 2, \dots, N-1$ , form an orthogonal basis over the set of  $N$ -dimensional complex vectors.

**Proof:** We have

$$\langle u_k, u_{k'} \rangle = \bar{u}_k^T u_{k'}$$

$$\bar{u}_k^T = \left( e^{-j2\pi k_0/N}, e^{-j2\pi k_1/N}, \dots, e^{-j2\pi k_{N-1}/N} \right)$$

$$u_{k'} = \begin{pmatrix} e^{j2\pi k'_0/N} \\ e^{j2\pi k'_1/N} \\ \vdots \\ e^{j2\pi k'_{N-1}/N} \end{pmatrix}$$

$$= \sum_{n=0}^{N-1} e^{j2\pi n(k'-k)/N}$$

$$= N\delta_{kk'}, \text{ where } \delta_{kk'} \text{ is Kronecker delta.}$$

$u_k = (e^{j2\pi kn/N})_{n=0}^{N-1} \in \mathbb{C}^N$   
 $u_{k'} = (e^{j2\pi k'n/N})_{n=0}^{N-1} \in \mathbb{C}^N$   
 $\bar{u}_k = (e^{-j2\pi kn/N})_{n=0}^{N-1} \in \mathbb{C}^N$

This orthogonal property is used to derive the formula for the IDFT from the definition of DFT.

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The inner product of  $U_k$  with; another vector say  $U_{k'}$  from the same set is equal 0. Now let us take two vectors  $U_k$  and  $U_{k'}$  from this set  $e^{j2\pi kn/N}$  where  $n$  takes value 0,1,2,3 and so on up to  $N-1$ . So then  $U_k$  is  $e^{j2\pi kn/N}$   $U_{k'} = e^{j2\pi k'n/N}$ . This  $N$  denotes the transpose. That means we are writing this vectors  $N$  vectors in the form of a column, okay.

So that is why we have written that as  $e^{j2\pi kn/N}$ , okay. They are  $N$  vectors, we writing them in the form of a column. So now inner product of  $U_k$  with  $U_{k'}$ ,  $U_k$  and  $U_{k'}$  belong to the  $N$  dimensional complex vector space that is  $\mathbb{C}^N$ , okay. So, they belong to  $\mathbb{C}^N$ ,

okay.  $U_k$  and  $U_k^*$  belong to  $C_n$  okay. So in  $C_n$  the inner product of  $U_k$  with  $U_k^*$  is defined as  $U_k^* \cdot U_k$ .

$U_k^*$  conjugate will be equal to  $e^{i k/n}$  where  $i$  will be replaced by  $-i$  so  $e^{-i k/n}$ , okay. So this will be  $U_k^*$  conjugate,  $U_k^* \cdot U_k$  conjugate transpose.  $U_k^*$  conjugate transpose means, we take the transpose of the vector. So, when we take the transpose of the vector, okay then  $U_k^*$  conjugate transpose will be row vector and that will be given by  $e^{i k/n}$ , okay where  $n$  takes values from 0,1,2,3 and so on up to  $n-1$ . So we can write like this.

Let us take  $n=0$  first, so we can write  $e^{i k_0/n}$  then  $e^{i k_1/n}$  - we have to take because of this  $-$ , okay so  $e^{-i k_1/n}$  second component and the last component is  $e^{i k_{n-1}/n}$ , okay. So this row vector,  $U_k^*$  conjugate transpose.  $U_k$  conjugate is the column vector actually, this column vector where we have  $n$  components. So let me write that as follows.

We write it as in the form of a column vector. So we have its component as  $e^{i k_0/n}$ ,  $e^{i k_1/n}$ ,  $e^{i k_{n-1}/n}$ ,  $n-1/n$  divided by  $n$ . So this is the column vector,  $U_k$  conjugate. When you take transpose of this conjugate vector actually these  $U_k$ 's are all belonging to  $C_n$ , okay they are belonging to  $C_n$ .  $k$  varies from; sorry  $n$  varies from 0 to  $n-1$ .

They all  $r$  vector in  $C_n$ , okay. So  $U_k^*$  conjugate transpose is this. When you multiply by  $U_k^*$   $U_k^*$  is what  $U_k^*$  is column vector, okay. So  $e^{i k_0/n}$ ,  $e^{i k_1/n}$  and so on to the power  $e^{i k_{n-1}/n}$ , okay. So when you multiply  $U_k^*$  conjugate transpose with  $U_k^*$  okay this being  $1/n$  matrix and this being  $n/1$  matrix when you multiply you get  $1/n$  matrix that is you get this summation.

$\sum_{n=0}^{n-1} e^{i k_n}$ , okay this 0,  $n$  varies from 0 to  $n-1$ , so that we write as  $n$ . So  $e^{i k_n}$  and that  $k_n = k/n$ , okay. So  $1/n \cdot n/1$  matrix gives you  $1/1$  matrix. And here when you sum this series, this is a geometric series, here you can see that, if  $k=k_n$  okay, if  $k=k_n$   $e^{i k_n}$  will be equal to 1. And this  $n$  is varying from 0 to  $n-1$  so this total sum will be then  $n$ , okay. So if  $k=k_n$  this is equal to; let us do it here.

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$$\begin{aligned}
 & \sum_{n=0}^{N-1} e^{2\pi i (k'-k)n/N} \\
 &= \frac{1 - (e^{2\pi i (k'-k)/N})^N}{1 - e^{2\pi i (k'-k)/N}}, \quad k \neq k' \\
 &= \frac{1 - e^{2\pi i (k'-k)}}{1 - e^{2\pi i (k'-k)/N}} = \frac{1-1}{1 - e^{2\pi i (k'-k)/N}} \quad \text{because } k, k' = 0, 1, 2, \dots, N-1 \\
 &= 0 \\
 &\text{If } k = k' \text{ then} \\
 &\sum_{n=0}^{N-1} e^{2\pi i (k'-k)n/N} = \sum_{n=0}^{N-1} 1 = N \\
 &\text{If } k, k' = 0, 1, 2, \dots, N-1 \\
 &\sum_{n=0}^{N-1} e^{2\pi i (k'-k)n/N} = N \delta_{kk'}
 \end{aligned}$$

$\delta_{kk'} = \begin{cases} 0, & k \neq k' \\ 1, & k = k' \end{cases}$

We have sigma n=0 to n-1 e to the power 2pi i k dash-k \* n/n, okay. So this is actually a geometric series, if k is not equal to k dash. If k; let us first find the sum when k is not equal to k dash. So this is equal to 1-e to the power 2pi i k dash - k raise to the; or divided by n raise to the power n, okay. The geometric ratio here is e to the power 2pi i k dash - k/n, okay.

And so this divided by e to the power n divided by 1-e to the power 2pi k dash - k/n if k is not equal to k dash, okay. And this will be equal to 1-e to the power 2pi k dash - k, all right divided by 1-e to the power 2pi i/ k dash - k/n. Now k and k dash are integers, so e to the power 2pi ik dash - k will be equal to 1, so we have 1-1/1-e to the power 2pi ik dash - k/n, okay. Because, k and k dash vary from 0 up to n-1 okay.

Now; so this is equal to 0. So when k is not equal to k dash sum of this series is 1 and if k=k dash then we have sigma n=0 to 2pi k dash - k/\*n/n=summation n=0 to n-1, 1, okay. So 1 will be summed n times and therefore this is equal to n, okay. So when k is not equal to k dash the value is 0, when k=k dash the value is n, so we can combine both the cases, okay.

So for all k, k dash varying from 0 up to n-1 we can say that, sigma n=0 to n-1 e to the power 2pi ik dash - k/n = n times delta kk dash, okay. Where delta kk dash is a Kronecker delta. Okay. And we know that delta kk dash = 1 when k is not equal to k dash and 0 when k=; sorry it is 0; it is 0

when  $k$  is not equal to  $k'$  and it is 1 when  $k = k'$ . Okay. So this is how we can find the value of this series, so we have this.  $N$  times  $\delta_{kk'}$ , where  $\delta_{kk'}$  is a Kronecker delta.

This orthogonal property; so now we can see that when  $k$  and  $k'$ , okay  $U_k$  and  $U_{k'}$  are any two vectors from this set of  $n$  vectors okay these  $n$  vectors are there, so from the set of  $n$  vectors then their inner product is 0 whenever  $k$  is not equal to  $k'$ , this means that these vectors form an orthogonal basis. So this orthogonal property of the  $e^{2\pi i j k / N}$  will be used to determine the inverse discrete Fourier transform using the definition of discrete Fourier transform. Let us see how we do this.

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**The inverse N-point DFT:**

Suppose  $u = \{u_j\}_{j=0}^{N-1}$  then  $U_k = \sum_{j=0}^{N-1} u_j e^{-2\pi i j k / N}$  we claim that

$$u_j = \frac{1}{N} \sum_{k=0}^{N-1} U_k e^{2\pi i j k / N} \text{ for } j = 0, 1, 2, \dots, N-1. \quad (1)$$

$u_j = \frac{1}{N} \sum_{k=0}^{N-1} U_k e^{2\pi i j k / N}$

(1) is called the inverse N-point discrete Fourier transform.

$$\begin{aligned} \sum_{k=0}^{N-1} U_k e^{2\pi i m k / N} &= \sum_{k=0}^{N-1} \left( \sum_{j=0}^{N-1} u_j e^{-2\pi i j k / N} \right) e^{2\pi i m k / N} \\ &= \sum_{j=0}^{N-1} u_j \left( \sum_{k=0}^{N-1} e^{2\pi i (m-j) k / N} \right) \\ &= \sum_{j=0}^{N-1} u_j N \delta_{mj} = N u_m, \quad m = 0, 1, 2, \dots, N-1 \end{aligned}$$

$= \sum_{j=0}^{N-1} u_j \langle u_m, u_j \rangle$

So let us arrive the formula for the inverse N-point DFT, okay. Suppose  $u = U_k$  or  $u_j$  where  $j$  varies from 0 to  $n-1$  then  $U_k$  is the discrete finite discrete transform; discrete finite transform of the vector  $u_j$ , so it is given by  $\sum_{j=0}^{N-1} u_j e^{-2\pi i j k / N}$ . We claim that, the inverse discrete Fourier transform IDFT is given by  $u_j = \frac{1}{N} \sum_{k=0}^{N-1} U_k e^{2\pi i j k / N}$  where  $j$  varies from 0 to  $n-1$ . So let us see how we prove this.

So we have  $U_k$ ,  $U_k$  is given to us, so  $u_j$  is the definition of DFT,  $U_k \sum_{k=0}^{N-1}$ , okay  $\sum_{k=0}^{N-1}$ .  $U_k$  will be  $\sum_{k=0}^{N-1}$  the value of  $U_k$  is  $\sum_{j=0}^{N-1}$ , okay  $u_j e^{-2\pi i j k / N}$ , okay  $e^{-2\pi i j k / N}$  that is  $U_k$  right  $\sum_{j=0}^{N-1}$ ;  $\sum_{j=0}^{N-1}$

to  $n-1$ ;  $u_j e$  to the power  $-2\pi i$  this. Now this can be written as; these are finite sums we can interchange them. So we can write as  $\sum_{j=0}^{n-1} u_j$ ,  $u_j$  can be written outside, okay,  $\sum_{k=0}^{n-1}$  okay,  $e$  to the power  $-2\pi i jk/n$ , okay.

Now, let us see, we have to show; now you okay; okay let us go to the previous slide. In the previous slide, we have shown that these vectors  $u$   $e$  to the power  $2\pi i$ ; this  $e$  to the power  $2\pi i k/n$  okay satisfies this orthogonal property that is  $U_k U_k^{\dagger}$  the other product of  $U_k$  with  $U_k^{\dagger}$  is equal to  $n \delta_{kk}$ . So here what will happen, here let us see; we will have  $\sum$  okay, we have to multiply  $U_k$  by  $2\pi i jk/n$  that we have not done.

So let me write this as; we can write like this so  $\sum_{k=0}^{n-1}$ , okay.  $U_k e$  to the power  $2\pi i jk/n$ , okay. So this is  $\sum_{k=0}^{n-1}$ , I have put the value of  $U_k$ , this is the value of  $U_k$ . I have to multiply by  $e$  to the power  $2\pi i jk$ ,  $j$  is already there. So what I should do, I should take some other index not  $j$ , let me take here  $m$ . Let us take here  $e$  to the power  $2\pi i$ , because  $j$  already I am using here so I should not write  $j$  here, so  $2\pi i m k/n$  let us write this. Okay. So  $e$  to the power  $2\pi i m k/n$ . Okay.

So, what I do, I consider this right hand side, okay. I consider  $\sum_{k=0}^{n-1}$ ,  $U_k e$  to the power  $2\pi i$ , instead of  $j$  I write  $m$ , okay. So  $2\pi i m k/n$ . Then this is equal to  $\sum_{k=0}^{n-1}$ , I put the value of  $U_k$  from here, okay.  $\sum_{j=0}^{n-1} u_j e$  to the power  $-2\pi i jk/n$  and then multiplied by  $e$  to the power  $2\pi i m k/n$ . Now let us see, I interchange the sums here because sums are finite, so  $\sum_{j=0}^{n-1}$ ;  $u_j$  I write outside then  $\sum_{k=0}^{n-1}$  to the power  $-2\pi i$ .

So I remove this here and write this as; I combine this in this and write  $e$  to the power  $2\pi i m-j * k/n$ . So I write like this. So I write  $e$  to the power  $2\pi i m-j * k/n$ , okay. So we have seen that, let us look at this one. We have seen that  $\sum_{n=0}^{n-1} e$  to the power  $2\pi i m k - k/n$  which is the inner product of  $U_k^{\dagger}$ , okay. So here we can say that we are having the enough product of  $U_m$  with  $U_j$  okay. So this is nothing but  $\sum_{j=0}^{n-1} u_j$  inner product of  $U_m$  with  $U_j$ . Okay. Okay, so this we have seen from the orthogonal property.

Now this is  $\sum_{j=0}^{n-1} u_j$  I write here =  $\sum_{j=0}^{n-1} u_j$  and then  $\delta_{mj}$ ;  $n$  times  $\delta_{mj}$ , okay. Because inner product of  $u_j$  is  $n$  times  $\delta_{mj}$  okay so this is equal to  $n$  times  $\sum_{j=0}^{n-1} u_j \delta_{mj}$ . And  $\delta_{mk}$  is the Kronecker delta. So it will be equal to 1 then  $m=j$  otherwise it is 0. So when this  $j$  will be equal to  $m$  this sum will be equal to  $U_m$  okay otherwise  $\delta_{mj}$  will be 0 for every other  $j$  which is not equal to  $m$ , so then other terms will be 0.

So this is  $n$  times  $U_m$ , okay. So we can see that  $\sum_{k=0}^{n-1} U_k e^{2\pi i k n/n}$  is  $= n$  times  $U_m$ , where  $n$  is taking any value from 0,1,2 up to  $n-1$ . Now we can replace  $m$  by  $j$  okay and then we have  $u_j$  from here  $u_j = 1/n$  okay. This  $n$  will go here, a  $\sum_{k=0}^{n-1} U_k e^{2\pi i k n/n}$  okay. So this is how we find the inverse  $N$ -point discrete Fourier transform. Okay.

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**Example**

Find  $N$ -point inverse DFT of  $\{X_k\}_{k=0}^{N-1}$  where



$$X_k = \delta[k - k_0] = \begin{cases} 1, & k = k_0 \\ 0, & \text{otherwise,} \end{cases}$$

for  $k_0 \in \{0, 1, 2, \dots, N-1\}$ .  
Then

$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k e^{2\pi i k n/N} = \frac{1}{N} e^{2\pi i k_0 n/N} \checkmark$$

$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} \delta(k-k_0) e^{2\pi i k n/N}$$

$$= \frac{1}{N} e^{2\pi i k_0 n/N}$$



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Now let us take example, find  $N$ -point inverse discrete Fourier transform of this frequency magnitude  $X_k$ ,  $k$  runs from 0 to  $n-1$ .  $X_k$  is given by  $\delta_{k-k_0}$ .  $\delta_{k-k_0}$  is equal to 1 when  $k=k_0$  otherwise it is 0. And  $k_0$  is any number from the set 0,1,2,3 and so on up to  $n-1$ , okay. Now let us use the  $N$ -point inverse DFT formula, okay. So  $N$ -point inverse DFT formula gives instead of  $u_j$  I am writing now  $x_n$  okay because our sequence is  $x_n$ .

For the sequence  $x_n$  we are using the discrete Fourier transform by  $X_k$ . So  $x_n = 1/n \sum_{k=0}^{n-1} U_k e^{2\pi i k n/n}$ , okay. So  $1/n \sum_{k=0}^{n-1} U_k e^{2\pi i k n/n}$ , okay. So

here, when you put the value of  $X_k$ . So  $X_n$ ,  $X_n$  will be equal to  $1/n \sum_{k=0}^{n-1} \delta_{k-k_0} e^{j 2\pi i k n/n}$ , okay. Now by definition of delta,  $\delta_{k-k_0}$  is 1 when  $k=k_0$ . So only the term where  $k$  becomes equal to  $k_0$  that term will survive other terms will all vanish.

So this will be  $1/n$ . And when  $k=k_0$  we have  $\delta_{k-k_0} = 1$ , so we get  $e^{j 2\pi i k_0 n/n}$ , okay. So we get this.

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**IDFT by matrix method**

In the matrix method, the IDFT can be calculated by using the matrix equation  $x = D_N^{-1} X$ , where  $x_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k e^{j 2\pi i k n/N}$ ,  $w_N = e^{j 2\pi i / N}$ ,  $X = D_N^{-1} x$

$$D_N^{-1} = \frac{1}{N} \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & w_N^{-1} & w_N^{-2} & \dots & w_N^{-(N-1)} \\ 1 & w_N^{-2} & w_N^{-4} & \dots & w_N^{-2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & w_N^{-(N-1)} & w_N^{-2(N-1)} & \dots & w_N^{-(N-1)(N-1)} \end{bmatrix}_{N \times N}$$

Handwritten notes:  $w_N = e^{-j 2\pi i / N}$ ,  $w_N = e^{-j 2\pi i k n / N}$

Okay, now we have discussed in the lecture, first lecture on DFT, how we can calculate DFT of a given sequence by using matrix method. Here also, the inverse discrete Fourier transform can be determined by using the matrix equation  $X = D_N^{-1} x$ . Okay, we had used the matrix equation  $X = D_N^{-1} x$  there. Okay. Here we have  $D_N^{-1}$  given by  $1/n$  \* this matrix, this matrix is  $n/n$  matrix okay.

So in the first row we have 1 1 1 each. Let us see how we get this matrix, we this sequence  $X_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k e^{j 2\pi i k n/N}$  let us use this one,  $X_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k e^{j 2\pi i k n/N}$  we have this sequence  $X_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k e^{j 2\pi i k n/N}$  okay,  $e^{j 2\pi i k n/N}$  okay. So you can see here, when this is  $k$  runs along the rows, okay.  $k$  runs along the rows, this is  $k=0$ ; this is  $k=1$ ; this is  $k=2$  and so on this  $k=N-1$ , okay,  $k=N-1$ . And your, we would be writing it as  $k n$ ; okay so  $n$  runs along the columns.

This is  $n=0, n=1, n=2$  and it is so on  $n=n-1$ . And we can write it as  $1/n \sum_{k=0}^{n-1} X_k W_n^{kn} - kn$ , okay. So we had; if you recall in the matrix method for DFT we had taken  $W_n$  to be equal to  $e$  to the power  $-2\pi i/n$ , okay,  $e$  to the power  $-2\pi i/n$ , so here  $W_n$  to the power  $-k$  and  $W_n$  to the power  $-kn$  will be  $W_n^{-kn}$  will be equal to  $e$  to the power  $2\pi i kn/n$ , okay. So we use same  $W_n$  here,  $W_n = e$  to the power  $2\pi i/n$ , okay.

So this is  $k=0, n=0$ ,  $W_n$  to the power 0 will be 1, so when  $k=0$   $W_n$  to the power  $-kn$  it will be 1, so we get this 1 here. And then when  $k=1$  okay, but  $n=$ ; yeah when  $k$  is taking any value from 0 to  $n-1$  but  $n=0$  again  $W_n = e$  to the power 0 will be equal to 1, so we get column first column also 1. This column will be  $k=1, n=1$ , so  $W_n$  to the power  $-1*1$  that is  $W_n$  to the power  $-1$ . Then  $k=1, n=2$ , so  $W_n$  to the power  $-2$  and so on.

Then  $W_n$  to the power  $-n-1$  when  $k=1, n=n-1$ . And similarly, we get the last row. So  $1/n$ , this  $1/n$  comes here, okay.  $X_k$  is the column vector, okay.  $X_0, X_1, X_2$  and so on  $X_{n-1}$  and this matrix  $W_n$  to the power  $-kn$ , this matrix is here, okay. So  $D_n$  inverses this.  $D_n$  inverse will be multiplied by  $X$  to get the sequence  $X$ .

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The Plancher theorem and Parseval's theorem



If  $X_k$  and  $Y_k$  are the DFTs of  $x_n$  and  $y_n$ ,  $n = 0, 1, 2, \dots, N-1$ , then the Plancher theorem states that

$$\sum_{n=0}^{N-1} x_n \bar{y}_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k \bar{Y}_k$$

**Proof:** By definition  $X_k = \sum_{n=0}^{N-1} x_n e^{-2\pi i kn/N}$  and  $Y_k = \sum_{m=0}^{N-1} y_m e^{-2\pi i km/N}$ .  
Hence,

$$\frac{1}{N} \sum_{k=0}^{N-1} X_k \bar{Y}_k = \frac{1}{N} \sum_{k=0}^{N-1} \left( \sum_{n=0}^{N-1} x_n e^{-2\pi i kn/N} \right) \left( \sum_{m=0}^{N-1} \bar{y}_m e^{2\pi i km/N} \right)$$

$Z_n = \sum_{k=0}^{N-1} X_k e^{2\pi i kn/N}$   
 $Y_n = \sum_{k=0}^{N-1} Y_k e^{2\pi i kn/N}$   
 $\bar{Y}_k = \sum_{m=0}^{N-1} \bar{y}_m e^{2\pi i mk/N}$



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Now let us discuss the Parseval's theorem. Plancher theorem and Parseval's theorem. Suppose  $X_k$  and  $Y_k$  are discrete Fourier transforms of  $x_n$  and  $y_n$ , okay. So  $X_n = \sum_{k=0}^{n-1} X_k e$  to the power  $2\pi i kn/n$ . And by  $n=$  similarly,  $\sum_{k=0}^{n-1} k e$  to the power  $2\pi i kn/n$ . Okay. So



those  $X_k Y_k$  are discrete Fourier transforms of the sequence  $X_n Y_n$ . Then the Plancherel theorem states that  $\sum_{k=0}^{N-1} X_k Y_k^* = \sum_{n=0}^{N-1} X_n Y_n^*$ . Okay.

So  $X_n Y_n^*$ ,  $X_n$  you multiply by  $Y_n^*$ , okay. What we will get? See,  $X_k$  from the definition of; we are taking the right hand side to prove this equation, to prove this theorem. So by the definition of discrete Fourier transform  $X_k$  is given by  $\sum_{n=0}^{N-1} X_n e^{-2\pi i n k / N}$  and  $Y_k$  will be given by  $\sum_{m=0}^{N-1} Y_m e^{-2\pi i m k / N}$ , okay. So let us take the right hand side of this equation.

So  $\frac{1}{N} \sum_{k=0}^{N-1} X_k Y_k^*$ , okay will be  $\frac{1}{N} \sum_{k=0}^{N-1}$  then you put the value of  $X_k$  you put the value of  $Y_k^*$ ,  $Y_k^*$  conjugate will be conjugate here, okay  $Y_k^* = \sum_{m=0}^{N-1} Y_m^* e^{2\pi i m k / N}$ , so we get  $\frac{1}{N} \sum_{k=0}^{N-1} \sum_{n=0}^{N-1} X_n Y_m^* e^{2\pi i (m-n)k / N}$ . Okay. So this is what we get and this, okay. Let us now see what happens to this.

**(Refer Slide Time: 27:06)**

The Plancherel theorem and Parseval's theorem cont...

$$\begin{aligned}
 &= \frac{1}{N} \sum_{k=0}^{N-1} \left[ \sum_{n=0}^{N-1} \sum_{m=0}^{N-1} X_n Y_m^* e^{2\pi i (m-n)k / N} \right] \\
 &= \frac{1}{N} \sum_{n=0}^{N-1} \sum_{m=0}^{N-1} X_n Y_m^* \left( \sum_{k=0}^{N-1} e^{2\pi i (m-n)k / N} \right) \\
 &= \frac{1}{N} \sum_{n=0}^{N-1} \sum_{m=0}^{N-1} X_n Y_m^* (N \delta_{mn}), \quad (\text{by orthogonality property}), \\
 &= \sum_{n=0}^{N-1} X_n Y_n^*.
 \end{aligned}$$

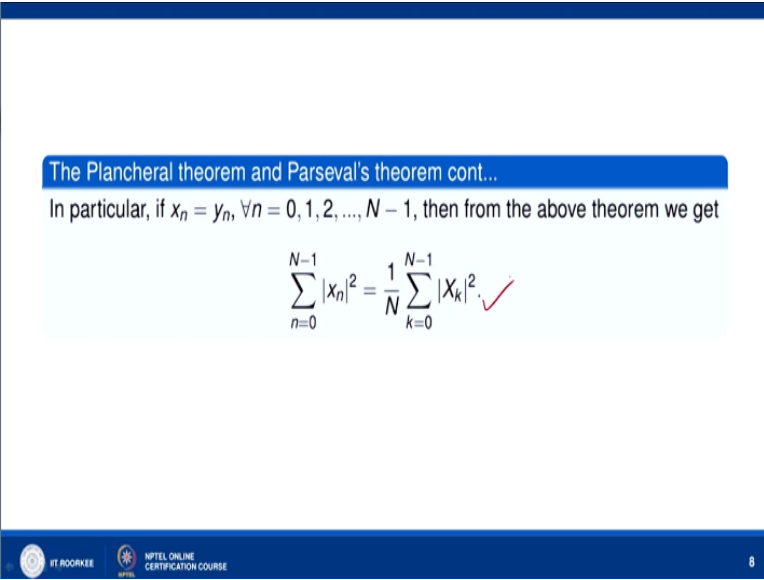
*Handwritten notes:  $\langle Y_m, Y_n \rangle = N \delta_{mn}$  and  $N \delta_{mn}$  with an arrow pointing to the sum over k.*

So  $\frac{1}{N} \sum_{k=0}^{N-1}$ , we can write it as  $\frac{1}{N} \sum_{k=0}^{N-1}$ , we put the two sigma together,  $\sum_{n=0}^{N-1} \sum_{m=0}^{N-1}$  then  $X_n Y_m^*$  okay  $X_n Y_m^* e^{2\pi i (m-n)k / N}$ , okay. So  $e^{2\pi i (m-n)k / N}$ , okay. Now, so now these are all finite sums, okay. We can interchange and write  $\frac{1}{N} \sum_{n=0}^{N-1} \sum_{m=0}^{N-1}$  the summation over  $k$  can be brought inside. So  $X_n Y_m^* \sum_{k=0}^{N-1} e^{2\pi i (m-n)k / N}$ .

Now this is again by using orthogonal property. This is nothing but in a product of sequence  $Y_m$  and  $Y_n$ , so this is equal to  $n$  times  $\delta_{mn}$ , okay. So we get  $\frac{1}{n} \sum_{n=0}^{n-1} \sum_{m=0}^{n-1} X_n Y_m \bar{\phantom{x}} \delta_{mn}$ , okay. So  $\delta_{mn}$  is Kronecker delta it will be equal to 1 when  $m=n$  otherwise it will be 0. So all those values of  $m$ , okay which are equal to  $n$  they give the value of  $\delta_{mn}$  as 1, okay otherwise it is 0.

So this  $n$ , this  $n$  will cancel with this  $n$  and  $X_n Y_m \bar{\phantom{x}} \delta_{mn}$  with this double summation to reduce to single sum, okay. So  $\sum_{n=0}^{n-1} X_n Y_n \bar{\phantom{x}}$ , okay when  $m=n$ . So this is how we get the left hand side. Okay. This left hand side. So this prove the Plancherel theorem. Let us see how we get the Parseval's theorem from here.

**(Refer Slide Time: 28:56)**



The Plancherel theorem and Parseval's theorem cont...

In particular, if  $x_n = y_n, \forall n = 0, 1, 2, \dots, N-1$ , then from the above theorem we get

$$\sum_{n=0}^{N-1} |x_n|^2 = \frac{1}{N} \sum_{k=0}^{N-1} |X_k|^2 \quad \checkmark$$

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So we go to particular case here, let us take the sequences  $X_n$  and  $Y_n$  to be same okay.  $X_n = Y_n$  for all  $n$ . So when  $X_n = Y_n$  for  $n=0, 1, 2, 3$  and so on up to  $n-1$  then from the Plancherel theorem what do we notice,  $X_k Y_k \bar{\phantom{x}}$  will be  $X_k \bar{\phantom{x}}$  now because at  $Y_n$  and  $X_n$  are sequences are same, so  $X_n X_k \bar{\phantom{x}}$ ,  $X_n X_k \bar{\phantom{x}}$  will be mod of  $X_k$  square. And here  $X_n \bar{\phantom{x}}$  will be mod of  $X_n$  square, so we have this result.  $\sum_{n=0}^{n-1} \text{mod of } X_n \text{ square}$  is equal to  $\sum_{k=0}^{n-1} \frac{1}{n} \text{ times this, okay} = \text{to this yeah.}$

So we have  $\sum_{n=0}^{n-1} \text{mod of } X_n \text{ square} = \frac{1}{n} \sum_{k=0}^{n-1} \text{mod of } X_k \text{ square, okay.}$   
So we get this result. So this is Parseval's theorem.

(Refer Slide Time: 29:56)

**Shifting property**

If  $F(\{x_n\})_k = X_k$  then  $F(\{x_n e^{-2\pi i m n / N}\})_k = X_{k-m}$ , where  $m$  is some integer.

**Proof:** Since  $F(\{x_n\})_k = X_k$ , we have

$$X_k = \sum_{n=0}^{N-1} x_n e^{-2\pi i k n / N}$$

Hence,

$$F(\{x_n e^{-2\pi i m n / N}\})_k = \sum_{n=0}^{N-1} x_n e^{2\pi i m n / N} e^{-2\pi i k n / N}$$

$$= \sum_{n=0}^{N-1} x_n e^{-2\pi i (k-m) n / N}$$

$$= X_{k-m}$$

Handwritten notes on the right side of the slide:

$$D(\{x_n\})^{(k)} = X_k$$

$$F(\{x_n\})^{(k)} = D(\{x_n\})^{(k)} = X_k$$

So if  $f(X_n)_k = X_k$  that means discrete Fourier transform of the sequence  $X_n = X_k$ . Earlier, when we started we had written the notation for the discrete Fourier transform as  $D$ ,  $D$  of  $X_n$   $k = X_k$  we had used earlier, but here in this Shifting property  $D$  is replaced by  $F$ , so there is just change of notation but it is the same thing. So  $f(X_n)_k$  is same as  $D$  of  $X_n$   $k$  it is equal to  $X_k$ . So then this Shifting property says that if you multiply  $X_n$  by  $e^{-2\pi i m n / N}$  then the discrete Fourier transforms shift to the right okay it becomes  $X_{k-m}$  where  $m$  is some integer.

So let us use the definition of discrete Fourier transform we can see that this  $X_k$  is not given by  $\sum_{n=0}^{N-1} x_n e^{-2\pi i k n / N}$ . Now multiply  $x_n$  by  $e^{2\pi i m n / N}$  okay. So we will replace here  $x_n$  by  $x_n e^{2\pi i m n / N}$  to get the discrete Fourier transform of this sequence. So we have  $\sum_{n=0}^{N-1} x_n e^{2\pi i m n / N} e^{-2\pi i k n / N}$  and this I can write as  $\sum_{n=0}^{N-1} x_n e^{-2\pi i (k-m) n / N}$ , okay.

And by definition of discrete Fourier transform, this is discrete Fourier transform of  $x_n$  sequence. Okay. So here this  $k$ , now corresponding to this  $k$  we have  $X_k$  here, okay. So here this  $k$  is replaced by  $k-m$ , okay so we have the discrete Fourier transform  $X_{k-m}$ , okay. So when you multiply  $x_n$  by  $e^{2\pi i m n / N}$  the discrete Fourier transform shifts to the right, it becomes  $X_{k-m}$ .

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**Another shifting property**

Let  $D(\{x_n\})(k) = X_k$  then  $D(\{x_{n-m}\})(k) = \underline{X_k e^{-2\pi i k m / N}}$ .

**Proof:** We have,

$$\begin{aligned}
 R.H.S &= \left( \sum_{n=0}^{N-1} x_n e^{-2\pi i k n / N} \right) e^{-2\pi i k m / N} \\
 &= \sum_{n=0}^{N-1} x_n e^{-2\pi i k (n+m) / N} \quad n+m=j \\
 &= \sum_{j=m}^{N+m-1} x_{j-m} e^{-2\pi i k j / N} \\
 &= \sum_{n=m}^{N+m-1} x_{n-m} e^{-2\pi i k n / N} = \underline{D(\{x_{n-m}\})(k)}.
 \end{aligned}$$

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Now another shifting property. Suppose the  $X_n(k) = X_k$  then the  $X_{n-m}(k)$  the  $X_{n-m}$  is equal to  $X_k e$  to the power  $-2\pi i k m / N$ , okay. Let us see how we get this. So we use this right hand side, right hand side by definition of the discrete Fourier transform it is  $\sum_{n=0}^{N-1} x_n e^{-2\pi i k n / N}$  and we are multiplying it by  $e^{-2\pi i k m / N}$ . So the summation is over  $n$  here so I can bring it inside,  $\sum_{n=0}^{N-1} x_n e^{-2\pi i k n / N} e^{-2\pi i k m / N}$  and we get  $\sum_{n=0}^{N-1} x_n e^{-2\pi i k (n+m) / N}$ , okay. And so let us now repeat  $n+m=j$ .

So when we write  $n+m=j$  then earlier was equal to 0, so now  $j$  will be equal to  $m$  and  $n$  was equal to capital  $N-1$  so  $j$  will be  $n+m-1$ , so  $n$  becomes  $j-m$  so  $x_{j-m} e^{-2\pi i k j / N}$ . Okay. And therefore, this is let us replace  $j/n$  so  $n=m$  to  $n+m-N$   $x_{n-m} e^{-2\pi i k n / N}$ . So this is nothing but the discrete Fourier transform of the sequence  $x_{n-m}$ .

(Refer Slide Time: 34:01)

**Convolution theorem**

Let  $x = \{x_n\}_{n=0}^{N-1}$  and  $y = \{y_n\}_{n=0}^{N-1}$  be two sequences of complex numbers then their convolution is defined as

$$(x * y)_n = \sum_{m=0}^{N-1} x_m y_{n-m \pmod N}$$

Then

$$D^{-1}(X_n Y_n) = \sum_{m=0}^{N-1} x_m y_{n-m \pmod N} = (x * y)_n$$

*Handwritten notes:*  
 $D((x * y)_n) = X_n Y_n$   
 $= \text{Product of the DFT of the sequences } x \text{ \& } y$

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Now Convolution theorem. Suppose  $X = X_n$   $n=0$  to  $n=N-1$  one sequence  $Y_n$   $n=0$  to  $n=N-1$  is the other sequence of complex numbers then their convolution is defined as  $x * y$   $n$  okay.  $\sum_{m=0}^{N-1} x_m y_{n-m \pmod N}$ , okay. This theorem says that, the discrete Fourier transform of this convolution of the sequence  $X_n$  and  $Y_n$ , okay.  $X$  and  $Y$ , this is the discrete Fourier transform of the convolution of the sequence  $X$  and  $Y$  which we denote by this notation is equal to  $\sum_{m=0}^{N-1} x_m y_{n-m \pmod N}$  okay.

So which is nothing but; yeah when you take the inverse discrete transform it is  $x_m * y_n$ .  $x_m$  is the discrete Fourier transform of  $X$  and  $Y_n$  is the discrete Fourier transform of the sequence  $Y$ . So discrete Fourier transform of, so we can otherwise, I mean in other words we can say that discrete Fourier transform of  $x * y$  okay, this equal to  $X_n * Y_n$  that is product of the DFT of the sequence, sequences  $X$  and  $Y$ . Okay.

So discrete Fourier transform of the convolution of the two sequences  $X$  and  $Y$  is equal to product of their discrete transforms, so this is; and I repeat the convolution of the sequences  $X$  and  $Y$  is defined as  $\sum_{m=0}^{N-1} x_m y_{n-m \pmod N}$  and then modulo  $N$ , modulo  $N$  is used so that the values of  $Y_n$  do not go beyond this  $N-1$  they remain in between  $0$  and  $N-1$ , okay.

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### Convolution theorem cont...

**Proof:** We shall show that

$$D\left(\sum_{m=0}^{N-1} x_m y_{n-m(\text{mod } N)}\right) = X_n Y_n,$$

where  $X_n$  and  $Y_n$  are the DFT of  $\{x_n\}_{n=0}^{N-1}$  and  $\{y_n\}_{n=0}^{N-1}$ .  
L.H.S

$$\begin{aligned} &= \sum_{n=0}^{N-1} \left( \sum_{m=0}^{N-1} x_m y_{n-m(\text{mod } N)} \right) e^{-2\pi i k n / N} \\ &= \sum_{m=0}^{N-1} x_m \left( \sum_{n=0}^{N-1} y_{n-m(\text{mod } N)} e^{-2\pi i k n / N} \right) \\ &= \sum_{m=0}^{N-1} x_m \left( \sum_{j=0}^{N-1} y_j e^{-2\pi i k (m+j) / N} \right) \end{aligned}$$

$$\begin{aligned} &= \sum_{j=0}^{N-1} y_j e^{-2\pi i k j / N} \left( \sum_{m=0}^{N-1} x_m e^{-2\pi i k m / N} \right) \\ &= Y_k X_k \end{aligned}$$

Now let us prove this. So discrete Fourier transform of, we are going to show that discrete Fourier transform of the convolution of the sequence  $X$  and  $Y$  = product of their discrete transforms. Okay. So let us take this left hand side, okay. Left hand side will be equal to sigma  $m=0$  to  $n-1$ ; okay sigma  $m=0$  to  $n-1$  this convolution okay we are writing here  $m=0$  to  $n-1$   $X_m Y_{n-m}$  modulo  $N$  and then multiplying by  $e$  to the power  $-2\pi i k n / n$ , okay.

So this is simply discrete Fourier transform of the convolution of  $X$  and  $Y$ . And then what we do; these are finites sums so we can interchange them, so sigma  $m=0$  to  $n-1$  this is written at outside and inside we write sigma  $m=0$  to  $n-1$   $X_m Y_{n-m}$  modulo  $N$   $e$  to the power  $-2\pi i k n / n$ . Okay. Now then, this  $X_m$  this  $X_m$  is independent of  $N$  so I can write it outside the summation over  $N$  and what we will get; then we will get, yeah so sigma  $m=0$  to  $n-1$   $X_m$  and this sigma  $m=$ ; we are interchanged, we have already interchanged okay.

So  $X_m$  will go outside the summation over  $n-1$ . Now we have sigma  $n=0$  to  $m-1$  by  $n-m$  okay modulo  $n$  into  $e$  to the power  $-2\pi i k n / n$ . So because of this modulo this is equal to sigma  $n=m$  to  $n+m-1$ , okay sigma  $n=m$  to  $n+m-1$  by  $n-m$  okay,  $e$  to the power  $-2\pi i k n / n$ , okay because of this modulo okay I can write this expression as sigma  $n=m$  to  $n+m-1$  by  $n-m$ . Now let me write  $n-m=j$ , okay.

So I will be getting  $j=0$  to  $n-1$  by  $j$  and here  $n$  will be  $n+j$  so  $e$  to the power  $-2\pi i m+j * n/n$ , okay which we can write in as follows.  $\sum_{j=0}^{n-1} e$  to the power  $-2\pi i m/n$  okay. So thus we write  $e^{jn/n} * Y_j$ . Okay, this  $Y_j$ . This is one thing multiplied by  $e$  to the power  $-2\pi i m/n$ . Okay. So this quantity  $e$  to the power  $-2\pi i m/n$  will be multiplied to this  $X_m$  which is outside the summation over  $n$ , okay, so this here, okay.

And we will get  $\sum_{m=0}^{n-1} X_m e$  to the power  $-2\pi i mn/n$ , okay,  $e$  to the power  $-2\pi i mn$ ; yeah you can write  $K$  okay what we have put here, we have put  $j$  right. So when we put  $j$  here, okay so I get here this sum as  $\sum_{j=0}^{n-1} Y_j e$  to the power  $-2\pi i kj$ . Where we getting  $kj$ ? Okay, this was; this is outside, okay. Inside, okay we are getting here  $2\pi i kn/n$ . And what was  $(( ))$  (40:49)?

Okay, so, I put here okay wait I put here of this not  $n$ ,  $n$  will be put as  $j$  okay. So this should be  $k$ , okay. So this should be  $k$  here not  $n$ , okay. So what we have here, actually this is now coming to this, this is nothing but  $Y_k$ , okay this is  $Y_k$  and this quantity okay multiplied to  $X_m$   $\sum_{m=0}^{n-1}$  gives you  $X_k$ , okay so we get this. This and this, okay their product is  $X_k Y_k$ . Thus the convolution of two sequences is the inverse transform of the product of the individual transforms. So this is how we prove the convolution theorem.

With this we have finished the discussed on discrete Fourier transforms. Thank you very much for your attention.