## INDIAN INSTITUTE OF TECHNOLOGY ROORKEE

## NATIONAL PROGRAMME ON TECHNOLOGY ENHANCED LEARNING (NPTEL)

#### **Discrete Mathematics**

Module-04 Discrete Probability Lecture-04 Information and mutual information

## With Dr. Aditi Gangopadhyay Department of Mathematics IIT Roorkee

The topic of the lecture is information and mutual information suppose a die has phases 1 and 4 with red color and phases 2, 3, 5 and 6 with black color now suppose this kind of die was rolled and we were told that the outcome was 4 clearly we were given all the information concerning the outcome of the experiment if we were told that the outcome was red we would agree that we were given some information but not all the outcome is narrowed to one of two possibilities on the other hand if we were told that the outcome was black we would feel that we were given some information.

But even less the outcome is narrowed to one of four possibilities another example suppose after 6 weeks of the semester students were told that there will be a one-hour examination clearly such an announcement contains a certain amount of information however if the students were told after 1 week of classes that there will be one hour examination we would say the announcement contains much more information because it is quite honest unexpected that an examination would be she jeweled after only one week of classes so this is a surprising statement so that is why it will contain more information.

So from these examples it is clear that it is very important to measure quantitatively how much information a certain piece of message carries if a statement tells us the occurrence of a certain event that is likely to happen we would say that the statement contains a small amount of information on the other hand if a statement tells us the occurrence of a certain event that is not likely to happen we would say that the statement contains a large amount of information.

So this observation suggests that the information contained in a statement asserting the occurrence of an event depends on the probability of occurrence of the event the information content of a statement asserting the occurrence of an event is - log P.

(Refer Slide Time: 04:54)

The information content of a statement asserting the occurrence of an event is -log P

where P is the probability of occurrence of that event. Since P is always less than or equal to 1, -log P is always positive. It is also clear that the smaller the value of P, the larger the quantity -log P. This implies if the probability of occurrence of an event is less the statement asserting the occurrence of an event will contain large amount of information.

\* Thus if we now consider the example by which we started, when we were told that the outcome of rolling a die was 4, the amount of information we received

 $= -\log \frac{1}{6} = \log 6 = 2.585.$ 

Where P is the probability of occurrence of that event since P is always less than or equal to  $1 - \log P$  is always positive it is also clear that the smaller the value of P the larger the quantity - log P this implies if the probability of occurrence of an event is less the statement asserting the occurrence of an event will contain large amount of information thus if we now consider the example by which we started when we were told that the outcome of rolling a die was for the amount of information we received =- log 1 / 6 and that is log 6 which is 2 . 585.

(Refer Slide Time: 06:16)

On the other hand, when we were told that the outcome was red, the amount of information we received

$$= -\log\frac{2}{6} = \log 3 = 1.585$$

Example : Suppose that we receive from the computer as output a binary digit that is either 0 or 1 with equal probability of occurrence. When we are told that the output is indeed 1, the amount of information we receive is  $-\log \frac{1}{2} = 1$ 

Similarly, when we are told that the output is 0, the amount of information we receive is also,

$$-\log \frac{1}{2} = 1$$

On the other hand when we were told that the outcome was red, the amount of information we received =  $-\log 2 / 6$  because if red is red has occurred then the probability will become 2 / 6 because there are 2 possibilities only 1 and 4 so that is why number of favorable cases will become 2 and that is why the probability will become 2 / 6 and as a result of this amount of information will be  $-\log 2 / 6$  and that will be log 3 which is 1. 585 another example suppose that we receive from the computer as output a binary digit that is either 0 or 1 with equal probability of occurrence.

When we are told that the output is indeed 1 the amount of information we receive is -  $\log \frac{1}{2}$  because there are only 2 possibilities 0 and 1 so that is why probability will be 1 / 2 and the amount of information will become -  $\log \frac{1}{2}$  that is 1 similarly when we are told that the output is 0 the amount of information we received is also -  $\log \frac{1}{2}$  so it will be 1 now suppose we receive 32 binary digits from the computer as output assuming all 2,  $2^{32}$  possibilities are equally likely the information we receive is -  $\log \frac{1}{2}^{32}$ .

(Refer Slide Time: 08:53)

Now, suppose we receive 32 binary digits from the computer as output. Assuming all 2<sup>32</sup> possibilities are equally likely, the information we receive is

$$-\log \frac{1}{2^{32}} = 32 \ b \ its$$

At this stage, I would like to introduce the notion of mutual information.

Suppose we were told that the outcome of rolling a die is red. How much does that help us to determine that the outcome is a 4? Suppose we were told that the professor will be out of town tomorrow. How much does that help us to determine that there will be a 1- hour examination tomorrow?

So there are all total  $2^{32}$  possibilities so the probability and there you equivalence oh so that is why the probability will become  $\frac{1}{2}$   $^{32}$  so the amount of information will become - log 1 /2<sup>32</sup> and that is 32 bits so the amount of information we are getting 32 bits at this stage I would like to introduce the notion of mutual information suppose we were told that the outcome of rolling a die is red how much does that help us to determine that the outcome is a 4 suppose we were told that the professor will be out of town tomorrow how much does that help us to determine that there will be one hour examination tomorrow.

So this kind of questions can be tackled by mutual information thus we want to know the amount of information concerning.

(Refer Slide Time: 10:43)

Thus, we want to know the amount of information concerning the occurrence of event A that is contained in the statement asserting the occurrence of event B, which we shall denote I(A,B). Since  $-\log P(A)$  is the amount of information contained in a statement asserting the occurrence of event A and  $-\log P(A/B)$  is the amount of information contained in a statement asserting the occurrence of A given that B has occurred the difference between these two quantities is the amount of information that B has occurred.

The occurrence of event A that is contained in the statement asserting the occurrence of event B which we shall denote by I (A, B) since -  $\log P(A)$  is the amount of information contained in a statement asserting the occurrence of event A and -  $\log P(A/B)$  is the amount of information contained in a statement asserting the occurrence of A given that B has occurred the difference between these two quantities is the amount of information on the occurrence of A provided by the assertion that B has occurred.

(Refer Slide Time: 12:01)

In other words, we need  $-\log P(A)$  bits of information to assert the occurrence of event A, and we still need  $-\log P(A/B)$  bits of information to assert the occurrence of event A after we were told that event B has occurred. Thus, the information provided by the occurrence of event B on the occurrence of event A is  $I(A, B) = [-\log P(A)] - [-\log P(A/B)]$  $= -\log P(A) + \log P(A/B).$ 

In other words we need - log P(A) bits of information to assert the occurrence of event A and we still need - log P(A/B) bits of information to assert the occurrence of event a after we were told that even B has occurred thus the information provided by the occurrence of event B on the occurrence of event A is I(A,B) that is I (A, B) that =- log P(A) that is - log P(A)-[- log P(A/B)] so this is the conditional probability so this =- log P(A) + log P(A/B).

(Refer Slide Time: 13:31)

For example, Let A be the event that 4 appeared and B be the event that red appeared when a die was rolled. Then  $I(A, B) = -\log P(A) + \log P(A / B)$  $= -\log \frac{1}{2} + \log \frac{1}{2}$ 

= 2.585 - 1

For example let a be the event that 4 appeared and B be the event that red appeared when a die was rolled then I(A, B) =- log P(A) + log P(A/B) so this will be = - log 1 / 6 because here the event A is that 4 has appeared so there are 6 possibilities 1, 2, 3, 4, 5,6 among that 4 has occurred with probability 1 / 6 and log P(A/B) what is this it is given that red has appeared so this event is given so in that case we will have only 2 possibilities 1 and 4 so that is why the probability that A/ B will become 1/2. So this will be log  $\frac{1}{2}$  so we will get 2.585 - 1 and this value will be 1.585 bits.

(Refer Slide Time: 15:18)

Now, if we replace the event B by the event C that an even number appeared. Then  $I(A, C) = -\log P(A) + \log P(A/C)$  $= -\log \frac{1}{6} + \log \frac{1}{3}$ = 2.585 - 1.585= 1 bit.

Now if we replace the event B by the event C that an even number appeared then what will happen, so B is replaced by C so that is why we have to find I(A, C) and that =- log  $P(A) + \log P(A/C)$  so - log P(A) that is same that is - log 1 / 6 + now what is P(A/C) if it is given that even number has occurred so that means there are only 3 possibilities 2,4 and 6 so in that case 4 will occur with probability 1 / 3 so that is why this will become log 1 / 3 so we will get 2.585 - 1.585 which is 1 bit.

(Refer Slide Time: 16:48)

So, from this we can observe that if P(A/B) is large, it means that the occurrence of *B* indicates a strong possibility of the occurrence of *A*. Consequently, I(A,B)is large. However, if P(A/B) is small, it means that the occurrence of *B* does not tell us much about the occurrence of *A*. Consequently I(A,B) is small. As a matter of fact, the occurrence of event *B* may mean that event *A* is less likely to occur. In that case P(A/B) is smaller than P(A) and I(A,B) is a -ve quantity.

So from this we can observe that if P(A/B) is large it means that the occurrence of B indicates a strong possibility of the occurrence of a consequently I(A, B) is large however if P(A/B) is small it means that the occurrence of B does not tell us much about the occurrence of A consequently I (A,B) is small as a matter of fact the occurrence of event B may mean that event A is less likely to occur in that case P(A/B) is smaller than P(A) and I (A, B) becomes negative.

Let us also examine some extreme cases suppose that B is a subset of A in S in that case intuitively the occurrence of B assures the occurrence of a since we have  $P(A \cap B) = P(B)$  it follows that P(A/B) = 1 and  $-\log P(A/B) = 0$  that is the mutual information provided by the assertion that B has occurred on the occurrence of A =the information provided by the assertion that A has occurred however suppose that B is the whole sample space in that case  $P(A \cap B) =$ P(A) and  $-\log P(A/B) = -\log P(A)$ .

(Refer Slide Time: 20:02)

However, suppose that *B* is the whole sample space. In that case,  $P(A \cap B) = P(A)$  and  $-\log P(A/B)$  is equal to  $-\log P(A)$ . Indeed that, I(A,B) = 0 means that occurrence of *B* tells us nothing about occurrence of *A*. Example: consider the problem of estimating the likelihood that there will be a 1 - hour examination when the professor is scheduled to go out of town. Let  $S = \{x_1, x_2, x_3, x_4\}$  be the sample space, where the samples represent the four possible outcomes.

And I (A, B) =0 and this means that occurrence of B tells us nothing about the occurrence of A so here we take one example consider the problem of estimating the likelihood that there will be one hour examination when the professor is scheduled to go out of town let s is the set containing  $x_1$ ,  $x_2$ ,  $x_3$  and  $x_4$  this is the sample space where the samples represent the 4 possible outcomes.

(Refer Slide Time: 21:11)



Now this can occur in 4 possibilities only professor is out of town and examination is given so that is considered as event  $x_1$  probability professor or out of state out of town and examination not given so that is considered as  $x_2$  professor in town and examination given that is  $x_3$  and  $x_4$  represents professor is in town and examination is not given so there are 4 cases which are possible here now another thing we are considering the probability is also probabilities of  $x_1$ ,  $x_2$ ,  $x_3$ ,  $x_4$  probability of  $x_1$  is considered as  $\frac{1}{2}$  probability of  $x_2$  is 1 / 16 probability of  $x_3$  is 3 / 16 probability of  $x_4$  is 1 / 4.

Let a denote the event that an exam is given and B the event that the professor is out of town note that  $P(A) = \frac{1}{2} + \frac{3}{16}$  so P(A) means probability that the that an exam is given probability what is the probability that an exam is given so that probability is P(A) and we know that we have 2 cases here  $x_1$  and  $x_3$  so that is why we have to add probability of  $x_1$  and probability of 3.

So that is why it is  $\frac{1}{2} + \frac{3}{16}$  which is 11 / 16 now next we are finding the conditional P(A/B) so that is P(A  $\cap$  B)/ P(B) now what is P(B) B is the event that professor is out of town and that is occurring in 2 places x<sub>1</sub> and x<sub>2</sub> so that is why to find the P(B) we have to add the probability of x<sub>1</sub> and probability of x<sub>2</sub> so that is here in the denominator and in the numerator it is  $\frac{1}{2}$  because it is P(A  $\cap$  B) that is both A and B will occur.

So here A and B that is exam is given and professor is out of town both these events is occurring here so probability of  $x_1$  we have to consider in the numerator and if we simplify this we will get the value 8 / 9.

(Refer Slide Time: 25:17)

The information needed that an exam will be given is  $-\log P(A) = -\log \frac{11}{16} = -\log 11 + \log 16$  = -3.46 + 4 = 0.54 bits.and the info. prov. by the fact that the Prof. is out of town on the fact that an exam will be given is  $I(A, B) = -\log \frac{11}{16} + \log \frac{8}{9}$  = 0.37 bits.

The information needed that an exam will be given is  $-\log P(A)$  so it is  $-\log 11 / 16$  because P(A) is 11 / 16 so this will be  $-\log 11 + \log 16$  and that will be -3.46 + 4 which is 0.54 bits and the information that professor by professor is out of town on the fact that an exam will be given is I(A, B) =  $-\log 11 / 16 + \log 8 / 9$  how is it coming here 8 is the event that exam and exam is given and B is the event that Professor is out of town so we are trying to find the mutual information of A and B.

So that is I(A, B) which is  $-\log P(A) + \log P(A/B)$  now  $-\log P(A)$  will become  $-\log 11 / 16$  now what about P(A/B) that we have already found that is 8 / 9 so it will be log 8/9 and that is why the resulting value will be 0.37 bits. (Refer Slide Time: 27:46)

Let C denotes the event that the Prof. is in town. Then  $P(A/C) = \frac{\frac{3}{16}}{\frac{3}{16} + \frac{1}{4}} = \frac{3}{7}$ We have  $I(A,C) = -\log \frac{11}{16} + \log \frac{3}{7}$  = -0.69 bits

Let C denotes the event that the professor is in town then P (A/C) will become P (A $\cap$ C) / P(C).

(Refer Slide Time: 28:22)



Now P(C) is the probability that the professor is in town and if we see the cases professor is in town in two cases  $x_3$  and  $x_4$  so we have to find the probability that professor is in town and for finding this we have to add P( $x_3$ ) and P( $x_4$ ) and that is what we are doing here.

(Refer Slide Time: 28:50)

Let C denotes the event that the Prof. is in town. Then  $P(A/C) = \frac{\frac{3}{16}}{\frac{3}{16} + \frac{1}{4}} = \frac{3}{7}$ We have  $I(A,C) = -\log \frac{11}{16} + \log \frac{3}{7}$  = -0.69 bits

So it will be 3 / 16 + 1 / 4 and in the numerator both the events will occur that probability so probability that A and C will occur so exam is given and professor is in town that probability if we consider the cases again.

(Refer Slide Time: 29:16)



So that is occurring here professor is in town and examination is given so  $x_3$  so that means we have to take the probability of  $x_3$ 

(Refer Slide Time: 29:31)

Let C denotes the event that the Prof. is in town. Then  $P(A/C) = \frac{\frac{3}{16}}{\frac{3}{16} + \frac{1}{4}} = \frac{3}{7}$ We have  $I(A,C) = -\log \frac{11}{16} + \log \frac{3}{7}$  = -0.69 bits

That will be 3 / 16 so if we simplify this we will get 3/7 so we will have I (A, C) that is the mutual information on taking this to events concerning this two events that is - log  $11 / 16 + \log 3 / 7$  and that is - 0.69 bits the fact that the professor is in town makes it less likely that an examination will be given so if we see so P (A/C) this is 3/7.

(Refer Slide Time: 30:37)

The information needed that an exam will be given is

$$-\log P(A) = -\log \frac{11}{16} = -\log 11 + \log 16$$
$$= -3.46 + 4 = 0.54$$
 hits

and the info. prov. by the fact that the Prof. is out of town on the fact that an exam will be given is

$$I(A, B) = -\log \frac{11}{16} + \log \frac{8}{9}$$
  
= 0.37 bits.

But P (A/B) it is 8 / 9 so from this we can conclude this that Professor is in town makes it less likely that an examination will be given consequently the mutual information provided by the presence of the professor on the occurrence of an examination is a negative quantity so that we have already observed that is -0.69. So the mutual information is very less now let us consider one example here.

(Refer Slide Time: 31:43)



Let us draw the figure first so this is the transmission into and this is the receiving end this is a simple model of a communication channel known as the binary symmetric channel at the transmission end either 0 or 1 is transmitted and at the receiving end either 0 or 1 is received specifically when 0 is transmitted 0 will be received with probability 1 -  $\varepsilon$  when 0 is transmitted 1 will be received with probability  $\varepsilon$  when 1 is transmitted 0 will be received with probability  $\varepsilon$  when 1 is transmitted 1 will be received with probability 1 -  $\varepsilon$ .

So this kind of communication channel we are considering suppose we have 2 equally likely messages  $m_1$  and  $m_2$  that will be transmitted over the channel using the representations 000 and 111 respectively if 010 was received we can compute the mutual information transmitted and the event that either part or the whole of the sequence 010 was received so if 010 was received we are trying to find the mutual information between the event that message  $m_1$  was transmitted and event that either part or the whole of the sequence 010 was received and that will be I ( $m_1$ , 0).

Now what is this will be - log the probability that  $m_1$  is transmitted + log the probability that 0 was received given that 0 was received the probability that  $m_1$  was transmitted so this value we can calculate in this way this - log P ( $m_1$ ) that will be - log  $\frac{1}{2}$  because  $m_1$  and  $m_2$  are equally likely so there are two possibilities  $m_1$  or  $m_2$  so that is why it will be  $\frac{1}{2}$  probability will be  $\frac{1}{2}$  so that is why this value will be - log  $\frac{1}{2}$ .

Now what about this conditional probability this we have to find so probability that  $m_1 / 0$  and this we know this probability that  $(m_1 \cap 0)$  both this events will occur  $m_1$  is transmitted as well

as 0 is received this probability and in the denominator probability that 0 is received and this can be written as probability that  $(m1 \cap 0) / P(m_1)$  into  $P(0 / m_1) + P(m_2)$  into  $P(0 / m_2)$  because this is this is P(0) this is nothing but the total probability so because there are 2 possibilities  $m_1$  and  $m_2$ so that is why either m1 or m2 will occur.

So P (m1) into P (0/m<sub>1</sub>) so this is the probability that  $m_1$  is transmitted that is given then what is the probability that 0 is disap and next is P (m<sub>2</sub>) into P (0, m<sub>2</sub>) now P (0, m<sub>2</sub>) is the probability that it is given that  $m_2$  is transmitted then what is the probability that 0 will be received.

(Refer Slide Time: 40:18)

So this numerator can be written as probability of  $m_1$  into  $P(0, m_1)/P(m_1)$  into  $P(0, m_1) + P(m_2)$  into  $P(0, m_2)$  so these probabilities we have to find. Now the numerator if you consider it will be  $P(m_1)$  into  $P(0, m_1) P(m_1)$  is  $\frac{1}{2}$  and  $P(0/m_1)$  that will be  $1 - \varepsilon$  so  $m_1$  given that  $m_1$  is transmitted 0 will be received with probability  $1 - \varepsilon$  okay so this will be  $\frac{1}{2}$  into  $1 - \varepsilon$  now in the denominator we will have  $P(m_1)$  so same whatever we are writing in the numerator this one is also same now the other one that will be  $\frac{1}{2}$  into  $\varepsilon$  so we will get  $I(m_1, 0) = -\log \frac{1}{2} + \log \frac{1}{2} 1 - \varepsilon / \frac{1}{2}$  into  $1 - \varepsilon + \frac{1}{2} \varepsilon$  so finally we are getting the value  $1 + \log 1 - \varepsilon$ .

(Refer Slide Time: 43:09)

$$\begin{split} \mathcal{I}\left(m_{1},0\right) &= -\log \mathcal{P}(m_{1}) + \log \mathcal{P}(m_{1}/o_{1}) \\ \mathcal{P}\left(m_{1}/o_{1}\right) &= -\frac{\mathcal{P}\left(m_{1}-\Omega_{1}\right)}{\mathcal{P}(m_{1})\mathcal{P}\left(\delta|_{2}\right)} \\ &= \frac{\mathcal{P}\left(m_{1},\Omega_{1}\right)}{\mathcal{P}(m_{1})\mathcal{P}\left(\delta|_{2}\right)} \\ &= \frac{\mathcal{P}\left(m_{1}\right)\mathcal{P}\left(\delta|_{2}\right)}{\mathcal{P}\left(m_{1}\right)\mathcal{P}\left(\delta|_{2}\right)\mathcal{P}\left(\delta|_{2}\right)\mathcal{P}\left(\delta|_{2}\right)} \\ &= \frac{\frac{1}{2}\left(1-\varepsilon\right)\varepsilon}{\frac{1}{2}\left(1-\varepsilon\right)\varepsilon} \\ &= \frac{1}{2} \cdot \\ \mathcal{I}\left(m_{1},0\right) &= -\log\frac{1}{2} + \log\frac{1}{2} = 0 \\ \mathcal{I}\left(m_{1},0\right) &= -\log\frac{1}{2} + \log\frac{1}{2} = 0 \\ \mathcal{I}\left(m_{1},0\right) &= -\log\frac{1}{2} + \log\frac{1}{2} \cdot \frac{1}{2}\left(1-\varepsilon\right)^{2}\varepsilon} \\ &= -\log\frac{1}{2} \cdot \\ \mathcal{I}\left(m_{1},0\right) &= -\log\frac{1}{2} + \log\frac{1}{2} \cdot \frac{1}{2}\left(1-\varepsilon\right)^{2}\varepsilon} \\ \mathcal{I}\left(m_{1},0\right) &= -\log\frac{1}{2} + \log\frac{1}{2} \cdot \frac{1}{2}\left(1-\varepsilon\right)^{2}\varepsilon} \\ &= -1 + \log\left(1-\varepsilon\right) \\ \mathcal{O} \quad \text{on old was necessed table us exactly the standard of the sta$$

Next we are considering I ( $m_1$ , 0 1) because we are trying to find the mutual information between the event that  $m_1$  is transmitted and the event that the whole sequence 010 or part of it is received so one part we have already considered that is 0 now we are considering another part that is 01 so I( $m_1$ , 01) this will be - log probability that  $m_1$  is transmitted + log probability that 0 1 is received then the probability that  $m_1$  is transmitted so this conditional probability log of this conditional probability.

So let us find  $m_1/01$  this will be probability  $(m_1 \cap 01) / P(m_1)$ into P(0/01) 1 sorry 0/01 given so let us write it freshly here  $m_1/01$  that will be probability that  $(m_1 \cap 01) / P(01)$  which is  $P(m_1)$ into  $P(0, m_1) / P(m_1)$ into  $P(01/m_1) + P(m_2)$ into  $P(01/m_2)$  now the numerator  $P(m_1)$  again  $\frac{1}{2}$  and this conditional probability will be  $1 - \varepsilon$  into  $\varepsilon$  so  $m_1$  given that  $m_1$  is transmitted that means 000 that is transmitted so from 0 will occur with  $P(1 - \varepsilon)$  and from 0, 1 will be received with  $P(\varepsilon)$ .

So that is why this we are getting in the denominator we will have  $\frac{1}{2}$  into  $1 - \varepsilon$  into  $\varepsilon +$  this also will give the same expression  $\frac{1}{2}$  into  $\varepsilon$  into  $1 - \varepsilon$  so this value will become  $\frac{1}{2}$  and as a result of this I (m<sub>1</sub>, 01) will be  $-\log \frac{1}{2} + \log \frac{1}{2}$  so this mutual information =0 in the same way if we find I (m<sub>1</sub>, 010) that will be  $-\log \frac{1}{2} + \log \frac{1}{2}$  into  $1 - \varepsilon$  whole square into  $\varepsilon / \frac{1}{2}$  into  $\varepsilon$  into  $1 - \varepsilon$  so this one will give  $1 + \log (1 - \varepsilon)$ .

So knowing that either 0 or 010 was received tells us exactly the same amount of information on the transmission on the transmission of message  $m_1$  however knowing that the sequence 01 was received.

(Refer Slide Time: 49:43)

$$= \frac{\rho(m_1)}{\rho(m_1)} \frac{\rho(o_1/m_1)}{\rho(m_1) + \rho(m_2)} \frac{\sigma(o_1/m_2)}{\sigma(m_1) + \rho(m_2)} = \frac{\frac{1}{2} (1-\varepsilon)\varepsilon}{\frac{1}{2}(1-\varepsilon)\varepsilon + \frac{1}{2}\varepsilon(1-\varepsilon)} = \frac{1}{2}$$

$$= \frac{1}{2}$$

$$I(m_1, 0) = -\log \frac{1}{2} + \log \frac{1}{2} = 0$$

$$I(m_1, 0) = -\log \frac{1}{2} + \log \frac{1}{2} \frac{1}{2}(1-\varepsilon)^{2}\varepsilon}{\frac{1}{2}\varepsilon(1-\varepsilon)^{2} + \frac{1}{2}\varepsilon^{2}(1-\varepsilon)} = \frac{1+\log(1-\varepsilon)}{\frac{1}{2}\varepsilon(1-\varepsilon)}$$

$$= 1+\log(1-\varepsilon)$$

$$O \text{ On OIO was necessed talls us exactly the same amount of information on the tomount of the same amount of information on the tomount of the tomount of medsage  $m_1$ .$$

Tells us nothing about the transmission off message  $m_1$  and this is why what we expect this is actually expected because the transmission of either  $m_1$  or  $m_2$  would yield the sequence 01 at the receiving end with the same probability, so now let us discuss one result of mutual information.

(Refer Slide Time: 51:08)

$$I(A,B) = -\log P(A) + \log P(A/B)$$
  
=  $-\log P(A) - \log P(B) + \log P(A \cap B)$   
=  $-\log P(B) - \log P(A) + \log P(A \cap B)$   
=  $-\log P(B) + \log \frac{P(A \cap B)}{P(A)}$   
=  $-\log P(B) + \log P(B/A)$   
=  $I(B,A)$ 

I(A, B) =- log P(A) + log P(A/B) and that =- log of P(A) - log P(B) + log P(A  $\cap$  B) because P(A/B) is P(A  $\cap$  B) / P(B) so we can write this as - log P(B) - log P(A) + log P(A  $\cap$  B) so it will be - log P(B) + log P(A  $\cap$  B) / P(A) and that =- log P(B) + log P(B / A) and this is nothing but I(B, A).

(Refer Slide Time: 52:28)



So this implies mutual information is a symmetric measure in the information concerning 2 events thus I(A, B) is a measure of the mutual information from B to A as well as from A to B that is all thank you.

**Educational Technology Cell** Indian Institute of Technology Roorkee

**Production for NPTEL** Ministry of Human Resource Development Government of India

For Further Details Contact

Coordinate, Educational Technology Cell Indian Institute of Technology Roorkee Roorkee-247667 E Mail: <u>etcell@iitr.ernet.in</u>, <u>etcell.iitrke@gmail.com</u> Website: <u>www.nptel.iim.ac.in</u>

> Acknowledgement Prof. Pradipta Banerji Director, IIT Roorkee

Subject Expert & Script Dr. Aditi Gangopadhyay Dept of Mathematics

IIT Roorkee

#### **Production Team**

Neetesh Kumar Jitender Kumar Pankaj Saini Meenakshi Chauhan

# Camera

Sarath Koovery Younus Salim

**Online Editing** Jithin.k

010111111

**Graphics** Binoy.V.P

**NPTEL Coordinator** Prof.Bikash Mohanty

An Educational Technology Cell IIT Roorkee Production @ Copyright All Rights Reserved WANT TO SEE MORE LIKE THIS SUBSCRIBE