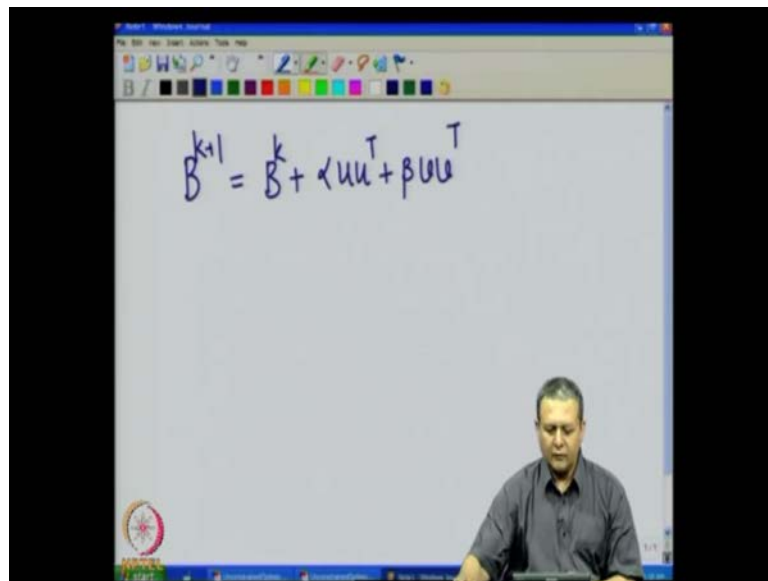


Numerical Optimization
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Lecture - 17
Quasi-Newton Methods - Broyden Family
Coordinate Descent Method

Welcome back to this series of lectures on numerical optimization. In the last class we discussed about symmetric rank one update, and rank two update. So, in particular we saw that there is no guaranty in symmetric rank one update, that the new matrix B_{k+1} is positive definite even though B_k is positive definite. So, we decided to use the symmetric rank two update.

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So, the matrix B_{k+1} was found using $B_k + \alpha uu^T + \beta vv^T$, transpose. Now, the matrix B_{k+1} should also satisfy Quasi-Newton condition, that is $B_{k+1} \gamma_k = \delta_k$. So, if you use that condition then we saw that we got the Davidon Fletcher Powell method.

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$$B_{DFP}^{k+1} = B^k + \frac{\delta^k \delta^{kT}}{\delta^{kT} \gamma^k} - \frac{B^k \gamma^k \gamma^{kT} B^k}{\gamma^{kT} B^k \gamma^k} \quad (\text{DFP Method})$$

- Is B_{DFP}^{k+1} a symmetric positive definite matrix, given that B^k is symmetric positive definite matrix?

B_{DFP}^{k+1} is a symmetric matrix.
 Let $x \neq 0, \gamma^k \neq 0, \delta^k \neq 0$.

$$x^T B_{DFP}^{k+1} x = x^T B^k x - \frac{(x^T B^k \gamma^k)^2}{\gamma^{kT} B^k \gamma^k} + \frac{(\delta^{kT} x)^2}{\delta^{kT} \gamma^k}$$

Since B^k is symmetric, $B^k = B^{k\frac{1}{2}} B^{k\frac{1}{2}}$ where $B^{k\frac{1}{2}}$ is symmetric and positive definite. Letting $a = B^{k\frac{1}{2}} x$ and $b = B^{k\frac{1}{2}} \gamma^k$,

$$x^T B_{DFP}^{k+1} x = \frac{(a^T a)(b^T b) - (a^T b)^2}{b^T b} + \frac{(\delta^{kT} x)^2}{\delta^{kT} \gamma^k}$$

And we got this formula for $B^k + 1$ which is nothing but $B^k + \delta^k \delta^{kT}$ by $\delta^k \delta^{kT} \gamma^k - B^k \gamma^k \gamma^k B^k$ by $\gamma^k \gamma^k B^k \gamma^k$, and then we started looking at the positive definiteness of this matrix. So, if we showed that this matrix is positive semi-definite. So that is clear from this, along with the fact that the numerator here is non-negative using in because of the quasi-strict inequality, this quantity is also non-negative because of the quasi-strict inequality. So, we saw that $x^T B_{DFP}^{k+1} x$ is greater than or equal to 0, and then we moved on to show that it is indeed the positive definite matrix.

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$$x^T B_{DFP}^{k+1} x = \frac{(a^T a)(b^T b) - (a^T b)^2}{b^T b} + \frac{(\delta^{kT} x)^2}{\delta^{kT} \gamma^k}$$

We now show that B_{DFP}^{k+1} is positive definite, that is,

$$x^T B_{DFP}^{k+1} x > 0, x \neq 0$$

We have already shown that $\delta^{kT} \gamma^k > 0$.
 Suppose $x^T B_{DFP}^{k+1} x = 0, x \neq 0$.
 Therefore, $(a^T a)(b^T b) = (a^T b)^2$ and $(\delta^{kT} x)^2 = 0$.

$$(a^T a)(b^T b) = (a^T b)^2 \Rightarrow a = \mu b \Rightarrow x = \mu \gamma^k \Rightarrow \mu \neq 0$$

$$(\delta^{kT} x)^2 = 0 \Rightarrow \mu \delta^{kT} \gamma^k = 0 \Rightarrow \delta^{kT} \gamma^k = 0 \quad (\text{contradiction})$$

Therefore, $x^T B_{DFP}^{k+1} x > 0, x \neq 0 \Rightarrow B_{DFP}^{k+1}$ is positive definite.

And for that purpose what we did was, we took $x^T B^k x$ which is non-zero and I will if you assume that $x^T B^k x + 1$ is 0, so which means that both this quantities are non-negative. So, if their sum is 0 that means each of the quantities has to be 0. So, which means $x^T B^k x$ is 0 and 1 is 0, which is a contradiction. So, this essentially means that x is a scalar multiple of b and which means that if you recall the definition of x and γ , we can write this x is a multiple of $\gamma^T B^k \gamma$ and since x is not 0. We can say that μ is not 0. We have assumed that x is not 0 and similarly, if you look at this now this quantity $\gamma^T B^k \gamma = 0$ implies that let us substitute the value of x here. So, μ into $\gamma^T B^k \gamma = 0$.

And since μ is not 0, this will be true only when $\gamma^T B^k \gamma = 0$ but we have already shown that $\gamma^T B^k \gamma > 0$ and that contradicts our earlier assumption that $\gamma^T B^k \gamma > 0$ and therefore, we have $x^T B^k x > 0$ for all x not equal to 0, which implies that B^k is positive definite. So, if B^k is symmetric and positive definite, we say that $B^k + 1$ is also symmetric and positive definite if we use the symmetric rank 2 correction or if we use the DFB method. So, let us see how the algorithm looks like if we use DFB method in our Quasi-Newton algorithm.

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Quasi-Newton Algorithm (DFP Method)

- (1) Initialize x^0 , ϵ and symmetric positive definite B^0 , set $k := 0$.
- (2) **while** $\|g^k\| > \epsilon$
 - (a) $d^k = -B^k g^k$
 - (b) Find $\alpha^k (> 0)$ along d^k such that
 - (i) $f(x^k + \alpha^k d^k) < f(x^k)$
 - (ii) α^k satisfies Armijo-Wolfe (or Armijo-Goldstein) conditions
 - (c) $x^{k+1} = x^k + \alpha^k d^k$
 - (d) Find B^{k+1} using DFP method
 - (e) $k := k + 1$

endwhile

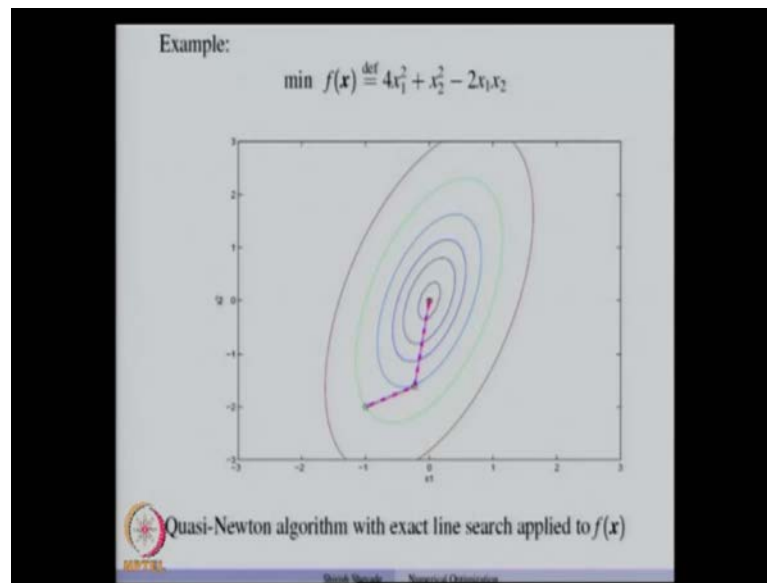
Output : $x^* = x^k$, a stationary point of $f(x)$.

So, compare to the other approaches that we have seen. So, for what we need here is a symmetric positive definite matrix B_0 to start with of course, we need the initial point x_0 , the stopping tolerance and then we said the iteration counter to 0. Now while the norm of the gradient is greater than epsilon, we said the direction d_k to be minus $B_k^{-1} g_k$. Now initially we have some matrix B_0 . So, suppose if we have no knowledge about the problem then we could use the B_0 matrix to be an identity matrix then we do the line search. So, that Armijo-Goldstein or Armijo-Wolf conditions are satisfied. Then we get x_{k+1} using this formula, that the formula which we have used in our earlier algorithms also and now is the time to find out B_{k+1} using b_k , x_k , x_{k+1} and g_k and g_{k+1} . So, if you have all these information we can use DFB method to update B_k to B_{k+1} .

So, we get B_{k+1} using suppose DFB method, and then increment the iteration counter and at the new point if the gradient norm of the gradient is less than or equal to epsilon we quit or stop the algorithm or otherwise continue this algorithm till this stopping condition is satisfied and at the end what we get is x^* to be x_k a stationary point of $f(x)$.

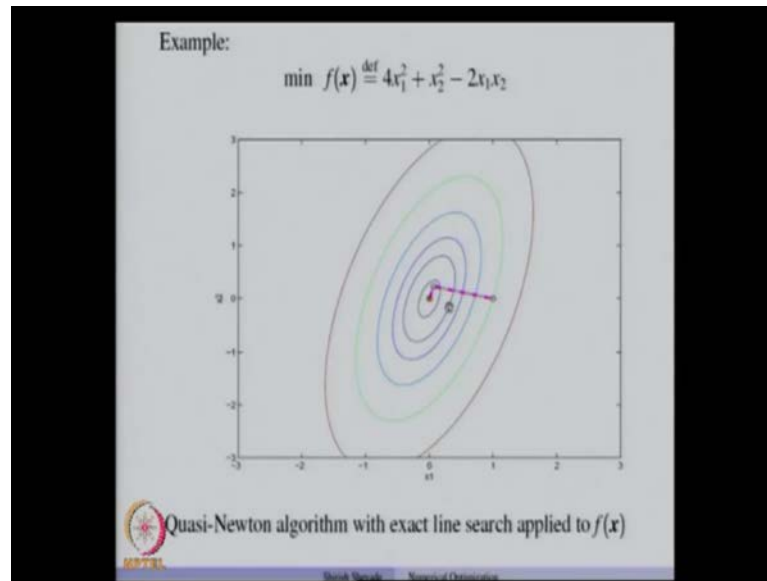
So, you will see that the main difference of this algorithm with Newton algorithm or modified Newton algorithm is that, there is no need to find out inverse of any matrix in fact that direction d_k is found by matrix where vector multiplication as compare to the direction that that we use in the Newton method the direction that was used in the Newton method was d_k to be minus $h_k^{-1} g_k$. So, that required inversion of a matrix and plus the multiplication of matrix and a vector while here there is only multiplication of a matrix by a vector. So, that saves lot of computational time per iteration in DFB method and this update is symmetric rank two update and it is easy to do this update. So, in all what we did need is that matrix vector multiplications or vector multiplications to in one iteration to get the update B_{k+1} .

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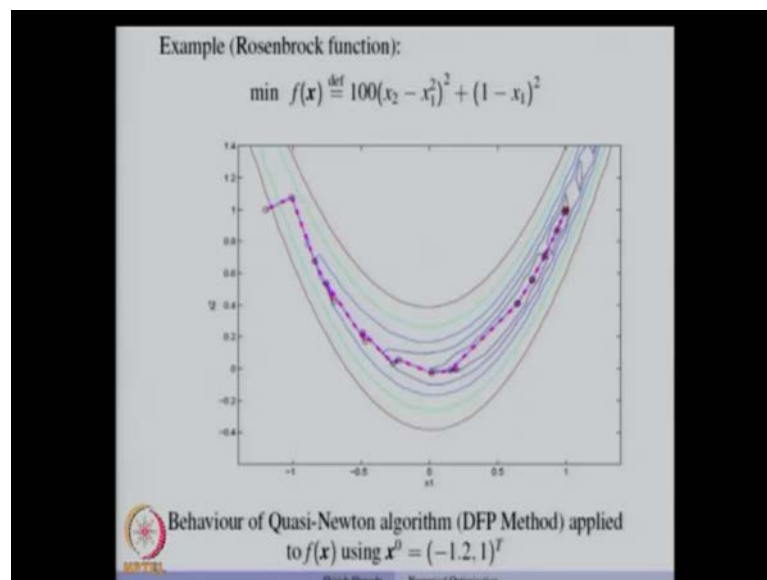
Now, let us take example consider quadratic function which is given here $f(x)$ equal to $4x_1^2 + x_2^2 - 2x_1x_2$. We have already seen when we discussed about other methods. So, these are the contours this is going to be the minimum of this f and when we apply Quasi-Newton algorithm with exact line search remember that, this is the quadratic function. So, is easy to work out these formulas for exact line search. So, if we use Quasi-Newton method the DFB method in the Quasi-Newton algorithm and exact line search. So, if we start from this point. So, this is our x_0 the x_1 is here and x_2 is here and x_2 is nothing but x^* . So, we require two steps to reach the solution of the given problem.

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Now we start from another point. So, if you start from here. So, here we will see that will require again two steps this is x_0 , x_1 and x_2 is nothing but x^* . So, we reach the solution in both the cases in exactly two steps.

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Now, let us look at the performance of Quasi-Newton algorithm with DFB method on the Rosenbrock function. So, if you start from this point you will see that initially it make some movement to this and then there are small steps along this and then finally it goes and converges to the local minimum.

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Example (Rosenbrock function):

$$\min f(x) \stackrel{\text{def}}{=} 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

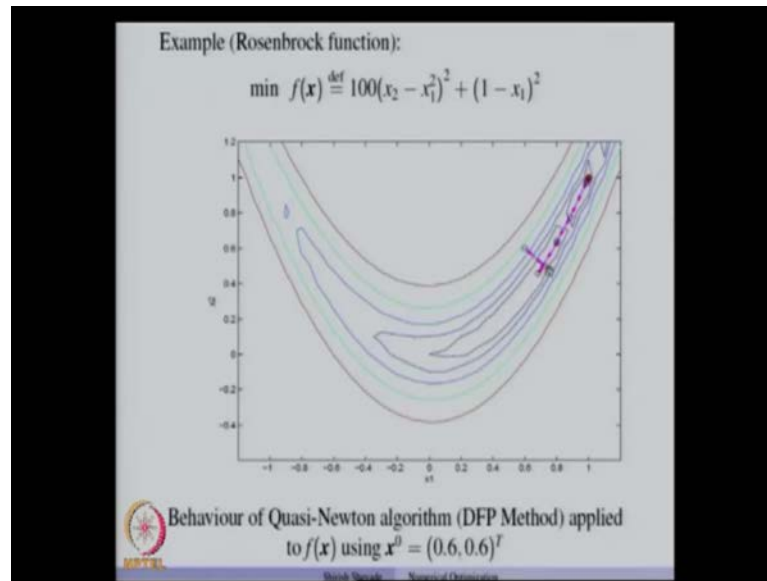
k	x_1^k	x_2^k	$f(x^k)$	$\ g^k\ $	$\ x^k - x^*$
0	-1.2	1	24.2	232.86	2.2
1	-1.01	1.08	4.43	24.97	2.01
2	-0.84	0.68	3.47	14.53	1.87
3	-0.70	0.42	3.33	25.61	1.79
4	-0.76	0.54	3.18	14.19	1.81
5	-0.47	0.17	2.37	14.80	1.69
10	0.20	-0.01	0.84	9.00	1.29
15	0.75	0.56	0.06	0.34	0.50
20	0.99	0.99	0.0002	0.69	0.02
24	0.99	0.99	5.72×10^{-12}	2.25×10^{-6}	5.35×10^{-6}

Table: Quasi-Newton algorithm (DFP Method) applied to Rosenbrock function, using $x^0 = (-1.2, 1.0)^T$.

Let us look at the iterations with that initial point. So, initial point was a minus 1.2 and 1. The value of the function at that point was 24.2 and the norm of the gradient was 232.86 and x^* is 1 1. So, this turns of x^k from x^* is 2.2. Note that, this distance is a Euclidean distance now from iteration. So, this was a initial point. So, when we go to 0.x1 which is minus 1.01, 1.02. You would see that there is a significant decrease in the objective function value from 24.22. It has come down to 4.43 and again there is a significant decrease in the norm of g^k allow the distance between the x^k and x^* has not decrease that much but then as the iteration progress you will see that and the end of 10th iteration the function value is 0.84.

The norm of the gradient is 9 and the distance has come down to 1.29 and at the end of 20 iterations you would see that the value of the function is very small. The norm is also very small and the distance of x^k from x^* is also small and if you use the stopping criteria of norm g^k less than or equal to epsilon where epsilon is 10 to the power minus 3, then we will see that after 24 iterations that norm g^k is indeed less than epsilon and the algorithm terminates.

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Now, if we start from some other point which is say 0.6, 0.6 and 0.6. So, this is how the algorithm behaves.

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Example (Rosenbrock function):

$$\min f(x) \stackrel{\text{def}}{=} 100(x_2 - x_1^2)^2 + (1 - x_1)^2$$

k	x_1^k	x_2^k	$f(x^k)$	$\ g^k\ $	$\ x^k - x^{k-1}\ $
0	0.6	0.6	5.92	75.60	0.57
1	0.72	0.50	0.12	6.32	0.572095
2	0.68	0.46	0.11	1.56	0.629112
3	0.80	0.63	0.06	4.65	0.421985
4	0.80	0.64	0.04	0.39	0.410591
5	0.88	0.78	0.02	2.67	0.252278
6	0.99	0.98	0.0005	0.94	0.0238278
7	0.98	0.97	0.0003	0.33	0.0348487
8	0.99	0.99	7.8×10^{-5}	0.23	0.02
9	0.99	0.99	5.3×10^{-7}	0.0044	0.0016
10	0.99	0.99	1.1×10^{-8}	0.0048	3.2×10^{-5}
11	0.99	0.99	3.1×10^{-13}	3.2×10^{-6}	1.2×10^{-6}

Table: Quasi-Newton algorithm (DFP Method) applied to Rosenbrock function, using $x^0 = (0.6, 0.6)^T$.

And if you look at the table so you will see that we started from 0.6, 0.6 the value of function of 5.92 norm of G^k is 75.6 and we are reasonably close to the solution compare to the previous initial point that we saw. Now, in the first iteration itself the value of the function came down to 0.12 there was a slight increase in the distance of x^k from x^* but norm of G^k has also come down significantly and as a iterations progress you will

see that in the 7th iteration itself the value of the function is 0.0003 and norm of the gradient is 0.33. So, value of the function is quite small the distance of x_k from x^* is quite small. So, you will see that we are almost close to the solution here at the end of the 7th iteration and although the x_1 and x_2 co-ordinates for 8, 9th, and 10th iterations would look like they are different but because of the lack of space they could not be shown exactly. So, the approximate numbers are given here but you would see that the values of the functions are different. So, although they look similar here actually the points are different and that is clear from the value of the function that you get here. So, you will see that at the end of 10th iteration norm of g_k is 0.0048 and the distance of x_k from x^* is 3.2×10^{-5} and the next iteration norm g_k becomes still less than 0.001 and the algorithm terminates. So, you would see that the value of the function is close to 0 in fact the value of the function is 0. When at the local minimum and that local minimum is x_1 equal to 1 and x_2 equal to 1 and the algorithm has found that particular point with certain accuracy.

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- Newton direction, $d_N^k = -(H^k)^{-1} g^k$
- Quasi-Newton direction, $d_{QN}^k = -B^k g^k$
 - Quasi-Newton condition: $B^{k+1} \gamma^k = \delta^k$
 $\Rightarrow \gamma^k = (B^{k+1})^{-1} \delta^k$
- Let $G^{k+1} = (B^{k+1})^{-1}$ approximate H^{k+1} .
Therefore, we get dual formulae:
$$B^{k+1} \gamma^k = \delta^k$$

$$G^{k+1} \delta^k = \gamma^k$$

$$B_{DFP}^{k+1} = B^k + \frac{\delta^k \delta^{kT}}{\delta^{kT} \gamma^k} - \frac{B^k \gamma^k \gamma^{kT} B^k}{\gamma^{kT} B^k \gamma^k} \quad (\text{DFP Method})$$

$$G_{BFGS}^{k+1} = G^k + \frac{\gamma^k \gamma^{kT}}{\gamma^{kT} \delta^k} - \frac{G^k \delta^k \delta^{kT} G^k}{\delta^{kT} G^k \delta^k} \quad (\text{BFGS Method})$$

Now, recall that the Newton direction that we used was minus H_k inverse g_k where H_k is the Hessian matrix at the iteration at the current iteration k and g_k is the gradient at the current iteration k . Now, when we use Quasi-Newton method, we use Quasi-Newton direction and the idea the reset approximate the H_k inverse matrix by a positive definite matrix B_k and use the direction minus $B_k g_k$. So, this helps us avoid the inverse calculations as compared to the Newton direction. Now, the Quasi-Newton direction d_k ,

which uses the matrix B_k we know that for the Quasi-Newton direction to work the matrix B_k should satisfy Quasi-Newton condition and that is $B_{k+1} \gamma_k$ is equal to δ_k .

Now, if B_{k+1} is positive definite. So, it is clearly invertible and therefore, we can set γ_k to be $B_{k+1}^{-1} \delta_k$. Now let us look at some matrix G_{k+1} , which is obtained using B_{k+1}^{-1} and let us assume that, this G_{k+1} approximates H_{k+1} . So, if we take the inverse of the matrix B_k , we get g_k to be B_k^{-1} and that g_k will be a good approximation of H_k .

So, if you combine this formula $B_{k+1} \gamma_k = \delta_k$ and write a formula in terms of $\gamma_k = B_{k+1}^{-1} \delta_k$ and B_{k+1}^{-1} is replaced by g_{k+1} then what we get is $g_{k+1} \delta_k = \gamma_k$. So, if you look at this two formulae, there is a dual relationship that exist between the variables for example, there is a relationship between δ_k and γ_k here γ_k and δ_k here and B_{k+1} , g_{k+1} here. So, that essentially tells us that if we have some update formula for the matrix B using γ_k and δ_k .

We can as well get an update formula for G by replacing γ in the B update formula by δ and δ in the B update formula by γ and the relationship is dual. So, on the other hand if we have a update formula for G , we can replace G by B in that formula, δ by γ in that formula and γ by δ and we get a update formula for B . So, this dual relationship can be exploited to get different formulas and note that here we in this formula we work with the matrix G which is an approximation of the matrix H at given iteration. So, if you look at DFB method.

Where we had this update formula $B_{k+1} = B_k + \text{this quantity} - \text{the other quantity}$ which is $B_k \gamma_k \gamma_k^T - b_k b_k^T$ by $\gamma_k \gamma_k^T - b_k b_k^T$ now, if you use this dual relationship we can write a formula in terms of G δ in γ by replacing B by G and γ by δ and δ by γ in this formula and what we get is called the BFGs update for G .

So, the BFGs stands for Broyden Fletcher Goldfarb and Shannon the inventors of this method update method. So, you will see that B is replaced by G δ is replaced by γ and γ is replaced by δ and we get update formula for G and we expect that.

This G is a good approximation of the actual hessian. Now remember that if we use the matrix G in this formula, then what we have to write is that $d_k^T Q d_k$ is equal to $-g_k^T G^{-1} g_k$, where the first G^{-1} corresponds to on the matrix G which approximates H_k . So, which again involves inversion of a matrix and that is going to be computational expensive. So, instead we would like to get a formula for G and then find a corresponding update formula for B and that formula can be use here instead of G^{-1} and with let us see how to do that.

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$$G_{BFGS}^{k+1} = G^k + \frac{\gamma^k \gamma^{kT}}{\gamma^{kT} \delta^k} - \frac{G^k \delta^k \delta^{kT} G^k}{\delta^{kT} G^k \delta^k} \quad (\text{BFGS Method})$$

- How to get B_{BFGS}^{k+1} from G_{BFGS}^{k+1} ?
- Use the condition,

$$B_{BFGS}^{k+1} G_{BFGS}^{k+1} = I$$

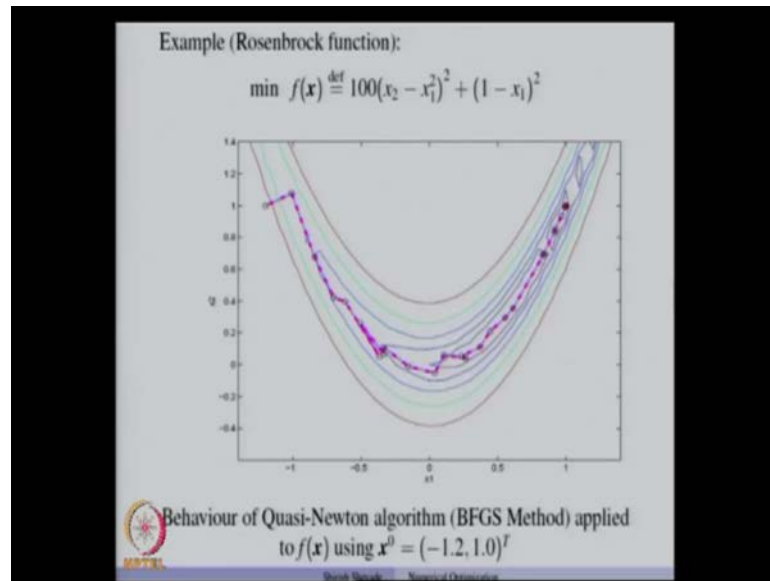
to get

$$B_{BFGS}^{k+1} = B + \left(1 + \frac{\gamma^T B \gamma}{\delta^T \gamma}\right) \frac{\delta \delta^T}{\delta^T \gamma} - \left(\frac{\delta \gamma^T B + B \gamma \delta^T}{\delta^T \gamma}\right)$$

So, the natural question that one would like to ask is that how to get B_{k+1} from G_{k+1} ? And that is obvious from this relationship that G_{k+1} is equal to B_{k+1} inverse. So, which means that B_{k+1} into G_{k+1} should be equal to an identity matrix. So, if we use that relationship that B_{k+1} BFGS into G_{k+1} BFGS equal to identity then using some matrix identities. What we get is the BFGS update formula for the matrix B and that would be something like this.

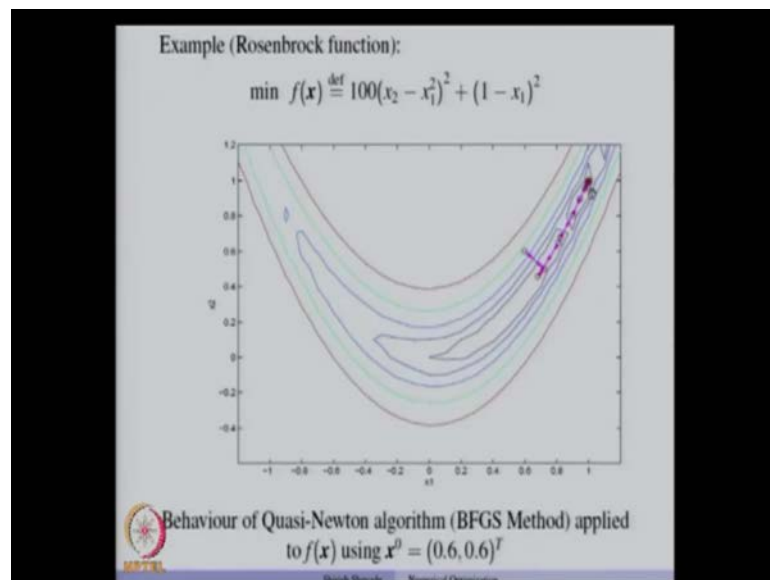
So, the right side to avoid notational clutter all the superscripts k has been suppressed but they do exist. So, all the right side terms involving B , γ and δ are the terms evaluated at the current iteration k . So, B is B_k , γ is γ_k and δ is δ_k in the right hand side of this expression. So, this is called BFGS update formula for the matrix B . So, one can derive similar formulas given any update formula for the matrix B .

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Now, let us apply this B F G S method to Rosenbrock function and this is the behavior that one gets of the Quasi-Newton algorithm where B F G S method is used.

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And so if you start from a different initial point we see a similar behavior as we saw in the DFB method.

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Broyden Family

- DFP method

$$B_{DFP}^{k+1} = B^k + \frac{\delta^k \delta^{kT}}{\delta^{kT} \gamma^k} - \frac{B^k \gamma^k \gamma^{kT} B^k}{\gamma^{kT} B^k \gamma^k}$$

- BGFS method

$$B_{BGFS}^{k+1} = B + \left(1 + \frac{\gamma^T B \gamma}{\delta^T \gamma}\right) \frac{\delta \delta^T}{\delta^T \gamma} - \left(\frac{\delta \gamma^T B + B \gamma \delta^T}{\delta^T \gamma}\right)$$

- Broyden Family

$$B^{k+1}(\varphi) = \varphi B_{BGFS}^{k+1} + (1 - \varphi) B_{DFP}^{k+1}$$

where $\varphi \in [0, 1]$.

So, so far we have studied two rank one update formulas one is DFB, two rank two update formulas one is the DFB formula and other one is the B F G S formula. Now these two formulae can be combined to get an update formula which belongs to what is called Broyden family and that would look something like this that... So, if we chose phi to be a positive fraction and B k plus 1 in the Broyden family is obtained using phi in to B k plus 1 B F G S plus 1 minus phi into B k plus 1 DFB. So, you will see that when phi equal to 0.

What we have is the DFB update formula and when phi equal to 1 we have the B F G S update formula. So, Broyden family consists of BFP formula as well as B F G S formula at the two extremes plus convex combination of the B F G S and DFB updates for the matrix B. So, this is also called pure Broyden family, because phi is kept constant here sometimes this phi can be also made a function of the iteration k. Now it is clear that if we chose phi in this range this B k plus 1 will be symmetric and positive definite, because we know that both B F G. And DFB update formulas for B are symmetric and positive definite. So, this gives a family of update rules for the matrix B and this is called Broyden family.

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Quasi-Newton Algorithm (Broyden Family)

- (1) Initialize x^0 , ϵ and symmetric positive definite B^0 ,
 $\varphi \in [0, 1]$, set $k := 0$.
- (2) **while** $\|g^k\| > \epsilon$
 - (a) $d^k = -B^k(\varphi)g^k$
 - (b) Find $\alpha^k (> 0)$ along d^k such that
 - (i) $f(x^k + \alpha^k d^k) < f(x^k)$
 - (ii) α^k satisfies Armijo-Wolfe (or Armijo-Goldstein) conditions
 - (c) $x^{k+1} = x^k + \alpha^k d^k$
 - (d) $B^{k+1}(\varphi) = \varphi B_{DFGS}^{k+1} + (1 - \varphi) B_{DFB}^{k+1}$
 - (e) $k := k + 1$

endwhile

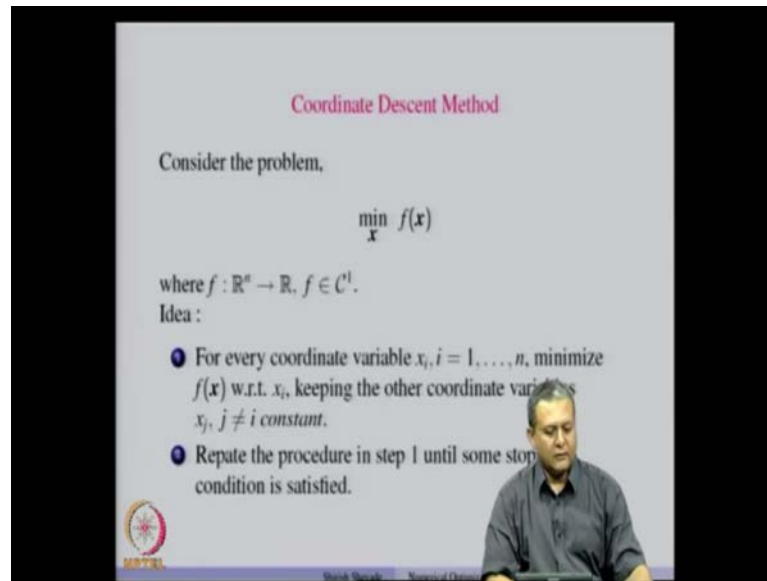
Output : $x^* = x^k$, a stationary point of $f(x)$.

So, the algorithm when we use Broyden family is given here. So, what we need is that apart from the symmetric positive definite matrix B naught we need the variable φ which is in the closing trouble 0 to 1. So, the first step is to get direction d^k which is minus B^k of φ into G^k and best on the value of φ will get B minus B^k φ that multiply by G^k would give us B^k the line search is then done.

So, has to satisfy Armijo Wolf or Armijo Goldstein Conditions, then we get x^{k+1} to be x^k plus $\alpha^k d^k$. The new point is obtained and then we have to update B^k and that is obtain using Broyden family update formula and that is B^{k+1} is a convex combination of B_{DFGS}^{k+1} and B_{DFB}^{k+1} . DFB and the iteration contours are increased as in the previous cases and finally, we get stationary point when the algorithm terminates. So, these are some of the update formulas in the Quasi-Newton method and will come back to this Quasi-Newton directions and Quasi-Newton methods some time later.

So, let us now look at some of the simplest techniques for minimizing a function and one of the simplest techniques is called Coordinate Descent Method.

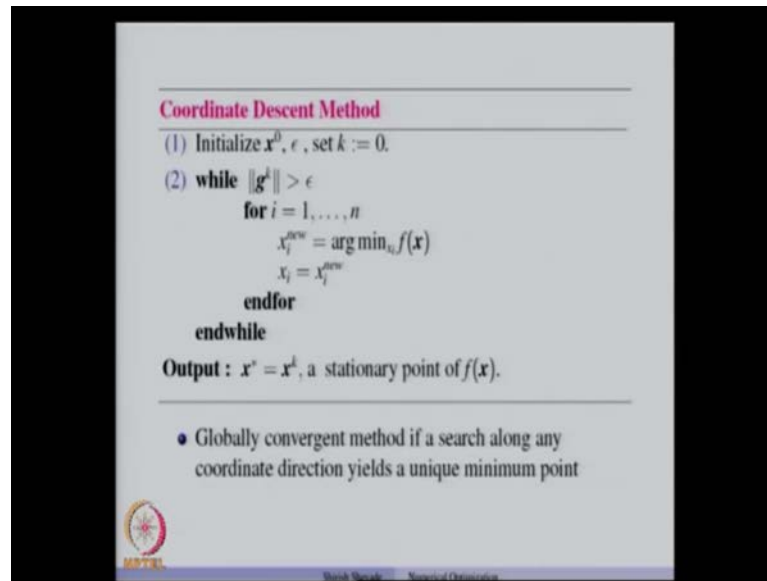
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The slide is titled "Coordinate Descent Method" in red text. Below the title, it says "Consider the problem," followed by the mathematical expression $\min_x f(x)$. Underneath, it states "where $f: \mathbb{R}^n \rightarrow \mathbb{R}$, $f \in \mathcal{C}^1$." and "Idea:". There are two bullet points: the first says "For every coordinate variable $x_i, i = 1, \dots, n$, minimize $f(x)$ w.r.t. x_i , keeping the other coordinate variables $x_j, j \neq i$ constant." and the second says "Repeat the procedure in step 1 until some stopping condition is satisfied." In the bottom right corner of the slide, a man in a dark shirt is visible, likely the lecturer. There is also a small logo in the bottom left corner of the slide.

So, as a name suggest the directions that are chosen are the coordinate directions and. So, we choose one coordinate direction at a time and optimize a function with respect to that coordinate alone keeping the other coordinates fixed and the procedure is repeated till we get an optimum point. So, let us see how to do that. So, consider this problem where we want to minimize $f(x)$, f is a real valued function whose domain is \mathbb{R}^n and f is differentiable. Now the idea of coordinate descent method is that, we consider every coordinate x_i there are n such coordinates and minimize the function $f(x)$ with respect to x_i alone keeping the other coordinate variables x_j where j is not equal to i constant. So, this becomes one-dimensional optimization problem which is many times easy to solve compare to n -dimensional optimization method. So, this procedure is repeated for every coordinate variable. So, after scanning all the n coordinates, once again goes back to the first coordinate and starts optimizing it and the whole procedure is repeated till some stopping criteria is satisfied.

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Coordinate Descent Method

(1) Initialize x^0, ϵ , set $k := 0$.

(2) **while** $\|g^k\| > \epsilon$
 for $i = 1, \dots, n$
 $x_i^{new} = \arg \min_x f(x)$
 $x_i = x_i^{new}$
 endfor
endwhile

Output : $x^* = x^k$, a stationary point of $f(x)$.

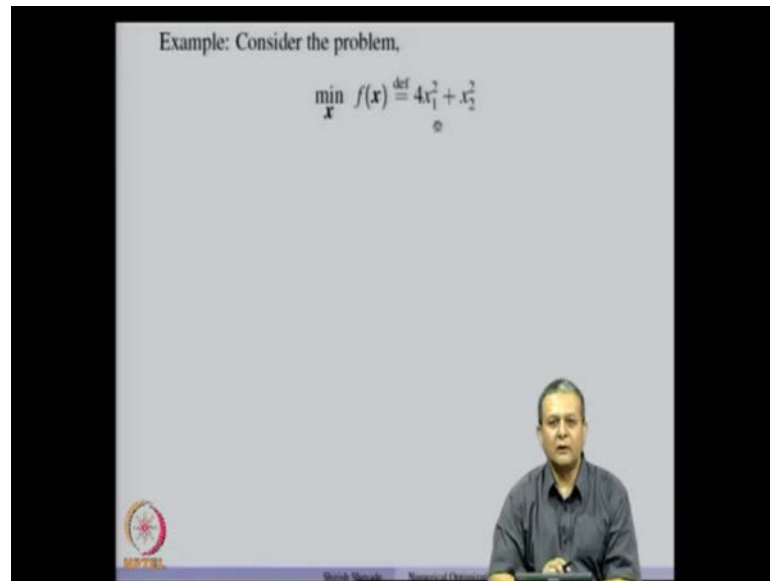
- Globally convergent method if a search along any coordinate direction yields a unique minimum point

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So, this procedure in step one is repeated till the stopping condition is satisfied. So, let us look at algorithm is very simple, what we need is a initial point some stopping tolerance which said the iteration contour to 0. So, while the norm of the gradient is greater than epsilon, we take each coordinate and minimize the function $f(x)$ with respect to that coordinate x_i and when we find the minimum we get the new value of x_i and that will be x_i^{new} then the variable x is set to x_i^{new} and we are ready to go to the next coordinate so this procedure.

So, once scan of all the coordinates is complete then we go back and check whether the norm of the gradient is less than or equal to epsilon, if it is not we continue again scanning all the n coordinates and bring the updates till norm of g^k becomes less than or equal to epsilon and that point we terminate the algorithm. So, a very simple way of optimizing a function and under certain assumptions one can show that, this method is globally convergent. That is if we assume that a search along any coordinate direction you will say unique minimum point then this method is globally convergent. So, a very simple method and under some mild assumptions of existence of unique minimum along coordinate direction right makes this method globally convergent. Now, let us take some example.

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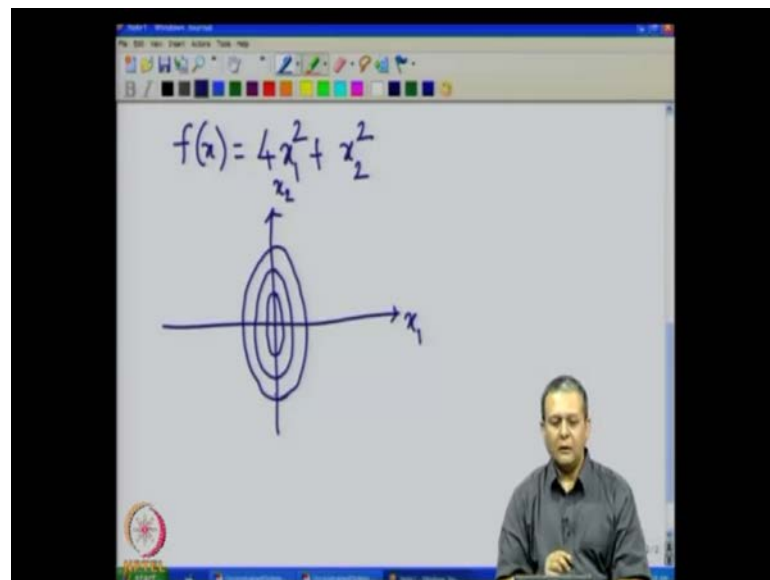
Example: Consider the problem,

$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

The slide features a man in a dark shirt at the bottom right, a logo in the bottom left, and a taskbar at the bottom.

So, let us consider the problem there we want to minimize $f(x)$ is the quadratic function in this case and $f(x)$ is the defined as $4x_1^2 + x_2^2$. So, let us look at this function first.

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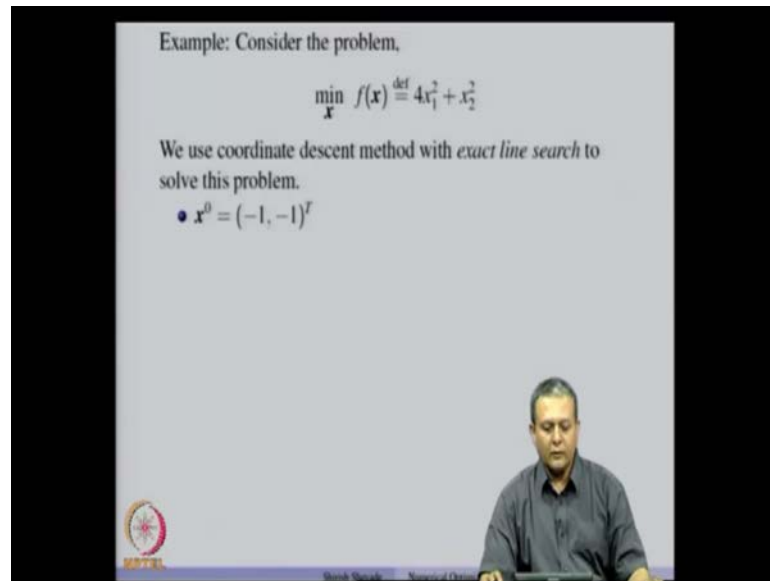


$f(\mathbf{x}) = 4x_1^2 + x_2^2$

The slide shows a hand-drawn graph of the function's contours in the x_1 - x_2 plane. The contours are concentric ellipses centered at the origin, elongated along the x_2 axis. The horizontal axis is labeled x_1 and the vertical axis is labeled x_2 . The slide also features a man in a dark shirt at the bottom right, a logo in the bottom left, and a software toolbar at the top.

So, we have $f(x)$ equal to $4x_1^2 + x_2^2$ now let us draw the contours of this function. So, this is x_1 and this is x_2 and a contours would look something like this. So, this is the function that we want to optimize.

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Example: Consider the problem,

$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

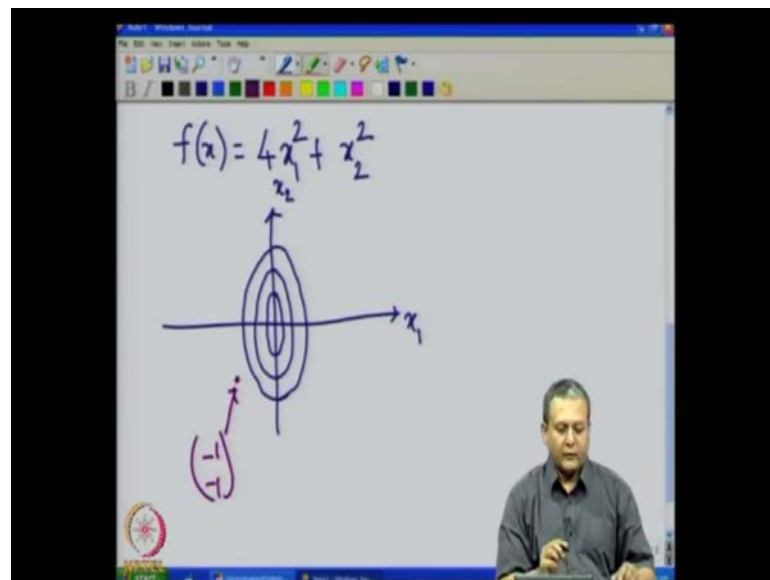
We use coordinate descent method with *exact line search* to solve this problem.

- $\mathbf{x}^0 = (-1, -1)^T$

The slide features a speaker in the bottom right corner and a logo in the bottom left corner.

And suppose we decide to use coordinate descent method and seen the function is quadratic it is easy to use the exact line search to solve this problem so let the initial point be minus 1 minus 1. So, let us consider some initial point.

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$f(\mathbf{x}) = 4x_1^2 + x_2^2$

The slide shows a graph of the function $f(\mathbf{x}) = 4x_1^2 + x_2^2$ in the x_1 - x_2 plane. The graph consists of several concentric ellipses centered at the origin, elongated along the x_2 axis. A point $(-1, -1)^T$ is marked in the third quadrant with a pink arrow pointing to it. The slide also features a speaker in the bottom right corner and a logo in the bottom left corner.

So, let me denote it suppose this point is minus 1 minus 1. So, this is the point from which we start and let us see how the algorithm works.

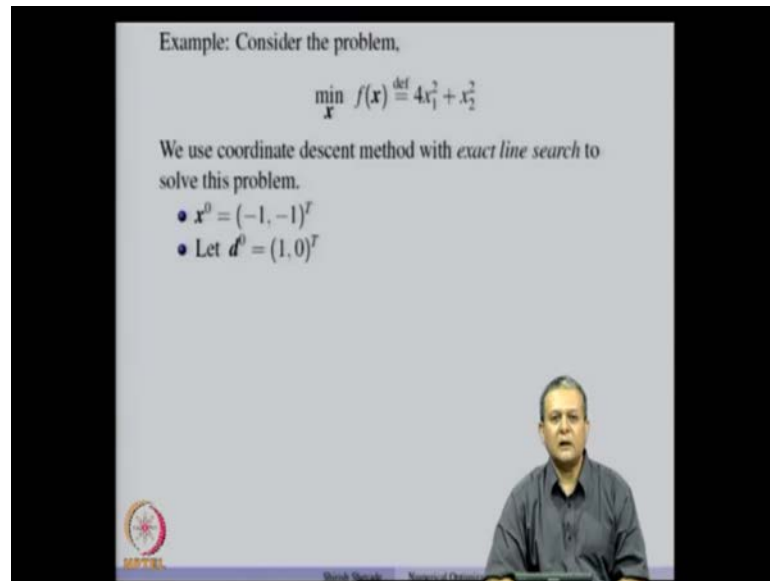
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Example: Consider the problem,

$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

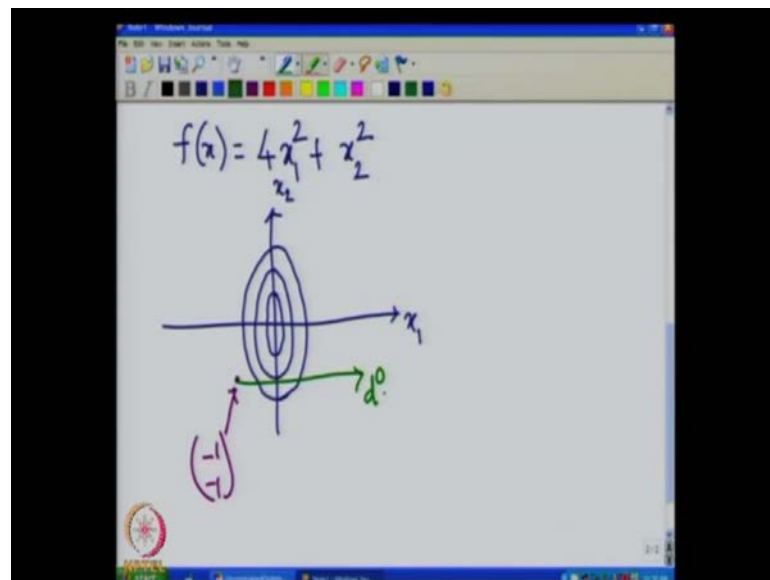
We use coordinate descent method with *exact line search* to solve this problem.

- $\mathbf{x}^0 = (-1, -1)^T$
- Let $\mathbf{d}^0 = (1, 0)^T$



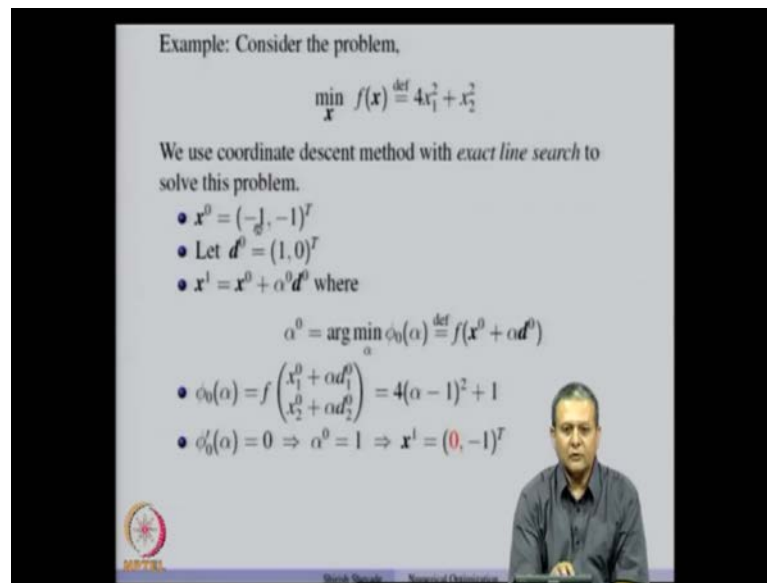
Now, let us choose the direction first direction to be 1 comma 0. So, that is the x_1 direction. So, initially we make a progress along the direction \mathbf{d}^0 and then we use the next coordinate direction which is 0 1. So, initially we choose the direction.

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So, this is \mathbf{d}^0 which is the x_1 coordinate.

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Example: Consider the problem,

$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

We use coordinate descent method with *exact line search* to solve this problem.

- $\mathbf{x}^0 = (-1, -1)^T$
- Let $\mathbf{d}^0 = (1, 0)^T$
- $\mathbf{x}^1 = \mathbf{x}^0 + \alpha^0 \mathbf{d}^0$ where

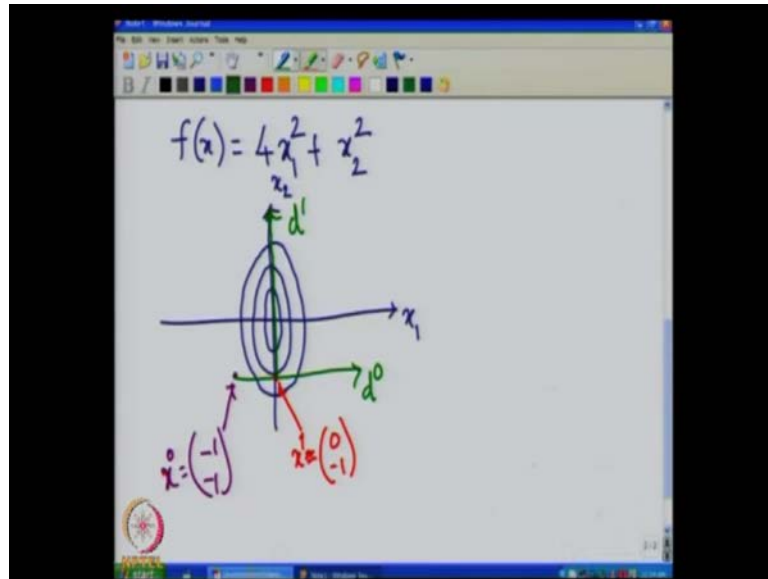
$$\alpha^0 = \arg \min_{\alpha} \phi_0(\alpha) \stackrel{\text{def}}{=} f(\mathbf{x}^0 + \alpha \mathbf{d}^0)$$

- $\phi_0(\alpha) = f \begin{pmatrix} x_1^0 + \alpha d_1^0 \\ x_2^0 + \alpha d_2^0 \end{pmatrix} = 4(\alpha - 1)^2 + 1$
- $\phi_0'(\alpha) = 0 \Rightarrow \alpha^0 = 1 \Rightarrow \mathbf{x}^1 = (0, -1)^T$

Then, we find \mathbf{x}^1 to be \mathbf{x}^0 plus $\alpha^0 \mathbf{d}^0$ using our usual update for \mathbf{x}^1 and α^0 is obtained by solving another optimization problem which is minimizing a function f of \mathbf{x}^0 plus $\alpha \mathbf{d}^0$. So, we know \mathbf{d}^0 remember that, this α is should be greater than 0 greater or equal to 0, when we do this minimization. So, by solving this problem we get α^0 . Now if we explicitly write the $\phi_0(\alpha)$ to be f of \mathbf{x}^0 plus $\alpha \mathbf{d}^0$ and x_2^0 plus αd_2^0 . So, these are the two coordinates for f and if we use this \mathbf{x}^0 and \mathbf{d}^0 substitute them here and plug them in this formula. So, what we get is an expression in terms of α .

So, remember that α is the only variable for this function ϕ and this ϕ is define using f like this where \mathbf{x}^0 and \mathbf{d}^0 are known. So, finally, we get an expression in terms of only α now it is clear that the minimum of $\phi_0(\alpha)$ with respect to α is obtained when α when the derivative of ϕ with respect to α vanishes at a point α^0 equal to 1 and therefore, what we get is \mathbf{x}^1 equal to 0 minus 1 on this is the point on the y axis. So, you will see that initially we started with minus 1, minus 1 and is which was this direction \mathbf{d}^0 to be our initial search direction and when we did the exact line search we move to the point \mathbf{x}^1 where this first coordinate become 0 second coordinate remained as it is.

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So, you will see that we reach the point. So, this was x^0 and this was x^1 to be 0 minus 1 . So, this is our current point and then we have to choose the direction d^1 . Now the second coordinate direction that we have to choose and that direction is d^1 . Now let us see how to do that optimization.

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Example: Consider the problem,

$$\min_x f(x) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

We use coordinate descent method with *exact line search* to solve this problem.

- $x^0 = (-1, -1)^T$
- Let $d^0 = (1, 0)^T$
- $x^1 = x^0 + \alpha^0 d^0$ where

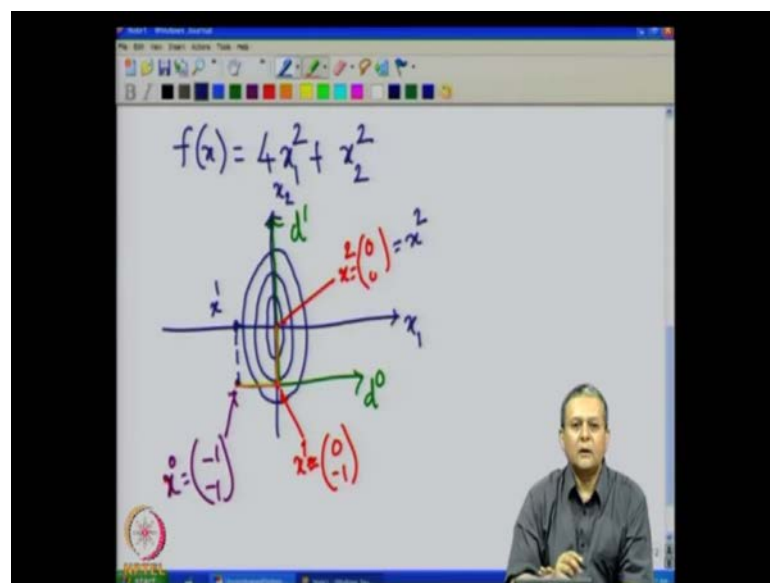
$$\alpha^0 = \arg \min_{\alpha} \phi_0(\alpha) \stackrel{\text{def}}{=} f(x^0 + \alpha d^0)$$

- $\phi_0(\alpha) = f \begin{pmatrix} x_1^0 + \alpha d_1^0 \\ x_2^0 + \alpha d_2^0 \end{pmatrix} = 4(\alpha - 1)^2 + 1$
- $\phi_0'(\alpha) = 0 \Rightarrow \alpha^0 = 1 \Rightarrow x^1 = (0, -1)^T$
- $d^1 = (0, 1)^T$, $x^2 = x^1 + \alpha^1 d^1$, $\alpha^1 = \arg \min_{\alpha} \phi_1(\alpha) \stackrel{\text{def}}{=} f \begin{pmatrix} 0 \\ \alpha - 1 \end{pmatrix} = (\alpha - 1)^2 \Rightarrow \alpha^1 = 1 \Rightarrow x^2 = (0, 0)^T = x^*$

Now, we use the next coordinate direction which is $0 \ 1$ and we know that x^2 is obtained by this formula x^1 plus $\alpha^1 d^1$ where α^1 is obtained by minimizing from function. So, let us see what that function is, so ϕ_1 of α is nothing but minimum

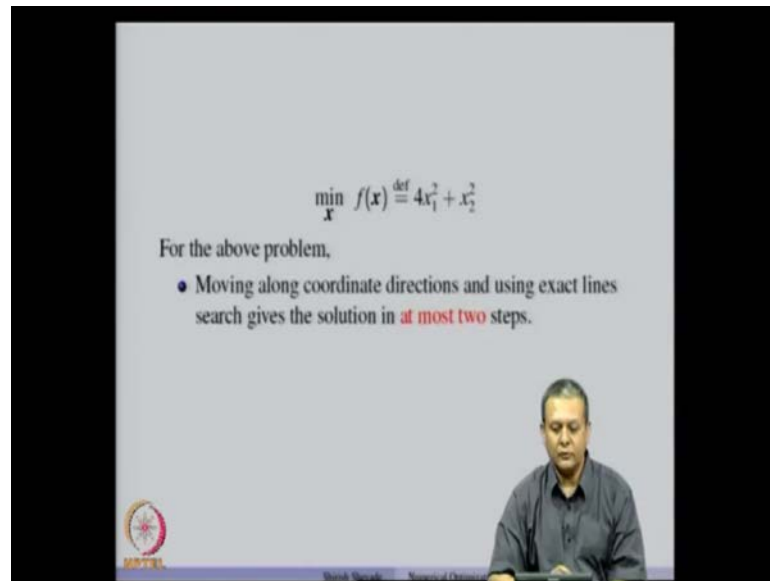
of. So, if we write x_2 to be x_1 plus αd_1 and we know x_1 and we know d_1 plug in those values of x_1 and d_1 in x_1 plus αd_1 we get the 2 coordinates as 0 and α minus 1. So, if you plug in these coordinates herein the original expression what we get is α minus 1 square and if we minimize this with respect to α that is take the derivative of with respect to α set it to 0. We get α equal to 1 and once we get α equal to 1 then substitute it here we have x_1 which is 0 minus 1 and d_1 which is 0 1. So, if you combine them what we get is x_2 equal to 0 comma 0 are the origin and we know that origin is the solution of this problem.

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So, we get a new point which is our x_2 and that is nothing but 0 0. So, the path truest by this method is. So, initially we move up to this point to x_1 and then we move to the point x_2 0 now suppose we had use the alternate directions like initially suppose we have we had moved along d_1 and then along d_0 then you would notice that we would we would have got the similar behavior. So, in that case we would have got this point as our x_1 and this point as our x_2 . So, you will see that even if you alternate our directions of search in this case d would still get the same point.

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The slide displays the following content:

$$\min_x f(x) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

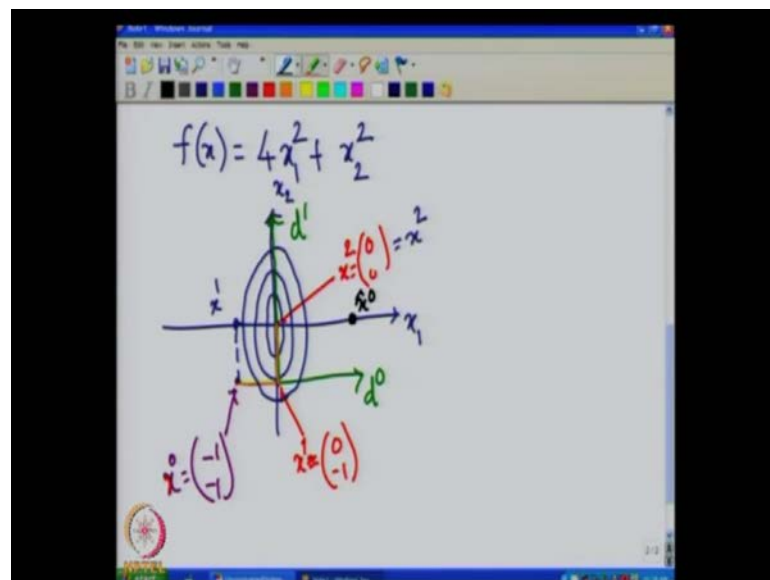
For the above problem,

- Moving along coordinate directions and using exact lines search gives the solution in **at most two steps**.

A man in a dark shirt is visible in the bottom right corner of the frame, appearing to be presenting the slide.

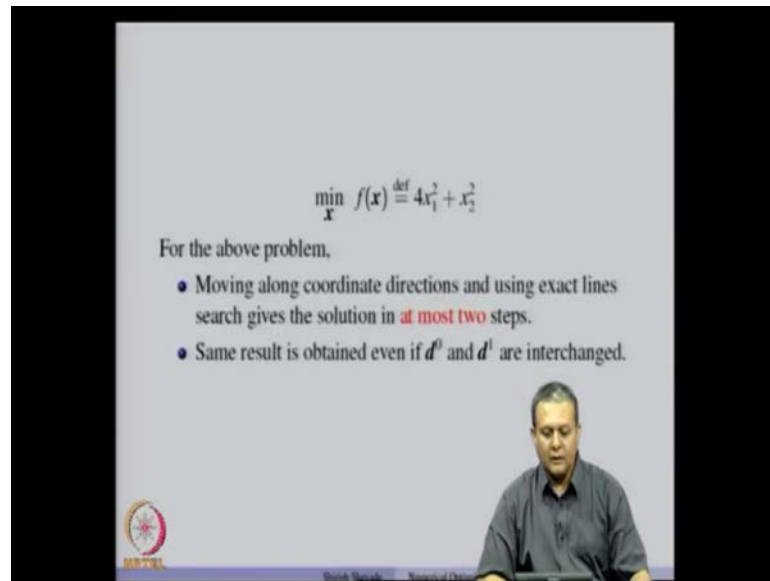
Important point to remember is that if you move along coordinate directions and used exact line search in this case we got the solution in at most two steps.

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Suppose, we are lucky and we start with initial point suppose this is our x^0 hat now you will see that in one iteration we could if you move along x_1 in one iteration we could reach the solution sometimes we could get the solution in one iteration if we are along the principle axes of this elliptical contours but otherwise we would require at the most two steps to reach the solution.

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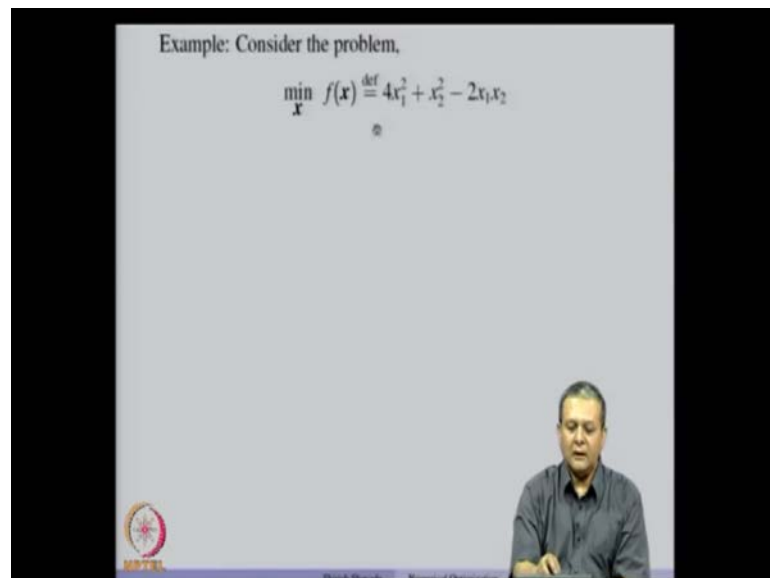
$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

For the above problem,

- Moving along coordinate directions and using exact lines search gives the solution in **at most two** steps.
- Same result is obtained even if d^0 and d^1 are interchanged.

Now, we will get the same solution x^* even if we alternate our search directions that is if we initially search along the direction d^1 and then along the direction d^0 we would still get the same final solution x^* so that order of the descent directions is not going to matter in this case. So, these are the some important points that one can remember.

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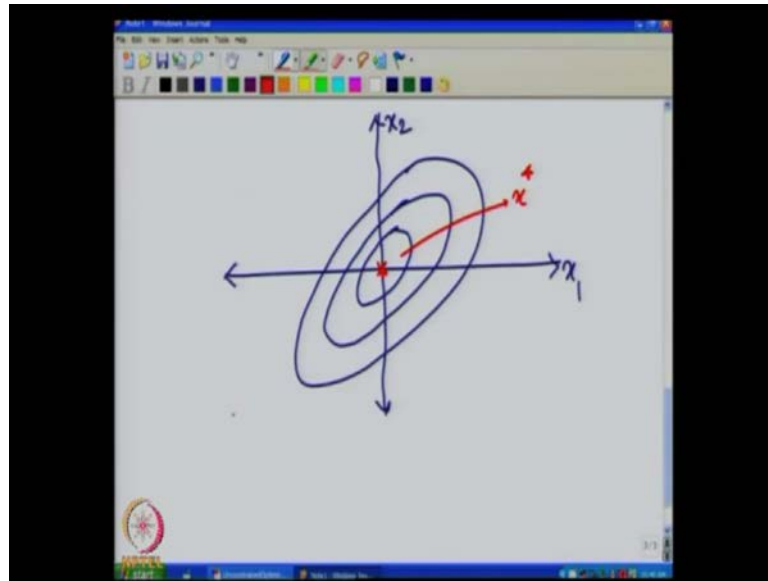


Example: Consider the problem,

$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2 - 2x_1x_2$$

Now, let us consider a case where we have a function $f(x)$ defined as $4x_1^2 + x_2^2 - 2x_1x_2$. So, let us draw the contours of this function.

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So, we have seen that the contours of this function they look like this x_1 and x_2 the contours look the like this and this is going to be the minimum of our problem and this is going to be our x^* . So, let us use the coordinate descent method to solve this problem.

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Example: Consider the problem,

$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2 - 2x_1x_2$$

We use coordinate descent method with *exact line search* to solve this problem.

- $\mathbf{x}^0 = (-1, -1)^T$
- Let $\mathbf{d}^0 = (1, 0)^T$
- $\mathbf{x}^1 = \mathbf{x}^0 + \alpha^0 \mathbf{d}^0$ where

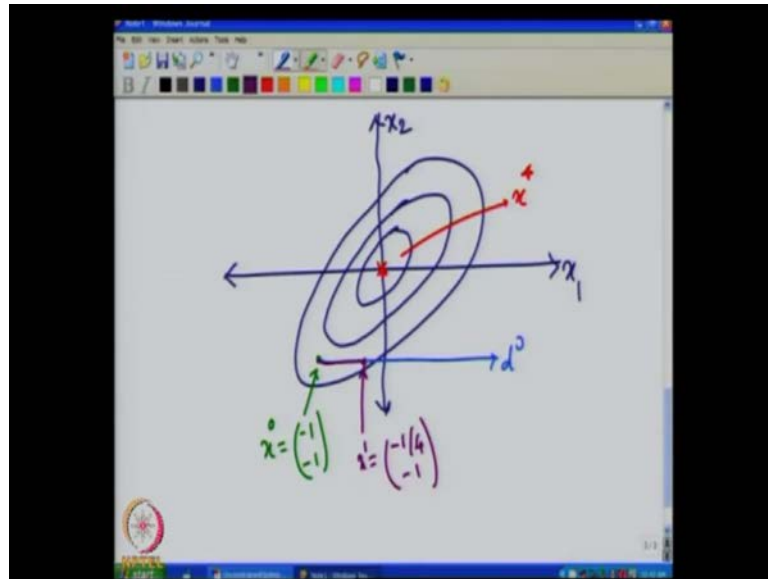
$$\alpha^0 = \arg \min_{\alpha} \phi_0(\alpha) \stackrel{\text{def}}{=} f(\mathbf{x}^0 + \alpha \mathbf{d}^0)$$

- $\phi_0(\alpha) = f \begin{pmatrix} x_1^0 + \alpha d_1^0 \\ x_2^0 + \alpha d_2^0 \end{pmatrix} = 4(\alpha - 1)^2 + 1 + 2(\alpha - 1)$
- $\phi_0'(\alpha) = 0 \Rightarrow \alpha^0 = \frac{3}{4} \Rightarrow \mathbf{x}^1 = \left(-\frac{1}{4}, -1\right)^T$

Again the function is quadratic we use the exact line search and suppose we start with the initial point minus 1 minus 1 and use the x_1 coordinate as the first coordinate along which we do the optimization first. So, we find x_1 to be using this formula x_0 plus $\alpha^0 \mathbf{d}^0$ where α^0 is obtained by minimizing $\phi_0(\alpha)$ which is nothing but f

of x_0 plus αd_0 . So, if you plug in these values of x_0 and d_0 . So, what we get is $\phi(\alpha)$ is the function of α and we will see that the minimum of this function occurs at α_0 equal to three-fourth by taking the derivative of this function derivative of this function with respect to α . And we get the point which is minus 1 by 4 minus 1.

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So, initial point suppose was somewhere here x_0 to be minus 1, minus 1 and when we moved along the direction. So, this was our d_0 . So, when we moved here and do the exact line search we got a point which is minus one-fourth minus 1. So, the point is some were here, so let us denoted by x_1 to be minus one-fourth minus 1. So, this is the initial path that was followed by this. Now, let us go to the next coordinate direction which is 0 1.

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Example: Consider the problem,

$$\min_{\mathbf{x}} f(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2 - 2x_1x_2$$

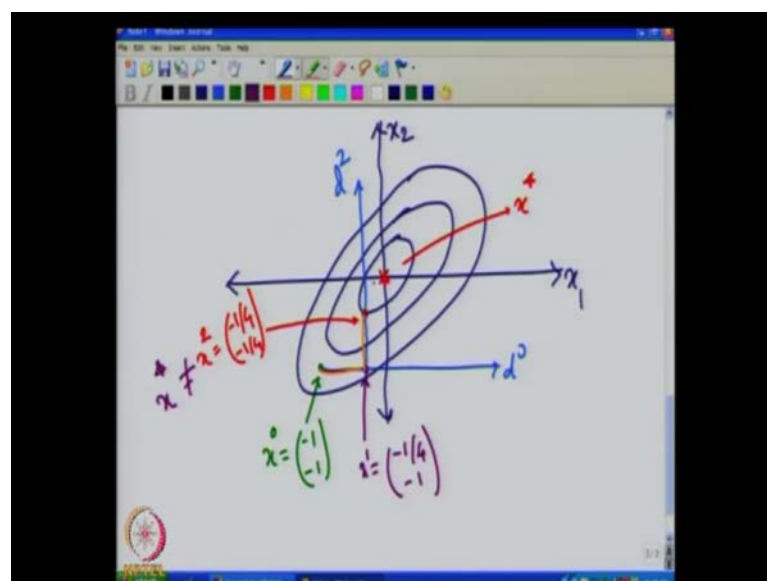
We use coordinate descent method with *exact line search* to solve this problem.

- $\mathbf{x}^0 = (-1, -1)^T$
- Let $\mathbf{d}^0 = (1, 0)^T$
- $\mathbf{x}^1 = \mathbf{x}^0 + \alpha^0 \mathbf{d}^0$ where

$$\alpha^0 = \arg \min_{\alpha} \phi_0(\alpha) \stackrel{\text{def}}{=} f(\mathbf{x}^0 + \alpha \mathbf{d}^0)$$
- $\phi_0(\alpha) = f \begin{pmatrix} x_1^0 + \alpha d_1^0 \\ x_2^0 + \alpha d_2^0 \end{pmatrix} = 4(\alpha - 1)^2 + 1 + 2(\alpha - 1)$
- $\phi_0'(\alpha) = 0 \Rightarrow \alpha^0 = \frac{3}{4} \Rightarrow \mathbf{x}^1 = (-\frac{1}{4}, -1)^T$
- $\mathbf{d}^1 = (0, 1)^T$, $\mathbf{x}^2 = \mathbf{x}^1 + \alpha^1 \mathbf{d}^1$, $\alpha^1 = \arg \min_{\alpha} \phi_1(\alpha) \stackrel{\text{def}}{=} f \begin{pmatrix} -\frac{1}{4} \\ \alpha - 1 \end{pmatrix} = (\alpha - 1)^2 + \frac{\alpha - 1}{2} + \frac{1}{4} \Rightarrow \alpha^1 = \frac{3}{4} \Rightarrow \mathbf{x}^2 = (-\frac{1}{4}, -\frac{1}{4})^T \neq \mathbf{x}^*$

And we know that \mathbf{x}^2 is obtained using this formula where α^1 is obtained by solving another minimization problem where $\phi_1(\alpha)$ is defined to be f of $\begin{pmatrix} -1/4 \\ \alpha - 1 \end{pmatrix}$ and if you plug in this value in the original formula. So, we get this expression for $\phi_1(\alpha)$. So, we take the derivative of this $\phi_1(\alpha)$ with respect to α set it to 0 and we get α^1 to be three-fourth and if you plug in this value of α^1 here then along with this value of \mathbf{x}^1 and this value of \mathbf{d}^1 what we get is \mathbf{x}^2 and \mathbf{x}^2 what we get is $\begin{pmatrix} -1/4 \\ -1/4 \end{pmatrix}$ which is not equal to \mathbf{x}^* .

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So, if you look at our. So, we use the direction d_2 ; d_2 is the direction here and we will see that we reach a point which is somewhere here and this point is x_2 that is minus one-fourth minus, one-fourth and we will see that we haven't reach the solution in two steps. So, the path test by this algorithm in the first two iteration is this and this and then you will see that we need some more iterations to reach the solution.

So, this x_2 is not equal to x^* in this case. So, obviously will need more iterations to reach the solution x^* in this case now we compare this with our previous function. So, you will see that in the previous function starting from 1 particular point we reach the solution in exactly two steps. Now, in the other case two steps were not enough to reach the solution if we use coordinate descent method. So, what is the problem? So, let us try to answer that.

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• Example 1:

$$\min_{\mathbf{x}} f_1(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2$$

- $H = \begin{pmatrix} 8 & 0 \\ 0 & 2 \end{pmatrix}$.
- x^* , attained in *at most two steps* using coordinate descent method

• Example 2:

$$\min_{\mathbf{x}} f_2(\mathbf{x}) \stackrel{\text{def}}{=} 4x_1^2 + x_2^2 - 2x_1x_2$$

- $H = \begin{pmatrix} 8 & -2 \\ -2 & 2 \end{pmatrix}$.
- x^* , could not be attained in two steps using coordinate descent method (if x^0 is not on one of the principal axes of the elliptical contours)

Take the function $f(x)$ the first function. So, we call it $f(x) = 4x_1^2 + x_2^2$ and if you look at the hessian matrix of this function the hessian matrix is 8 and 2 along the diagonals and non-diagonal elements 0 and in this case the solution x^* the attained at most two steps using coordinate descent method now if you look at the other function that we saw which is $4x_1^2 + x_2^2 - 2x_1x_2$ we will see that the hessian matrix here there are some off diagonal elements as well and these off diagonal elements are there because of this coupling between x_1 and x_2 . So, such a term did not exist in the earlier case and therefore, the hessian matrix was nice diagonal matrix while

here it is not a diagonal matrix and we saw that unless we choose x_1 to be one of the principle axes of the elliptical contours we cannot raise the solution in two steps if we use coordinate decent method. So, this term results in the non-diagonal hessian matrix and if we look at this two functions $f_1(x)$ and $f_2(x)$ and compare them you would see that it is easy to minimize x_1 keeping x_2 fixed here and similarly, minimize x_2 keeping the x_1 fixed because the terms involving x_1 and x_2 are separable.

So, here is a term which contains only x_1 here is a term which contains only x_2 . So, we have separable objective functions separable in terms of x_1 and x_2 while here you will see that this term does not make the objective function separable in terms of x_1 and x_2 and it is this term which makes the hessian matrix also non-diagonal. So, the reparability of the objective function in terms of its variables is lost when we use when we have such a term which couples the different coordinates and that because of that we are not able to achieve the minimum in exactly two steps in this two-dimensional case if you use coordinate method.

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Consider the problem:

$$\min_x f(x) \triangleq \frac{1}{2} x^T H x + c^T x$$

where H is a symmetric positive definite matrix.

- Let $\{d^0, d^1, \dots, d^{n-1}\}$ be a set of linearly independent directions and $x^0 \in \mathbb{R}^n$
- Any $x \in \mathbb{R}^n$ can be represented as

$$x = x^0 + \sum_{i=0}^{n-1} \alpha^i d^i$$

- Given $\{d^0, d^1, \dots, d^{n-1}\}$ and $x^0 \in \mathbb{R}^n$, the given problem is to minimize $\Psi(\alpha)$ defined as,

$$\frac{1}{2} \left(x^0 + \sum_{i=0}^{n-1} \alpha^i d^i \right)^T H \left(x^0 + \sum_{i=0}^{n-1} \alpha^i d^i \right) + c^T \left(x^0 + \sum_{i=0}^{n-1} \alpha^i d^i \right)$$

So, let us consider a general function again a quadratic function which is half of x transposes h x plus c transpose x where h is a symmetric positive definite matrix now. So, remember that the hessian matrix h need not be diagonal. So, let us choose some n directions which are linearly independent and let us choose some initial point x_0 now any x in the inputs space which is n -dimensional space can be written using x_0 and

linear combinations of these n independent directions because they form a basis for this n -dimensional space as these vectors are linearly independent. So, they form a basis. So, any x in the n -dimensional space can be written as x_0 plus some linear combination of these vectors.

So, suppose x is written as $x_0 + \sum_{i=0}^{n-1} \alpha_i d_i$ where i varies from 0 to $n-1$. Now let us rewrite this problem in terms of α . So, let us assume that d_0 to d_{n-1} are known quantities, x_0 is also a known quantity. So, d_0 to d_{n-1} are known, x_0 is known and what is unknown is α . So, let us rewrite this original problem by substituting x in this objective function by $x_0 + \sum_{i=0}^{n-1} \alpha_i d_i$ where, i varies from 0 to $n-1$ and let us that problem will be a problem where the variables are α s and not x . So, we are given d_0 to d_{n-1} and x_0 in \mathbb{R}^n we rewrite the original problem as minimize $\psi(\alpha)$ where $\psi(\alpha)$ is defined as this. So, you will see that every x in this original equation is replaced by $x_0 + \sum_{i=0}^{n-1} \alpha_i d_i$ now this is the quadratic function. So, h is constant. So, h does not depend on x . So, h remains as it is then c is the scalar that also does not depend on x . So, that remains as it is and we have a function $\psi(\alpha)$ in terms of the variable α .

So, the original problem becomes the problem of minimizing $\psi(\alpha)$ with respect to α . Note that there are n α s. So, the number of variables here was n the number of variables in this new problem is also n and suppose we solve this problem with respect to α will get a solution α^* and if we use that α^* in this equation. So, if we find out $x_0 + \sum_{i=0}^{n-1} \alpha_i^* d_i$ then we get x^* which is going to be the solution of this problem. Note also that we are working with the symmetric positive definite matrix. So, this function is strictly convex and therefore, they exist a strict local minimum and that local minimum x^* is in this case obtained by solving another problem $\psi(\alpha)$ with respect to α and getting α^* now the reason for doing this will become clear when we expand this quantity take out the constant terms and make some assumptions on d 's. So, we will do that in the next class.

Thank you.