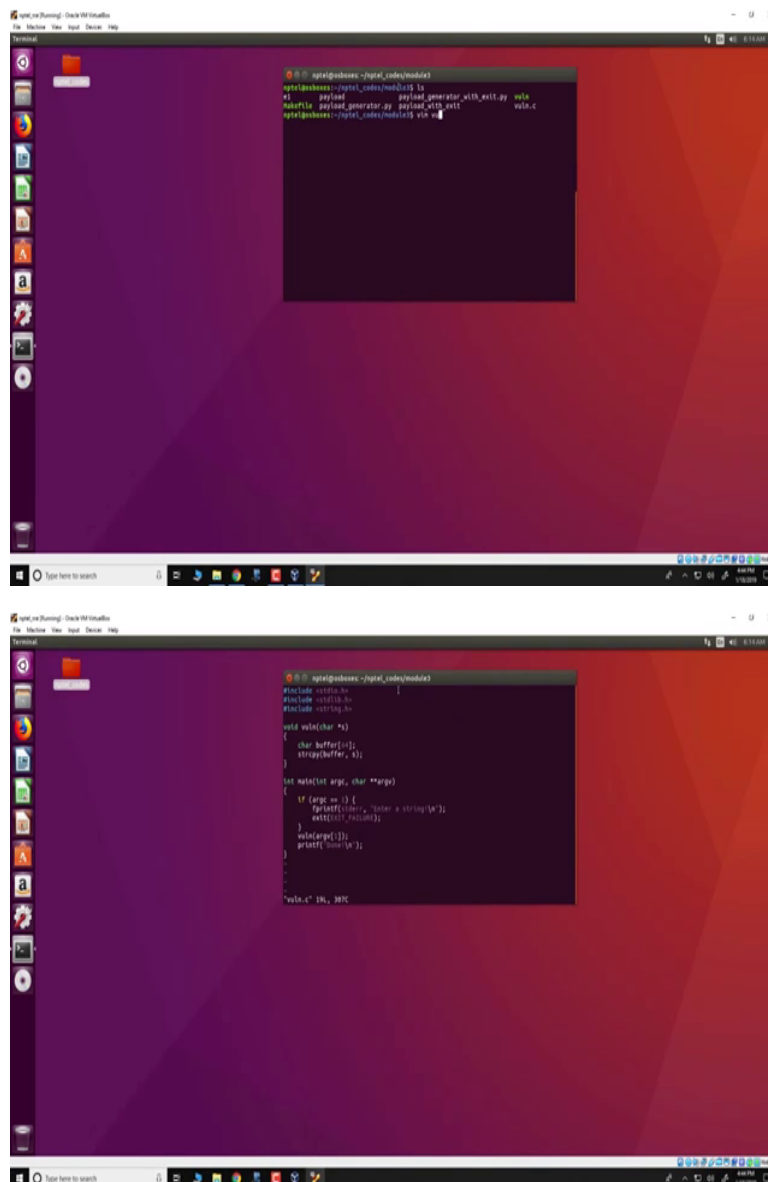


Information Security - 5 - Secure Systems Engineering
Prof. Chester Rebeiro
Indian Institute of Technology Madras
Demonstration of a Return Oriented Programming (ROP) Attack

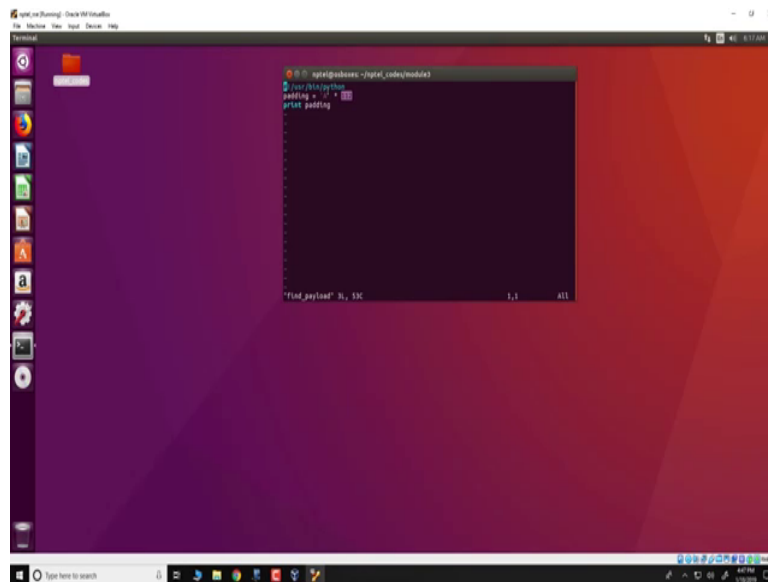
Hello and welcome to this demonstration in the course for secure system engineering, in one of the previous videos we seen how attackers use the return to libc attack to overcome the fact that this stack was made non-executable with a return to libc attack even though we have a stack which is non-executable, still an attacker would be able to overflow a buffer and do something malicious, in the demonstration that we are see today, we will be creating a shell using the return to libc attack.

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So see this program working first, so we will run it as follows give it a short string and we see the fact that it is actually completed, all the other hand if we give the program a very large string like this, which is much larger than 64 bytes, then as expected buffer will overflow due to the long string copy which is done and we would get a segmentation fault, now a very similar to the what we have seen in the previous demonstration, where we seen a buffer overflow on the stack, in order to generate the payload you will use the Python strip which you have used in the previous demonstration, so we just copy it from there on model 2, find_payload over here and as using before this particular strip would generate strings of any arbitrary length.

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```
april@ubuntu:~/ipat_code/module3
└─$ python3 find_payload.py
padding = 112
print padding

"find_payload" N = 512      1,1      413
```

For example over here since we have specified 112, so this particular python strip would print A 112 A, so by varying this 112, we could actually get different lengths for the string.


```
agari@ubuntu:~/ngpt_codes/modules$ ./p
GNU gdb (Ubuntu 7.11.1-ubuntu14.1) 7.11.1
Copyright (C) 2016 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <http://gnu.org/licenses/gpl.html>
This is free software; you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law. Type 'show copying'
and 'show warranty' for details.
This GDB was configured as 'x86_64-linux-gnu'.
Type 'show configuration' for configuration details.
For bug reporting instructions, please see:
<http://www.gnu.org/software/gdb/bugs/>.
Find the GDB manual and other documentation resources online at:
<http://www.gnu.org/software/gdb/documentation/>.
For help, type 'help'.
Type 'apropos word' to search for commands related to 'word'...
Reading symbols from ./vuln_1.c...done.
(gdb) b main
Breakpoint 1 at 0x808040c: File vuln_1.c, line 13.
(gdb) r AAA
Starting program: /home/agari/ngpt_codes/modules/vuln_1.c AAA
Breakpoint 1, main (argc=2, argv=0xfffff8be) at vuln_1.c:13
warning: Source file is more recent than executable.
13      if (argc == 1) {
14          p system
15      }
(gdb) p system
```

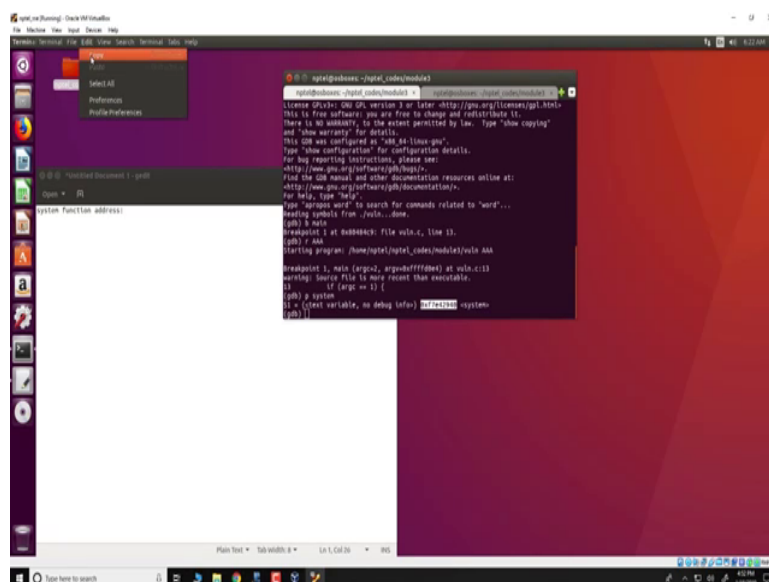
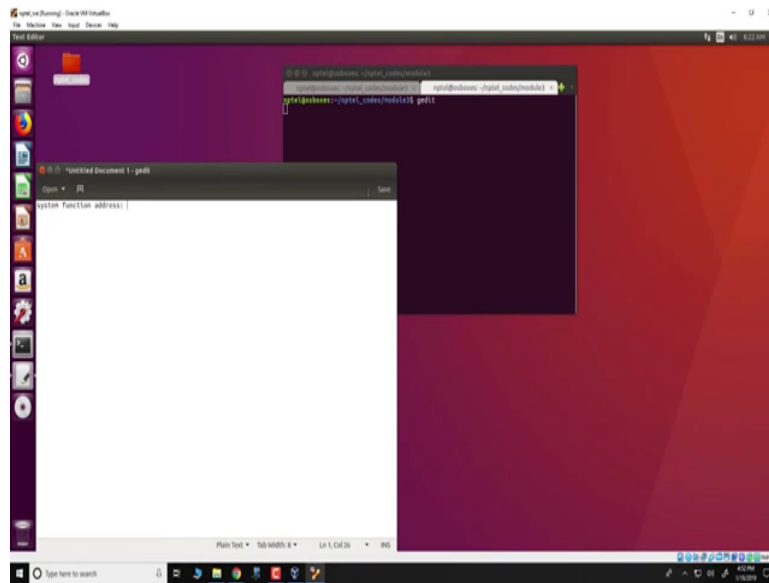
```
agari@ubuntu:~/ngpt_codes/modules$ ./p
License GPLv3+: GNU GPL version 3 or later <http://gnu.org/licenses/gpl.html>
This is free software; you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law. Type 'show copying'
and 'show warranty' for details.
This GDB was configured as 'x86_64-linux-gnu'.
Type 'show configuration' for configuration details.
For bug reporting instructions, please see:
<http://www.gnu.org/software/gdb/bugs/>.
Find the GDB manual and other documentation resources online at:
<http://www.gnu.org/software/gdb/documentation/>.
For help, type 'help'.
Type 'apropos word' to search for commands related to 'word'...
Reading symbols from ./vuln_1.c...done.
(gdb) b main
Breakpoint 1 at 0x808040c: File vuln_1.c, line 13.
(gdb) r AAA
Starting program: /home/agari/ngpt_codes/modules/vuln_1.c AAA
Breakpoint 1, main (argc=2, argv=0xfffff8be) at vuln_1.c:13
warning: Source file is more recent than executable.
13      if (argc == 1) {
14          p system
15      }
(gdb) p system
$1 = 0x00000000, no debug info) 0xfffff8be <system>
(gdb)
```

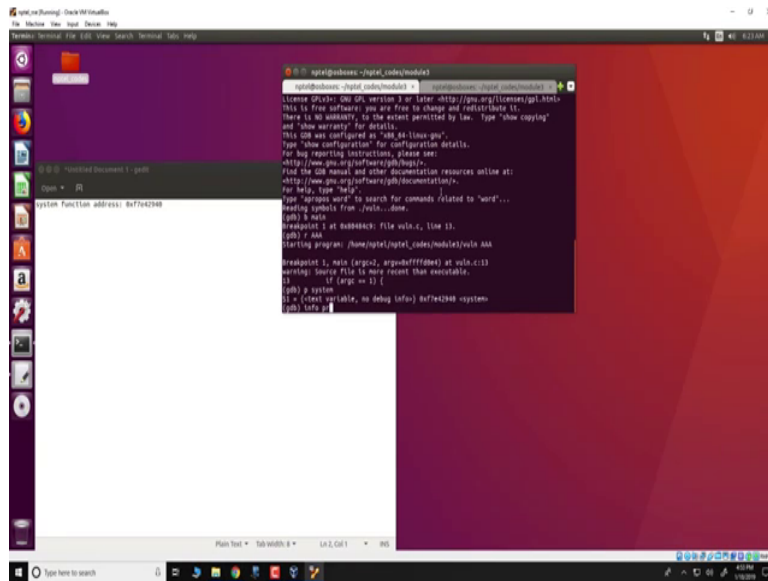
So for example if I increase this A from 72, where 72 actually just modifies the frame pointer present on the stack and I make it 76, then it is the firm pointer that gets modified as well as the return address, so in order to create this payload for the return to libc attack as we have seen in the previous lecture, what we would require is to fill the buffer so that the return address is modified and the return address is modified with a pointer to the system function, which is present in the libc, also what is required is a string argument to this system function, so we will required string such as /bin/sh or /bin/bash, which is a string comprising of an executable.

So what is expected to happen is that when system executes, it would tick the string from the memory and execute that corresponding executable okay, so let us try to build this return to libc attack, so there are two things that we would require, we would require the address of

where system resides the entire process and we had also need to find somewhere in the process of the string/bin/sh is present, so let us do this we will do this by using the GDB debugger, will do, run GDB./1 and put a breakpoint at main and run this file with some argument, we obtained the breakpoint but what we are really interested in is to find the address of these libc function for system and that we can obtain as follows, we can do a P system that print system and we see that this address over here 0XX7E42940 is the location of the system function in the entire process.

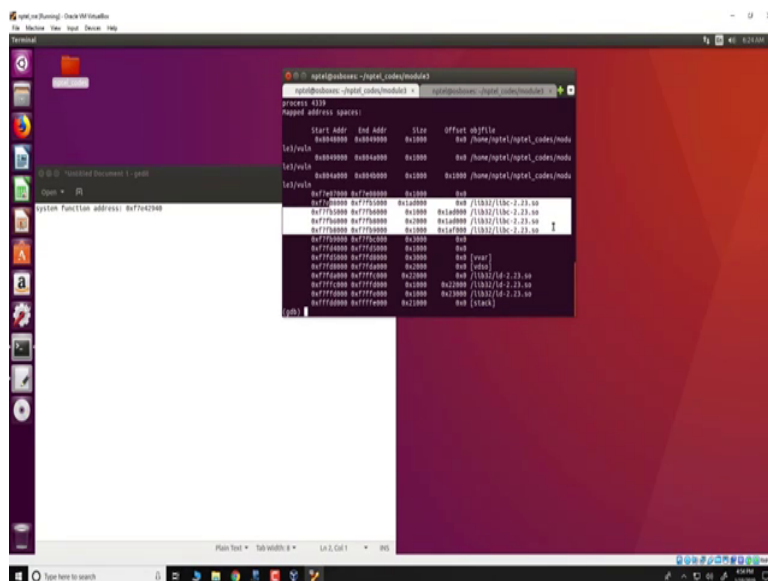
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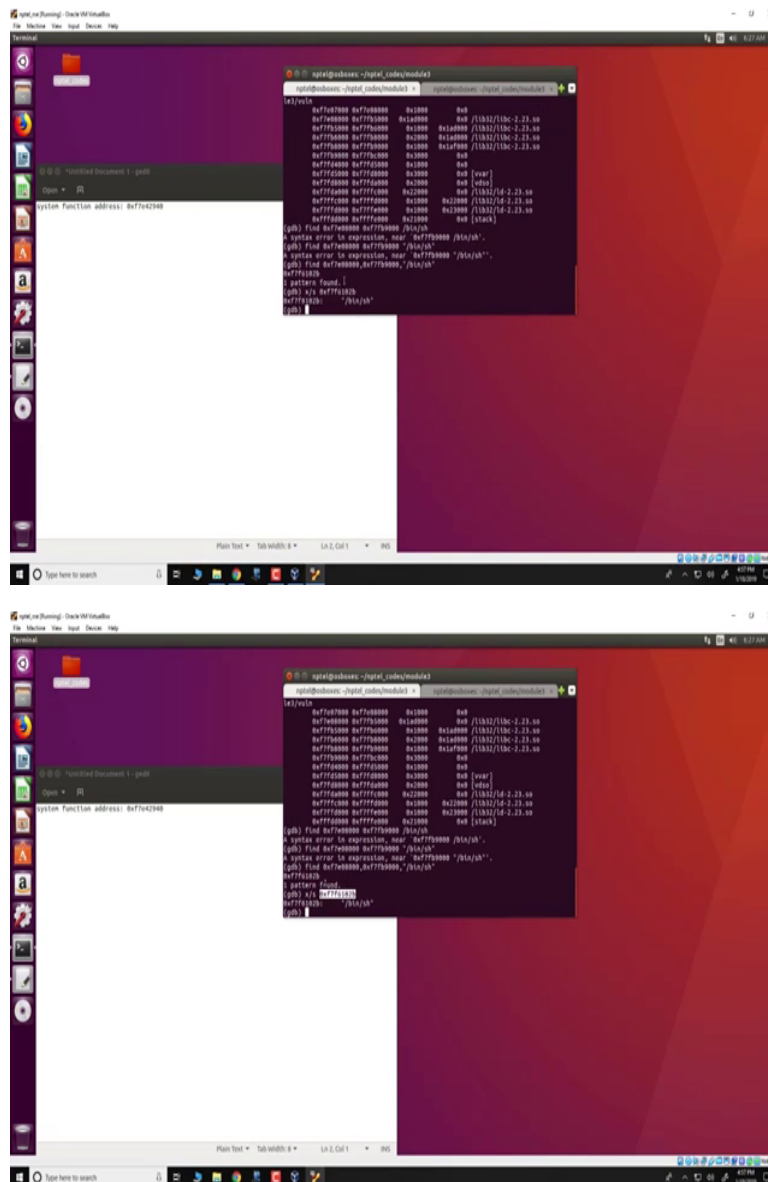
So let us create a copy of this, so that is, okay the next thing we need is the address of the string/bin/sh, we could also use the string/bin/bash which is a pointer, which is a string comprising of the bash executable but making it work that we would be a little more difficult, on the other hand we actually use and create a SH shell that is /bin/sh and we need to find somewhere in the entire process, where /bin/sh is present, so we do is we try to search in the various paths of this process memory, so one thing what we know for sure is that /bin/sh is present in the libc, so if you know the memory range where a libc is present in the entire process space, we could search that memory area and identify the address of /bin/sh okay, so the first thing we need to do is find where libc is present.

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So we do this using the command `info proc map` and we see that `libc` is present starting at location `F7E08000` to `F7FB5000` and in same similar way there are other areas, other 4kd pages comprising of `libc`, so we will try to search in this entire memory range for the required string, GDB supports search searching across memory by using the `find` command, so command starts like this, so it is `find 0XF7E08000`, so this is the starting address for `libc` that is obtained from here and will search up to the end address this one `F7FB9000`, yes you can see as there was small a syntax error essentially `find` requires the start address, the end address and the string to be search all separated by commas, so what `find` would do is search in this memory range for this acquired string and what we see here is that it has found one pattern located at this address `F7F6102B`.

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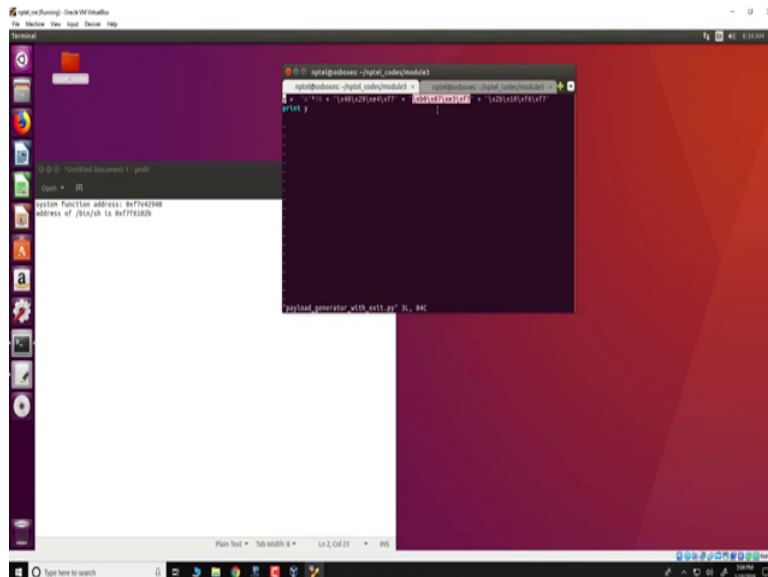
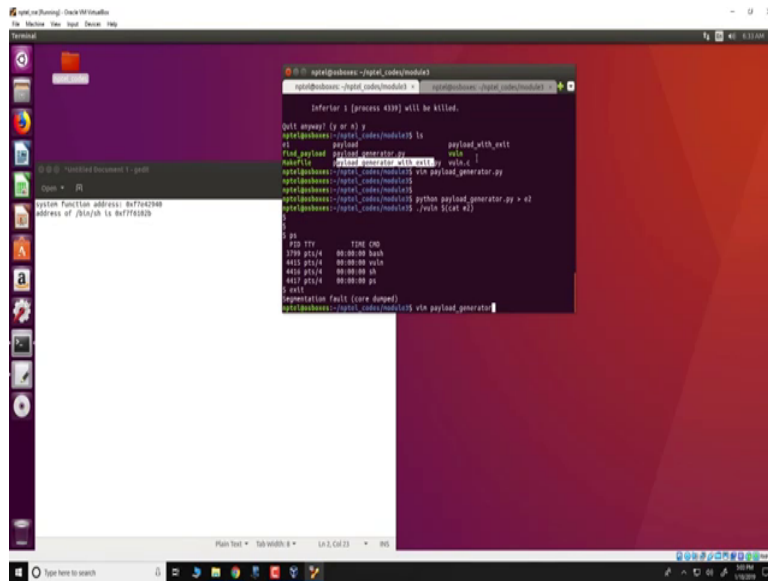


Let us verify that this memory location actually has the required string, so we can do that by dumping the string in using this command `X/S0XF7F6102B` and we see that the string present `/bin/bash` or `/SH`, will record the string, address of `/bin/sh` is here okay, so this is all that we can need for a very basic return to libc attack, so how we will exit GDP, next thing we will do is to create a payload, so in order to create a payload we written a small Python script known as payload generator.

So will open that and what we see here is that this particular payload generator is print a string `Y` which comprises of `76 A`, so the `76 A` are used to overflow the buffer and as we and you can reconnect that if we overflow by `76 A`, the next four strings would be overriding the return address, so in little Indian format will put the address of the system function that is `F7E42940`.

As you can see here this matches the address of system function which we have just found, next really were of 4 bytes in the stack, so we just fill it with some arbitrary values, in this case we put `A` but we could actually to fill it with anything and then we put the address of the string `/bin/sh`, so this address `F7F6102B` which matches the address that we have actually found for `/bin/sh`.

So what is going to happen now is that when the function actually returns is going to return to this particular address which we are specified over here and which would be present in the return address location in the stack and this would trigger the system function in libc to execute, the system function would then tick from the stack, the only argument that it requires and this argument is a character pointer and it is expecting a pointer to a string comprising of an executable, so it uses this address as the argument and picks the string `/bin/sh` and executes it.



So the first thing we need to do is run the payload generator and create the required payload, so we do this as follows, so we run Python payload generator.py and store the result in this file call E2 and then we run vull again and give the input using act E2 and give the input from the file E2, so what we see now is due to this maliciously formed string, it would trigger the system to execute and create the shell as before is missing with the regular buffer, this shell is what is present over her with the PID 4416.

And also what you see is that there are other processes present in the background, first is the bash which is the original process comprising of this particular tunnel, then you have the vull executable which has a process ID 4415 and it is still running in the background, we have SH and of course this particular process PS which we have just used.

So what the system function does is that it forks a process and then executes that particular process, so essentially the child of this one process is the shell that we have just created, now the one process is locked until the SH process completes its execution, so when we terminate this SH process what we see is that the one process will continue to execute, so let us see this happening so we execute the SH but what we will see is that the one process will have a segmentation fault okay.

So this is because the one process is looking in these stack and it takes out some garbage or some arbitrary values from the stack and tries to execute that and it is causing this, the process to actually crash and as we have seen in the theory, a of better way of doing it is to safely exit that particular program, as an assignment you could try to modify this wonderbility that we have created, this wonderbility that we have exploited and try to safely exit from the shell, so in order to do that we also created this Python script payload generator with exit, which creates the string not only comprising of the system function that we have here and its corresponding argument but also the exit function, so this address you would require to fill with the address of the exit function, which would be present somewhere in the libc memory region of the process. Thank you.