

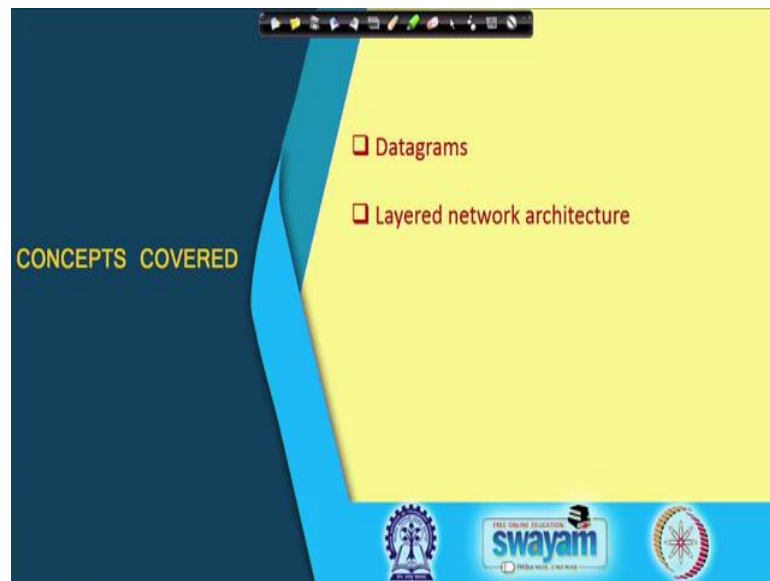
Ethical Hacking
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Lecture - 03
Basic Concepts of Networking (Part II)

So, we continuing with our discussion on the Basics of Computer Networking. So, this is the second part of this lecture on basic concepts of networking. Now, if you recall in our last lecture, we had talked about the broad approaches of data transmission over a network.

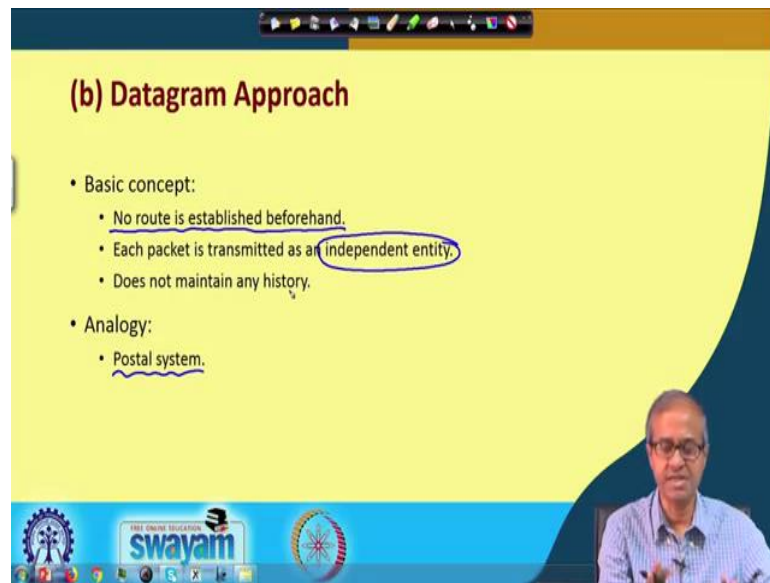
We are talking about something called packet switching, we mentioned the virtual circuit approach, where some dedicated path gets established, a fixed path get established called a virtual circuit, all packets will be following that path. Today we shall be starting with the other alternative which is the so called datagram approach.

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So, in this lecture we shall be talking about first the datagrams, some issues relating to that and then we shall be talking about layered network architecture, which is a very commonly used popular architecture for designing networking systems ok.

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(b) Datagram Approach

- Basic concept:
 - No route is established beforehand.
 - Each packet is transmitted as an independent entity.
 - Does not maintain any history.
- Analogy:
 - Postal system.

The slide features a yellow background with a blue header and footer. The footer includes the 'swayam' logo and other educational icons. A man with glasses is shown in a small video window in the bottom right corner of the slide.

Let us look at the datagram approach, well, if you compare with the virtual circuit approach, the main difference comes here that no route is established beforehand. This is one difference or advantage you can say, you have here advantage in the sense that there is no initial delay. In virtual circuit because of the route establishment, there was an initial delay and once the route is established, then only the packets can be transmitted. Now, here because there is no route established, so each packet is transmitted as independent entities.

Just like you think of a postal system as an analogy, I am talking about a postal system. When you drop a letter in a letter box, each letter is sent or is being you can say despatched as an independent entity. So, you do not know which path the letter is following or actually how it is going ok. So, that you do not know, but one thing you know is that, that when you post a letter, you specify some kind of an address, their destination address. So, here also when you transmit a packet, the packet must carry the address of the destination node along with it.

Because, it is being sent as an independent entity, you must have the address ok. And, it does not maintain any kind of history like, there is no concept of a path that is established and so on. So, in that sense, there is no history of data communication or transmission that is maintained in this case. Each packet is being routed as an independent entity fine.

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Datagram Approach (contd.)

- Every intermediate node has to take routing decisions dynamically.
- Makes use of routing table
- Every packet must contain source and destination addresses.
- Problems:
 - Packets may be delivered out of order.
 - If a node crashes momentarily, all of its queued packets are lost.
 - Duplicate packets may also be generated.

The diagram shows three nodes: A, B, and C. Node A is connected to Node B, which is then connected to Node C. Node B has multiple outgoing links, indicating multiple possible paths. A packet header is shown with a box labeled 'H' containing 'SA' (Source Address) and 'DA' (Destination Address).

So, here the difference is, here every intermediate node will be taking some dynamic decisions, it will be maintaining something called a routing table. So, you see if you think of a scenario, where there is a node A, there is a node B, and there is a node C. So, you are sending data from node A to node B and then node B to node C.

Now, this node B will be having some kind of a table that it maintains, the table contains some relevant information. And what kind of information it contains? It contains information like, like whenever an incoming packet comes, so, it will contain information like, where this packet has to be forwarded to ok.

There can be multiple outgoing links from B and to take this decision, what is required is that when you are thinking about 1 packet, I mentioned that each packet will be containing a header. In this header there must be a source address and also a destination address. Just like a letter, when you post a letter, you give the address where you want to send the letter to and in case the letter cannot be delivered, you also give your own address, so that the letter can come back to you.

In a same way in the header there are other information also, but the most important one is source address, the person who is sending and the destination address where you are sending right. Now, just like the normal postal system, here there can be some problems. Like, when you post letters, there is no guarantee that the order you are posting the letters, the letters will get delivered in the same order, there is no guarantee regarding that,

courtesy of our postal system or the way the letters are sorted and delivered. So, packets may be delivered out of order, why? Because each packet is being routed as an independent entity, there is no guarantee that all packets are flowing or following the same path. Some packet may be following a shorter path, some packet may be following a longer path. So, the delays of each packet can be different ok.

So, the delivery of the packet can be out of order. Similarly, if some of the intermediate node temporally goes down, all packets which are stored in buffer, will get lost. So, some packets might get lost also. This also happens in a postal system, sometimes a letter never gets delivered, but something which happens here, which is not there in a postal system, is that, sometimes duplicate packets may get generated.

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Datagram Approach (contd.)

- Every intermediate node has to take routing decisions dynamically.
 - Makes use of a *routing table*.
 - Every packet must contain *source and destination addresses*.
- Problems:
 - Packets may be delivered out of order.
 - If a node crashes momentarily, all of its queued packets are lost.
 - Duplicate packets may also be generated.

The slide features a hand-drawn diagram in blue ink showing three nodes (circles) connected by arrows, illustrating a non-sequential path. A video inset in the bottom right corner shows a man with glasses speaking. The slide also includes logos for 'swayam' and 'Free Online Education' at the bottom.

Like for every link, suppose there is a link between a node to the next node. This node will be sending some packets to the next node and will be expecting some acknowledgement to come, that well I have received the packets correctly.

If the acknowledgement does not come, this sender will assume, there is some something wrong, let me send it again ok. But it may so happen that because of some other problem, the acknowledgement got delayed, but if this sender sends the same packet again, there will be duplicate which will be generated right. So, some duplicate packets might get generated ok.

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Datagram Approach (contd.)

- Advantages:
 - Faster than virtual circuit for smaller number of packets.
 - ❖ No route establishment and termination.
 - More flexible.
 - Packets between two hosts may follow different paths.
 - ❖ Can handle congestion/failed link.

The diagram shows a network topology with nodes A, B, C, D, E, F, G, and H. Node A is the source on the left, and node H is the destination on the right. There are multiple paths from A to H: A-B-D-F-H, A-B-D-G-H, A-C-E-G-H, and A-C-E-H. The slide is part of a Swayam presentation, as indicated by the logo at the bottom.

Now, talking about the advantages of this approach, datagram approach; first thing is that if the number of packets is not very large, this will be faster than the virtual circuit approach, because in the virtual circuit approach, there is an initial overhead of connection establishment. But once the connection establishment is over, packet transmission will be very fast. So, if there are large number of packets, then virtual circuit can be faster overall ok.

Datagram approach does not need any route establishment in the beginning or termination at the end, it does not require these. So, in that sense it is more flexible. Packets can follow different paths, if sum of the links get overloaded, then the packets can follow an alternate path, these flexibilities are there ok. So, this as I mentioned, it can handle congestion in the links or some of the links may fail altogether, then an alternate link may be used which was not there for a virtual circuit all right.

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Comparative Study

- Three types of delays must be considered:
 - a) Propagation Delay
 - Time taken by a data signal to propagate from one node to the next.
 - b) Transmission Time
 - Time taken to send out a packet by the transmitter.
 - c) Processing Delay
 - Time taken by a node to process a packet.

The slide also features a diagram of a communication link between two nodes, with a handwritten label '1Mbps' in a circle in the middle of the link. At the bottom of the slide, there is a banner for 'swayam' (Free Online Education) and a small video inset of a man speaking.

Now, if you just compared these approaches which we have talked about, now for comparison we usually compare with respect to some parameters; one is called propagation delay. Like you see when we talk about a communication between two nodes; propagation delay is the signal delay from source to the destination. How much time a single signal transition takes to move from this place to this place.

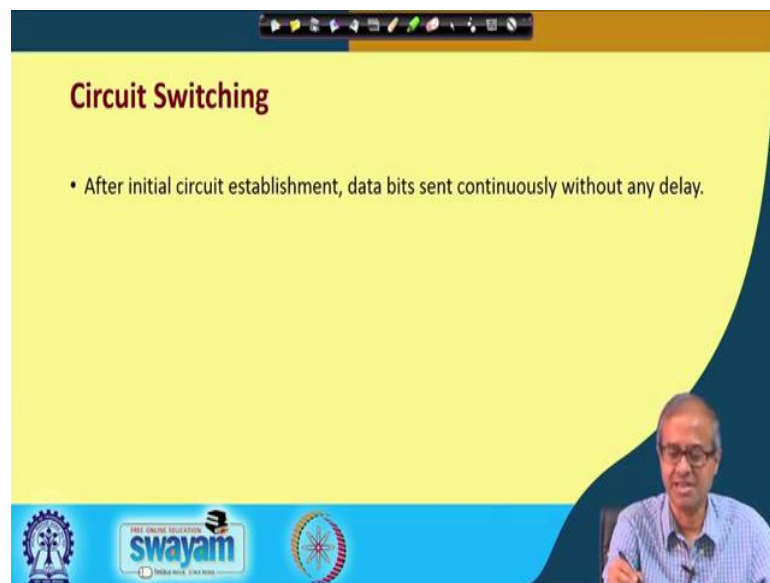
If, it is a copper link, it can be something, if it is a fibre optic link, it can be different, if it is a satellite link, it can be different. So, depending on the type of link, how much time suppose you send one bit of data? How much time that one bit will take to reach the other end, it depends on the medium. Suppose, you are using a satellite, then it will be taking a fraction of a second to go up to the satellite and again come back ok.

So, it depends on the medium, but transmission time is different. Transmission time depends on the bandwidth of the link, well you know that links are categorised by the speeds, kilobytes per second, megabits per second, gigabits per second. Suppose it is a 1 megabit per second, link 1 Mbps. So, 1 million bits can be sent per second. So, if there are 1000 bits you are trying to send, so, here how much time will it take? Here every bit will take 1 microsecond 1 by mega.

So, 1000 bits will take one millisecond that is the transmission time. It depends on the size of the message, how many bits are there and the bandwidth of the link. It does not depend on the propagation delay; whether it is a satellite link or a copper link whatever.

So, these two are two separate things and finally, processing delay. If your communication is going via intermediate nodes and you know, here we use something called store and forward approach. So, these intermediate nodes will also take some time to receive a packet, store it, consult the routing table, take some decision and then forward it. So, it will take some small time for that. So, that is the processing delay. So, these are the different components of delay in a typical network communication.

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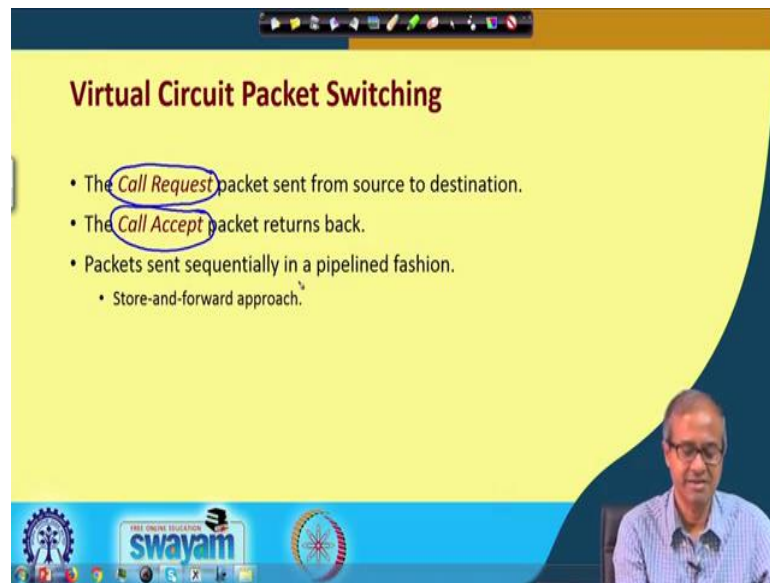
Circuit Switching

- After initial circuit establishment, data bits sent continuously without any delay.

Logos at the bottom: IIST, swayam, and a circular logo.

Now, talking about circuit switching, here only this initial circuit establishment phase is there, but once this phase is over, this delay is over; you have a dedicated connection from source to destination, data transmission can proceed in the maximum speed continuously without delay. Because, you know, you have a guaranteed bandwidth, you can send the bits continuously.

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Virtual Circuit Packet Switching

- The Call Request packet sent from source to destination.
- The Call Accept packet returns back.
- Packets sent sequentially in a pipelined fashion.
 - Store-and-forward approach.

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Now, talking about virtual circuit, well during the initial virtual circuit establishment phase a special packet is usually sent, I am not going into the detail, which is called call request packet. This call request packet is responsible for establishing the virtual circuit and setting up the routing tables of all the intermediate nodes, corresponding to this particular virtual circuit number ok.

Now, once the virtual circuit has been established and acknowledgement signal or packet will be traversing along the same path back, this is called a call accept packet. So, that the sender and all the intermediate nodes will know that the circuit has been established. Now, the packets can flow, now the packets can be sequentially one after the another in a pipeline fashion using store and forward approach, this is how virtual circuit will work.

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The image shows a presentation slide with a yellow background and a dark blue header and footer. The title 'Datagram Packet Switching' is in a bold, dark red font. Below the title, there is a bulleted list of characteristics. In the bottom right corner, there is a small video inset showing a man with glasses and a blue shirt. The footer contains the 'swayam' logo and other icons.

Datagram Packet Switching

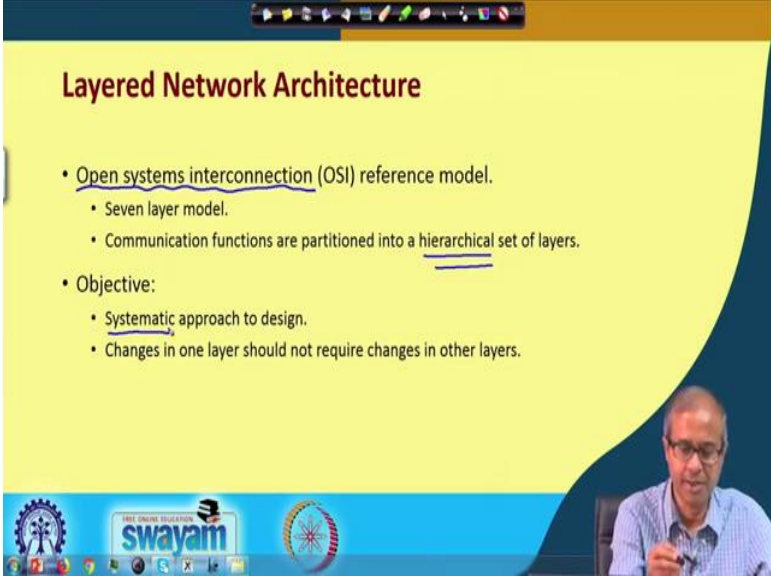
- No initial delay.
- The packets are sent out independently.
 - May follow different paths.
 - Also follows store-and-forward approach.

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Now, talking about the datagram packet switching, because there is no connection establishment phase, so there is no initial delay ok. The packets are sent out independently, may follow different path and this also follows store and forward concept, but if the number of packet is very large, then the total time required here maybe more as compared to virtual circuit, but here you have lot of flexibility.

You can share the links, if there are some link failures, link congestion, you can overcome those problems. So, those flexibilities are there and because of that, this datagram approach is the most widely used approach in the internet today ok.

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Layered Network Architecture

- Open systems interconnection (OSI) reference model.
 - Seven layer model.
 - Communication functions are partitioned into a hierarchical set of layers.
- Objective:
 - Systematic approach to design.
 - Changes in one layer should not require changes in other layers.

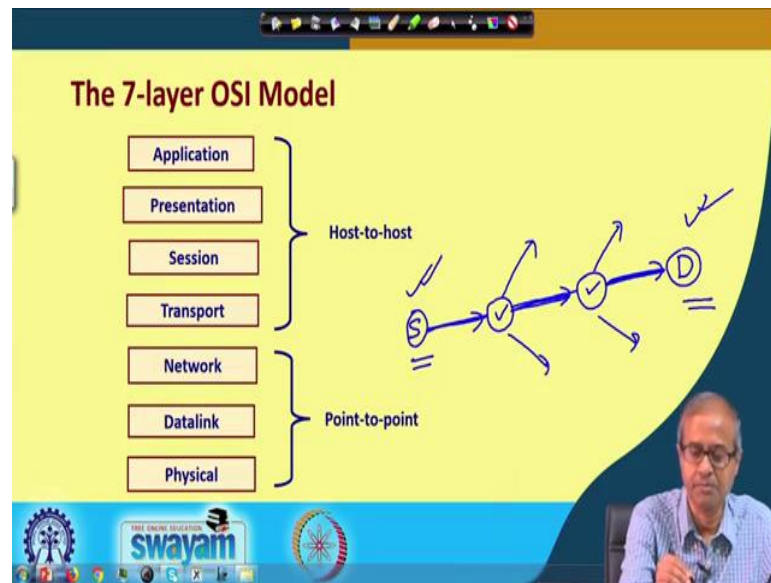
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Now, let us talk about something called layered network architecture. Now, see a networking system has a number of responsibilities. Broadly speaking it is data communication, but if you look into the detail, there are a number of steps which are involved in the process. Now, I mean in order to design a networking system in a flexible way, a conceptual layering has come in. Layering means you divide the total functionality into a number of levels, where the function of each of the level is very well defined ok.

Now, a very well-known model is the so, called OSI model. This is open systems interconnection model, this was proposed by international systems organisation, ISO. This is called ISOs OSI model, it is called, there are 7 layers in this model. I shall show the layers very shortly and there is a hierarchical relationship among these layers. The main objective of this layering is, as I mentioned, if you divide the functionality into well-defined layers, it becomes much more systematic.

Your design of the networking software becomes much more systematic. Not only that you may choose to modify the design of one of the layers without touching the other layers. So, one of the layers may use a different technology, but you do not modify the other layers, that way you can say the maintainability of the entire networking system becomes also better ok. These are the some of the advantages.

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Now, talking about these 7-layers, you see these are the 7 layers; starting from the lowest point physical, data link, network, transport, session, presentation, application. Now you see, if you think of a network like this, well I am only showing one path, there can be other paths also, I am just showing one path.

Let us from a source S to a destination D this is the path which is followed. Now, you see some of these layers are called host to host and some of the layers are designated as point to point, host to host means the upper 4 layers, they will exist only in the source and the final destination.

They will not exist in the intermediate nodes. Point to point means 2 nodes which are directly connected by a link is called a point to point connection. So, there is 1, 2 and 3. There are 3 point to point connections here. There are 3 point to point links and the lower 3 layers physical data link and network, they will be active in all these intermediate nodes also, but the upper 4 will be active only in the source and the destination, I will show a diagram ok.

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Layer Functions

- **Physical**
 - Transmit raw bit stream over a physical medium.
- **Data Link**
 - Reliable transfer of frames over a point-to-point link (flow control, error control).
- **Network**
 - Establishing, maintaining and terminating connections.
 - Routes packets through point-to-point links.

Application
Presentation
Session
Transport
Network
Datalink
Physical

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Now, first let me briefly talk about the function of the different layers. The lowermost layer the physical layer, this actually concerns about the electrical or means in whatever way you are sending, either using electrical signals or using optics, fibre optics, using optical signals, this concerns the transmit of the raw bits that constitutes your data over the channel, whatever channel it is. Some kind of physical medium, it can be a copper, it can be a fibre optic cable, it can be a microwave link, it can be Bluetooth, it can be anything ok. This is a physical layer.

Physical layer establishes a physical connection between two directly connected link points. Data link layer ensures that link is reliable to some extent, it tries to detect some errors in communication, if some error takes place, it will try to retransmit the data again ok. Data link layer is responsible for reliable transfer of data, now at this level, the data unit is called a frame over a point to point link well. It also takes into account flow control; that means, if there is a speed mismatch, it can slow down or make it faster, the rate of data transfer.

If there are some errors, it will try to handle error, it will try to retransmit that frame again and so on. Now, in the next higher layer which is the network layer, here the packet routing or the transmission of the packets is the main concern. Establishing, maintaining terminating connections and routing packets, there may be several point to point links, but I want to transmit a packet from let us say myself to your computer, but there can be 100

intermediate computers. So, all the network layers in this 100 computers must be responsible to forward the packet in the right direction, so that ultimately it can reach you ok. That is the responsibility of the network layer.

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Layer Functions (contd.)

- **Transport**
 - End-to-end reliable data transfer, with error recovery and flow control.
- **Session**
 - Manages sessions.
- **Presentation**
 - Provides data independence.
- **Application**
 - Interface point for user applications.

Application
Presentation
Session
Transport
Network
Datalink
Physical

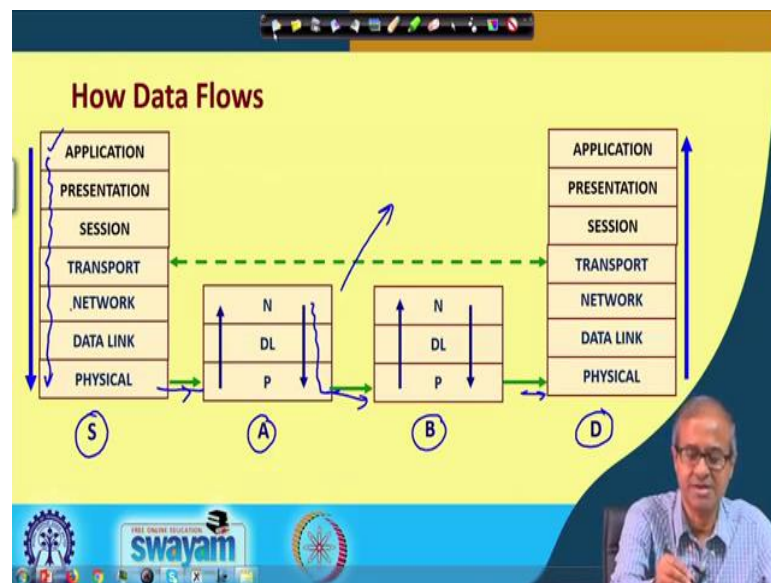
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Now, the upper 4 layers, these are end to end layer; that means, the point to point links are not aware of this layers, it is directly between the source and the final destination. Transport layer talks about end to end reliable data transfer like, I am concerned with, I am trying to send a data to you, I am asking you, have you received the data correctly, you tell me no, then I will send the data again. I am not worried about what intermediate nodes are there. This is a direct host to host link. Session, session layer is responsible for managing sessions; you know that there are many applications where there is a concept of a session.

You do some kind of a login, the session remains active during the period of login, then you logout, the entire period is called, referred to as a session. The session layer is responsible for maintaining such sessions. Presentation layer is used, is an optional layer of course, sometimes if you are using some kind of codes which are not compatible, you can do some code conversion, you can do some error detection correction, error detecting correcting code, some encryption, you can, you can add something extra in this presentation layer before you can transmit the data.

Like I want that my data transmission secure, I can do some encryption, encrypt my data before I can actually send, that will be done in the presentation layer. And application layer is almost any application you can think of some application, which wants to use the network for data transmission right.

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Now, the way data flows, it can be clear from this picture, you see this is the source, this is the final destination. So, the data starts from the application, it flows through these 7 layers, application to presentation, to session, to transport, network, data link, physical. And, once it reaches the physical layer, it actually gets physically transmitted to the next node A.

In the next node, it goes to the data link layer to check whether there is an error to network layer, to check what to forward it next, then it takes the decision again, comes back to the physical layer and forwards it to the right next one, because there can be some other node also connected to this. So, it will decide what to forward it.

So, again this node B will have this physical, data link and network layer, it will go up to that again, come down and it will in this way finally, reach the destination, in the destination data will flow in the reverse direction from physical to data link to network up to application.

So, roughly this is how data flows through this network so called stack. Because, the order in which the data is going down, the reverse way it is going up that is it is like a stack. So, this is sometimes also referred to as a network stack.

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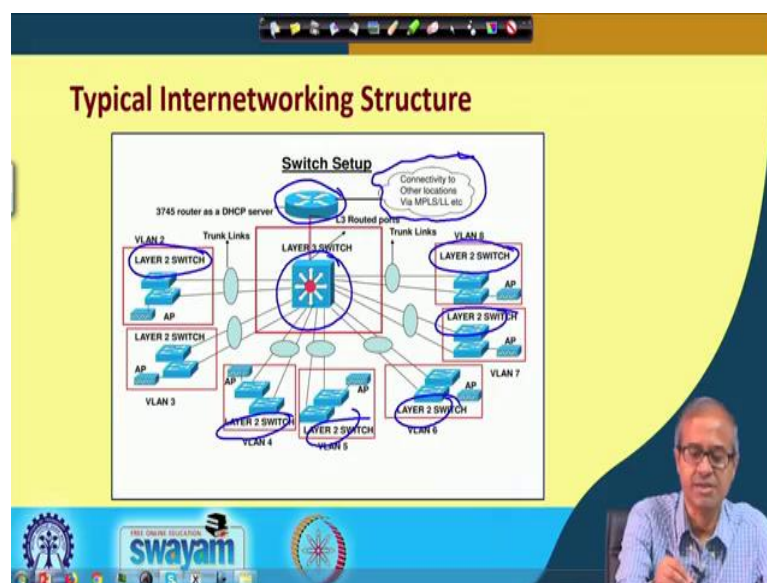
Internetworking Devices

- **Hub**
 - Extends the span of a single LAN.
- **Bridge / Layer-2 Switch**
 - Connects two or more LANs together.
 - Works at data link layer level.
- **Router / Layer-3 Switch**
 - Connects any combination of LANs and WANs.
 - Works at network layer level.

Now, talking about some of the internetworking devices, well we use hubs, bridges or layer-2 switches, routers or layer 3-switches, we shall come to these three things later. They are different, you can say equipments, their pictures are shown, you see pictorially they all look similar.

But, their functions are different. The first one is a hub, this is a layer 2 switch and this is a router or a layer 3 switch, their functions are different, layer 3 switch works at the IP layer level. So, we shall be discussing this later. Bridge works at the data link layer level and hub works at the physical link level.

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So, here I am showing a typical diagram that, how this networking devices can be interconnected in a real scenario? So, you see, this is your outside world, your connections to the outside world and you will be having a router here.

Typically routers are used to connect a network with other networks ok. And, inside your organisation there can be a main switch, this is a layer 3 switch, layer 3 switch is very similar to a router, this is internal to your organisation. And, in the different sections of the departments, you can have smaller switches, these are the layer 2 switches and the hubs I am not showing. So, under this layer 3 switches, you can have hubs and finally, you can make connections to the actual individual computers or the nodes. So, this is how a typical network looks like.

There will be hubs at the lowest level, layer 2 switches at the next level, layer 3 switches or routers at the next level, and to connect to the outside world, you will be finally having some routers which have connections to the outside world right. So, with this we come to the end of this lecture. We basically mentioned and talked about the networking layer, the 7 layer protocol and so on. In our next lecturer we shall be seeing some specific network protocol stack that is followed in the internet which is called TCP/IP.

TCP/IP is a very well-known protocol stack that is widely used in internet. In fact, the entire internet is based on this standard. So, we shall be looking into some details on that and other related issues.

Thank you.