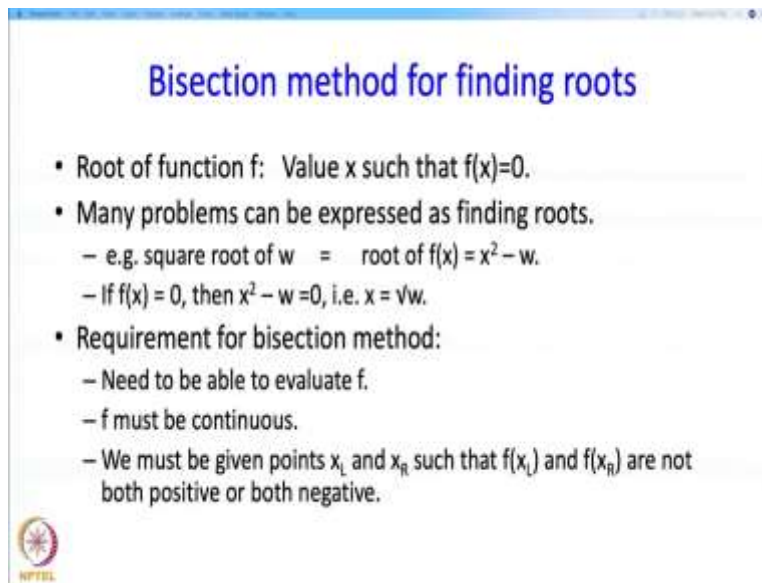


An Introduction to Programming through C++
Professor Abhiram G. Ranade
Department of Computer Science and Engineering
Indian Institute of Technology Bombay
Lecture No. 8 Part – 3
Computing Mathematical Functions
Bisection Method


Welcome back! In the previous segment, we discussed numerical integration. In this segment, we are going to discuss the bisection method for finding roots.

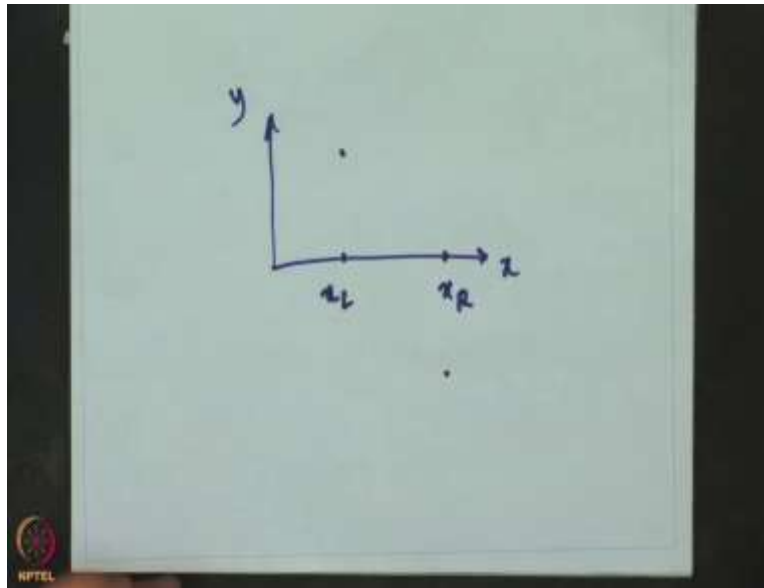
(Refer Slide Time: 0:31)



Bisection method for finding roots

- Root of function f : Value x such that $f(x)=0$.
- Many problems can be expressed as finding roots.
 - e.g. square root of w = root of $f(x) = x^2 - w$.
 - If $f(x) = 0$, then $x^2 - w = 0$, i.e. $x = \sqrt{w}$.
- Requirement for bisection method:
 - Need to be able to evaluate f .
 - f must be continuous.
 - We must be given points x_l and x_r such that $f(x_l)$ and $f(x_r)$ are not both positive or both negative.





So, the root of a function F is a value x such that $F(x)$ equal to 0. Or in other words, it is the point, where $F(x)$, the graph of $F(x)$ crosses the x -axis. Now, many problems can be expressed as finding roots. For example, if I want the square root of some number W , then I claim that is the same thing as the root of the function $F(x)$ equal to $x^2 - W$. Why is that? Well, if we do find root, that is if we find $F(x)$ such that, if we find x , such that $F(x)$ equal to 0, then we know that $x^2 - W$ is equal to 0. Or, x must be square root of W .

Now, the bisection method has relatively few requirements as to when it can be applicable. So for example, we require the ability to evaluate F , given an x . So, we should be able to calculate $F(x)$ given an x . Then, we want F to be continuous. And then, we must be given points x_l and x_r , such that $F(x_l)$ and $F(x_r)$ are not both positive or not both negative. So these are three relatively simple requirements, but once we have these then we can use the bisection method. And the bisection method can be used for finding roots. But, as we just remarked, finding roots is really itself useful for finding, for calculating functions, such as say the square root in this case, okay. So, again to reiterate, we need to be able to evaluate F . F must be continuous. And, we must be given points x_l and x_r , such that $F(x_l)$ and $F(x_r)$ are not both positive or both negative.

Or in other words, what must be the case? So, if we plot the graph of F , graph of F . So, this is the X -axis. This is the Y -axis. And, let us say, this is the point, this is the point with X co-ordinate x_l . This with x_r . Then, the points $F(x)$, the value of F at x_l , say it is

here, then x_r should be on the other side. Or, it could be at 0. But, both of them cannot be on the same side, okay. So that is what the point is. That is what the requirement is.

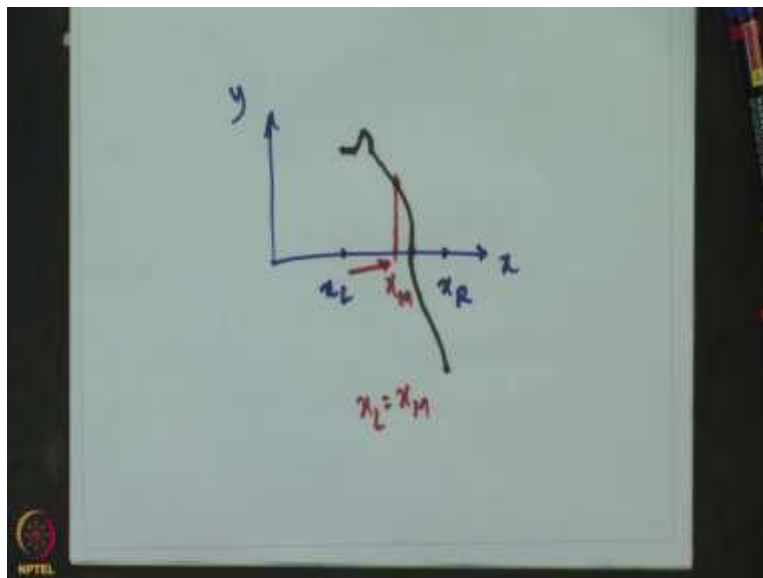

So, I am going to describe the bisection method now. But, we will assume these three properties. So, in other words, we will assume that we are given x_l and x_r satisfying this property.

(Refer Slide Time: 3:36)

Bisection method: basic iteration

- Basic iteration will bring x_l, x_r closer, while maintaining invariant:
 x_l, x_r have different signs or are 0.
- Invariant is true initially.
- Invariant + Continuity \Rightarrow root exists between x_l and x_r (both inclusive).
- We iterate till $x_r - x_l \leq \epsilon$, our desired error bound
- We declare x_l as the root.
- Error in declared root is at most ϵ .

Iteration
Let $x_M = (x_l + x_r)/2$
midpoint of interval (x_l, x_r) .
If $f(x_M)$ has same sign as $f(x_l)$,
then set $x_l = x_M$
Sign of x_l did not change
Sign of x_l continues to remain different from sign of x_r
else set $x_r = x_M$.
Sign of x_r does not change...
Invariant holds even here



So the basic iteration of the method is as follows. So we will start with the user given x_l and x_r . And then, we will bring them closer. As we bring them closer, we will maintain

the invariant, then, that x_l and x_r must have different signs or they are 0. If they are 0, it does not really matter. Now, I want to observe that at the beginning this invariant is true. That is because the user gave us two such, two such numbers. Now, before proceeding, let me discuss the implication of this. So, suppose $F(x_l)$ is indeed positive and $F(x_r)$ is negative. Now, because this function is continuous, what does it mean? The graph of the function must start somewhere over here. Maybe it will increase. Maybe it will do whatever it needs to, but it is continuous. So this graph, this line must be continuous. And being and while it is continuous, it must still reach this point. So, what does it mean? Then that means it must cross this X axis at some point. Otherwise it cannot go to the other side. But the point at which it crosses is exactly the root. So, what we know is that the root must be contained in this interval. So effectively when we said to the user please give us x_l and x_r . The user is effectively giving us an interval, which contains the root. So, invariant plus continuity implies that the root exists between x_l and x_r . And the root might be exactly at x_l or exactly at x_r . So, $F(x_l)$ or $F(x_r)$ are allowed to be 0. So, we are going to execute this basic iteration until the distance between x_r and x_l , becomes smaller than some error bound. Let us call it epsilon. If, epsilon is really small, then these two numbers are really close, and the root is somewhere in between them. Which really means that we can declare one of those numbers, say x_l as the root. So, we will have error in this. But, the error is at most epsilon.

So, that is the whole, that is the whole; that is the gist of the argument, okay. So we are going to get; we are going to get an approximate answer. But you will see that the approximation can be as good as you want. No matter what epsilon you give us, we will be able to iterate until that point.

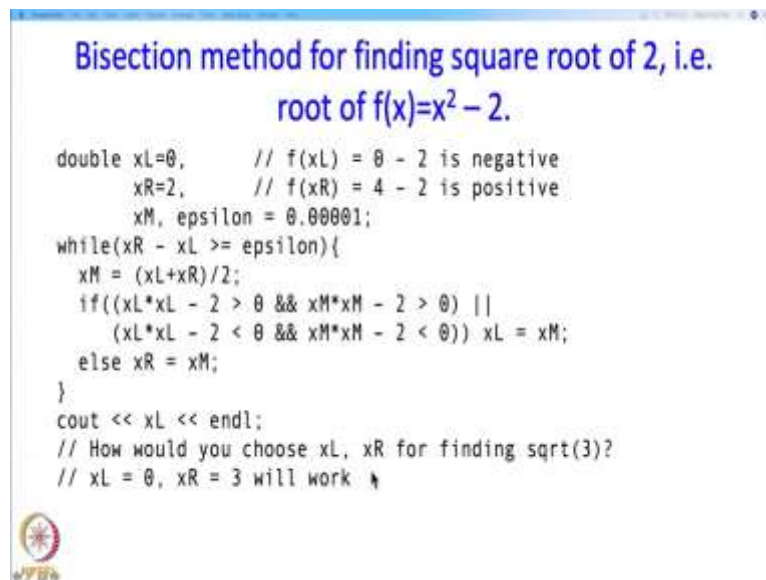
So, let me, let me tell you now, how this iteration is going to happen. How we are going to shrink this interval, okay. So, first we calculate x_m , which is $(x_l + x_r)/2$. So, what is that? Well, that is the midpoint of this interval x_l to x_r . Next, we check the signs of x_m and x_l . If, these two signs are the same, okay, then we are going to set x_l equal to x_m . So again, let us come back to this picture. So, this is x_l , x_r . We find the mid point, which is over here. Then, we look at $F(x_l)$, okay. So the sign of $F(x_l)$ and the sign, the sign of $F(x_m)$. So this is x_m . So the sign of $F(x_m)$ is positive. And the sign of $F(x_l)$ is positive.

So that means, $F(x_l)$ and $F(x_m)$ have the same sign. So in which case, we will execute x_l equal to x_m . Or in other words, we will move this point over to this point.

So, I claim now, that this is going to keep our invariant without violation. So, why is that? Well, x_l and x_m had the same signs. So the signs did not change as we moved x_l to x_m . As we assigned x_l equal to x_m , the signs did not change, okay. So that means the sign of x_l continues to remain different from the sign of x_r . So, if it was different before, then it will remain to continue to remain different. But, we are assuming that in, that this invariant was true in the beginning of the iteration. And therefore, what we have proved that in this case the invariant will hold even after the iteration. Or in other words, the root will now be contained in this smaller interval, which is true in this case, in our picture.

If, this condition is not holding that is if x_m and x_l do not have the same sign, then we are going to set x_r equal to x_m . So, why is that? Well, in that case x_r must have the same sign as x_m . And so again, we can conclude that the sign of x_r has not changed. And so even in this case, we can conclude that the invariant is going to hold. So basically what we have is a simple method by which we can shrink this interval.

(Refer Slide Time: 9:37)



So, let us now see the code file. So, we are going to write this code for the problem that we mentioned earlier, which is finding the square root. And we will make the problem

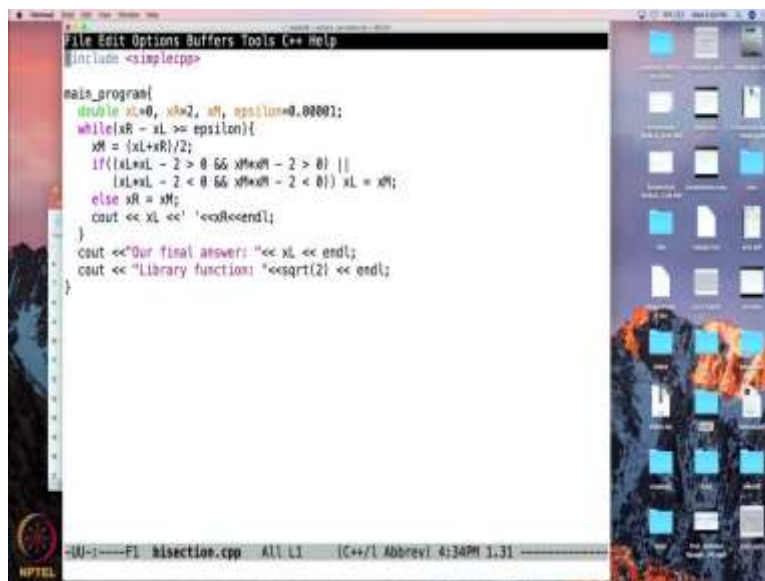
very concrete by saying that we want to find the square root of 2. So in that case, as we discussed earlier, we should be finding the root of $F(x)$ equal to x^2-2 . Because if x is root of this, then x^2-2 must be equal to 0 or x^2 must be equal to 2 or x must be square root of 2. So the root of this equation will indeed give us this square root of 2.

So first we need to supply x_l and x_r . So I claim that x_l equal to 0 will work. So let us see what that means. So x_l equal to 0 means, $F(x)$ equal to 0 minus 2. So $F(x_l)$ is negative. So, so long as we can supply an x_r , such that $F(x_r)$ is positive, we will have satisfied of requirement. So we are going to set x_r equal to 2. So $F(x_r)$ is 2 square or 4 minus 2 and therefore, it is positive. So our x_l and x_r are really satisfying, what we set out that their signs, the signs of $F(x_l)$ and $F(x_r)$ are indeed different. So then we need to have a variable to store this x_m value. And we need an epsilon as well. And let us say, just for, just as an example that we pick epsilon to be 10 to the power minus 5, 10 to the power minus 4. I am sorry, 10 to the power minus 5, okay. Anyway it does not matter. Now, we are going to just have to write our basic iteration. And, what is our basic iteration? Well, we check whether the interval x_r to x_l or the length of this interval which is x_r-x_l is bigger than epsilon. If, it is bigger than epsilon, then that means that our interval is still larger than what we want. So in that case, we should execute our basic step. So, what is our basic step? So, we are going to find x_m , which is the midpoint of x_l and x_r . And, so it is $x_l+x_r/2$. Then, we are going to check whether x_m and x_l have the same signs, okay.

So, what is this checking? This is checking whether $F(x_l)$ is bigger than 0. So, does x_l have the positive sign? And, this checks whether x_l has, x_m , $F(x_m)$ also have the positive sign. So, this is checking one side of the condition that we want, that $F(x_l)$ and $F(x_m)$, both have positive side, positive signs. So, either they both; so if they both have positive signs or if they both have negative signs, which is the second condition over here, then, that means they have the same signs. x_l and $F(x_l)$ and $F(x_m)$ have the same sign. In which case, we are going to set x_l equal to x_m . Otherwise, we are going to set, x_r equal to x_m . This is exactly what we said on the previous slide. That is it. So we are going to repeat this. We are going to repeat while this condition holds. We are going to stop as soon as this condition stops holding or in other words, when x_r-x_l becomes smaller than epsilon. Or in other words, our interval have become smaller than epsilon. So at that point

we are going to print out x_l as the root, alright. So, this is, this is what we said on the last slide and this is what we have here. So, this should work, and we will yeah. So, before moving on, I just want to ask you a simple question. So suppose instead of finding the square of 2, square root of 2, I want the square root of some other number. So, say square root of 3. How, could you choose x_l and x_r ? So, we can follow; we can see what is happening over here and maybe we will try to do something that is over here. So here we have 2. So let us try if we put, what happens if we put 3 over here. So indeed we will see that it works. Because, again this will be negative and this will be $9-3$, so this will be positive, okay. So, you can use the same idea to solve this for any number larger than 1, for any number yeah. And, we are, we are, yeah. So, let me leave it at that.

(Refer Slide Time: 14:49)

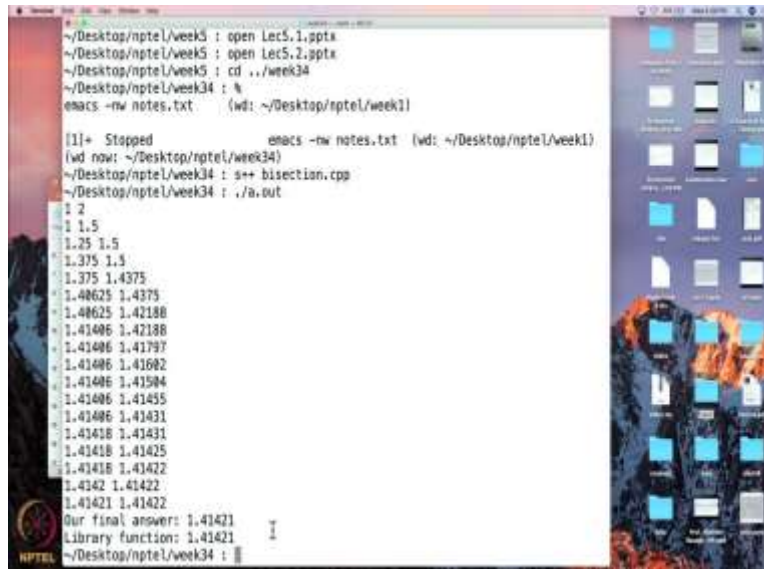
A screenshot of a C++ IDE window titled "bisection.cpp". The code implements a bisection method to find the square root of 2. It defines variables x_l , x_r , and x_m , and a constant ϵ (epsilon) set to 0.00001. A while loop continues until $x_r - x_l > \epsilon$. Inside the loop, x_m is calculated as $(x_l + x_r) / 2$. A conditional statement checks if $(x_l * x_l - 2) > 0$ and $(x_m * x_m - 2) > 0$. If true, x_l is updated to x_m . Otherwise, x_r is updated to x_m . A cout statement prints x_l and x_r at the end of each iteration. After the loop, a final cout statement prints "Our final answer: " followed by x_l . Another cout statement prints "Library function: " followed by $\sqrt{2}$ (using the sqrt function). The status bar at the bottom shows "F1 bisection.cpp All L1 [C++/I Abbrev] 4:34PM 1.31".

```
File Edit Options Buffers Tools C++ Help
#include <simplecpp>

main_program{
    double x_l=0, x_r=2, x_m, epsilon=0.00001;
    while(x_r - x_l > epsilon){
        x_m = (x_l+x_r)/2;
        if((x_l*x_l - 2 > 0 && x_m*x_m - 2 > 0) ||
           (x_l*x_l - 2 < 0 && x_m*x_m - 2 < 0)) x_l = x_m;
        else x_r = x_m;
        cout << x_l << ' ' << x_r << endl;
    }
    cout << "Our final answer: " << x_l << endl;
    cout << "Library function: " << sqrt(2) << endl;
}
```

So, let us do a quick demo of this. So, this is the program that we had on the slides. There is a slight difference. So we are printing our final answer that we calculated over here. But, in addition, in each iteration, we are also printing out the current interval. So the current values of x_l and x_r . And we are doing this at the end of iteration. So at the end of the first iteration, we will calculate what x_l and x_r are and so on for each iteration. So we are calculating the intervals. We are calculating the final answer. And then we know that there already exists a library function $\sqrt{}$ which gives us the square root. So we will print that also. So that will tell us, how good our answer is.

(Refer Slide Time: 15:40)

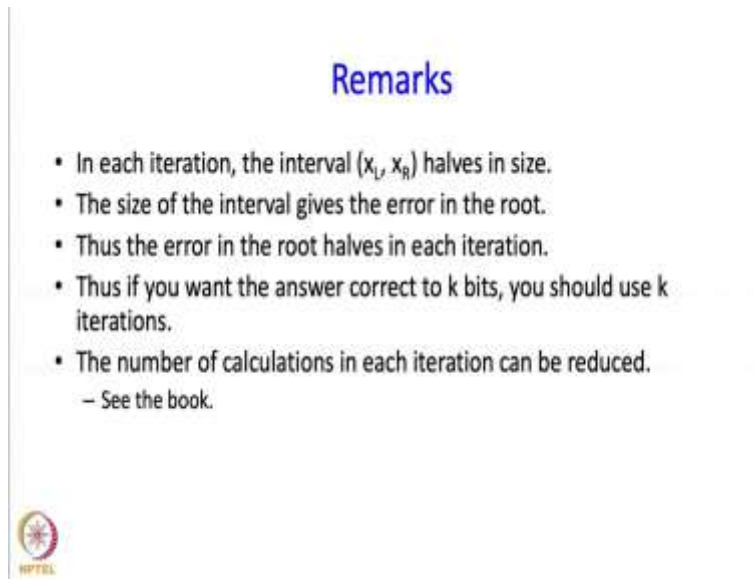


```
~/Desktop/npTEL/week5 : open Lec5.1.sptx
~/Desktop/npTEL/week5 : open Lec5.2.sptx
~/Desktop/npTEL/week5 : cd ../week34
~/Desktop/npTEL/week34 : %
emacs -nw notes.txt (wd: ~/Desktop/npTEL/week1)

[1]+  Stopped                  emacs -nw notes.txt (wd: ~/Desktop/npTEL/week1)
(wd now: ~/Desktop/npTEL/week34)
~/Desktop/npTEL/week34 : s++ bisection.cpp
~/Desktop/npTEL/week34 : ./a.out
1 2
1 1.5
1.25 1.5
1.375 1.5
1.375 1.4375
1.40625 1.4375
1.40625 1.42188
1.41406 1.42188
1.41406 1.41797
1.41406 1.41602
1.41406 1.41504
1.41406 1.41455
1.41406 1.41431
1.41418 1.41431
1.41418 1.41425
1.41418 1.41422
1.4142 1.41422
1.41421 1.41422
Our final answer: 1.41421
Library function: 1.41421
~/Desktop/npTEL/week34 : %
```


So, let me compile this. Okay, so let us see, what it has printed. So at the end of the first iteration the value, the interval was 1 to 2, okay. Does not makes sense? Well, we started off with 0 to 2, okay. x_l was 0 and x_r was 2. But then we moved x_l to the midpoint, so it was 1 to 2. And so the interval keeps on shrinking. And eventually it has got down to this. So it take some amount of work, number of iterations. But it eventually did get to this. And as you can see this is indeed identical to what the library function tells us. Of course, there may be additional digits after this, which are not being printed over here. And maybe the final answer is different. So, if you want to check that, you can print the additional digits also. But anyway to 5 digits, this is certainly the same.

(Refer Slide Time: 16:55)



Remarks

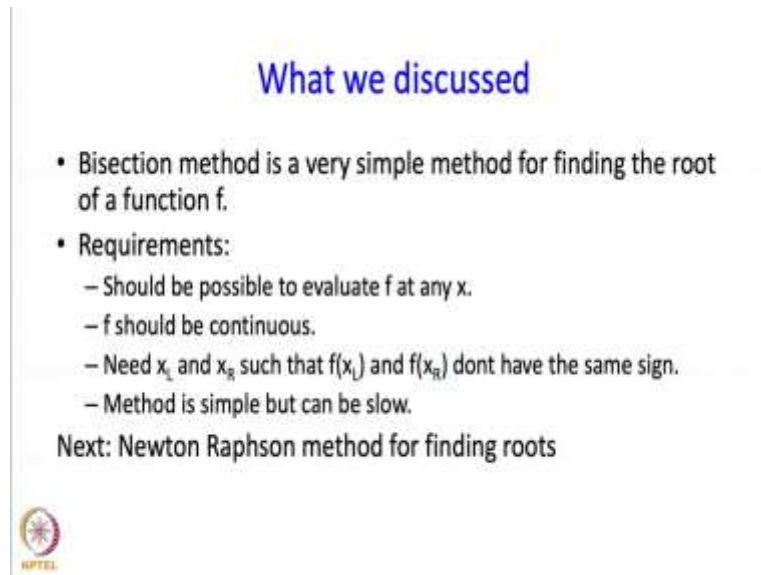
- In each iteration, the interval (x_l, x_r) halves in size.
- The size of the interval gives the error in the root.
- Thus the error in the root halves in each iteration.
- Thus if you want the answer correct to k bits, you should use k iterations.
- The number of calculations in each iteration can be reduced.
 - See the book.



Okay, so let us get back to the slides. So a few remarks. So because we are picking the midpoint in each iteration, this interval x_l to x_r is halving. So in each iterations the interval halves. So the uncertainty in our, in our estimate goes down by the factor of 2. Or in other words, we get one more bit of our final answer. And the size of interval gives the error in the root, we said. And so the error halves. And so if you want to answer correct to K bits you should use K iterations.

Now, the code that I have given over here is doing a few additional calculations over and above what is strictly needed. So I do not want to get into that detail right now in this lecture. But, you can take a look at the book and see how some of the work that you are doing could be reduced. And let me leave you with an exercise. I would like you to modify the program, so that it calculates the cube root of any number W . So the number W should be read in from the keyboard. And your program should correctly initialize x_l and x_r now. So think about how you would do, how you would have to do that. And then we should print out the cube root this time, not the square root.


(Refer Slide Time: 18:27)



What we discussed

- Bisection method is a very simple method for finding the root of a function f .
- Requirements:
 - Should be possible to evaluate f at any x .
 - f should be continuous.
 - Need x_l and x_r such that $f(x_l)$ and $f(x_r)$ don't have the same sign.
 - Method is simple but can be slow.

Next: Newton Raphson method for finding roots



Okay, so what did we discuss? So we discussed the bisection method, which is a very simple method for finding the roots of a function. So it requires that we should be able to evaluate F for any x and that F should be continuous. And we need to have x_l and x_r , such that $F(x_l)$ and $F(x_r)$ do not have the same sign. So if these three requirements are met, then we can use the bisection method. And I should point out that the method is simple but it is slow, in the sense that there are faster methods available. And next we will see the Newton-Raphson method, which is one such method. But before that we will take a break.