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Lecture –72 Monte Carlo Simulations (Part - 04)

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| Pseudo random number generators: Lehner algorithm The linear congruential generator: $\begin{aligned} \mathbf{x}_{j} = \alpha x_{j_{i_{1}}} & \text{mod} m \\ \mu_{j_{i_{1}}} x_{j_{m}}^{T} \\ \mathbf{x}_{j_{m}} = 2^{k_{j_{1}}} \\ \mathbf{x}_{m_{j_{m}}} = 2^{k_{j_{m}}} \\ \mathbf{x}_{m_{j_{m}}} = 2^{k_{j_{m}}}$ | tural Reliability Lecture 8 Monte Carlo simulations | lations | Carlo simu | Monte | |
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| The linear congruential generator: $\begin{aligned} \overline{x}_{i} = ax_{i-1} \mod m \\ \underline{u}_{i} = x_{i}/m \\ Generally, \\ m^{-2^{3}}, b = bits/word \\ a - \sqrt{m} \\ \underline{m}_{p} - 2^{i-2} \end{aligned} $ $\begin{aligned} m & \text{is prime} \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is a primitive root of } m \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible by } m. \\ & a^{-1} & \text{is exactly divisible b } m. \\ & a^{-1} & \text{is exactly divisible b } m. \\ & a^{-1} & \text{is exactly divisible b } m. \\ & a^{-1} & \text{is exactly divisible b } m. \\ & a^{-1} & \text{is exactly divisible } m \\ & a^{-1} & is exactly d$ | Simulations | rithm | erators: Lehmer alg | dom number gen | Pseudo rando |
| $ \begin{array}{c} x_i = ax_{i-1} \mod m \\ u_i = x_i/m \\ \hline \\ generally, \\ m - 2^k, b = bits word \\ a - \sqrt{m} \\ m_i - 2^{b-2} \end{array} $ a is a primitive root of <i>m</i> iff are multiplier $m_0 = \operatorname{priod} df$ generator, $a^i \operatorname{mod} m \neq 1$ for $n = 1, 2,, m - 2.$ That is, no integer from the sequence $a - \sqrt{m} \\ m_i - 2^{b-2} \\ \hline \\ m_i - 2^{b-2} \\ \hline \\ \hline \\ \hline \\ m_i - 2^{b-2} \\ \hline \\ \hline \\ \hline \\ \hline \\ m_i - 2^{b-2} \\ \hline \\ $ | | tor is a full period generator if and only if: | This genera | al generator: | The linear congruential |
| This value for a way first arranged by Lamis Condman or A Clu-in | | rimitive root of m a is a primitive root of m iff $a^{n} \mod m \neq 1$ for $n = 1, 2,, m - 2$. That is, no integer from the sequence $a - 1, a^{2} - 1,, a^{m^{n}2} - 1$ is exactly diviable by m . Fermat's ling theorem, $a^{n} = a \pmod{m} \Leftrightarrow a^{n-1} = 1 \pmod{m}$ | • a is a prator, $x_{i_1a_2} = x_i$ *Best* RN Generator: $m = 2^{11} - 1$ $a = 7^5 = 16807$ | a = multiplier $m_0 =$ period of gene if for smallest m_0 , | $u_i = x_i / m$ Generally, $m \sim 2^{b}$, $b = \text{bits/word}$ $a \sim \sqrt{m}$ |
| Derick Henry Lehner (1905 – 1991) developed the linear congruential generatorian in 45% wills extense threadings with a charter threadings. Random Number Generators: good ones are hard to find, by Stephen K Park and Kenner Computed to Mile, in Computing Practices, Communications of the ACM, vo. 31, October. 1988. | 6 | iginal work. | 1969 while Lehmer himsel for the modulus <i>m</i> in his or Further reading: Random Number Generators: g Keith W Miller, in Computing Pr | congruential generator | developed the linear con |

We conclude the topic of random number generation with a look at the Lemur algorithm the linear congruential generator it is what you see on the screen it starts with a seed and then recursively generates the next number in the sequence with modular arithmetic and being the modulus a being the multiplier and the number xi is normalized by m to obtain the uniform deviate, m0 is the period of generator and the numbers m and a are typically chosen depending on the word size of the machine.

And it is it is found that this generator is a full period generator only if the modulus is a prime number and a is a primitive root of m. Now this generator has been studied extensively and the best set of parameters MNA the best RNG it is found to be when m is the Mercen prime to the power 31 - 1 and a is 7 to the power of 5. In this case this generator is a full period generator and m 0 is m-1 and these values were first proposed by for aid was by Lewis et al in 1969.

And the value of m was by Lemur himself in his original work in 1948. I have listed this excellent paper by Park and Miller which discusses this algorithm and others. So, now let us just end this discussion with one or two small examples of the Lemur algorithm.

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| $x_i = \alpha x_{i-1} \mod m$ $u_i = x_i / m$ Given a = 2, m = 2 ³ - 1 What is the period of this generator? 1 3 2 6 4 5 9 Period=3 Period=3 Given a = 3, m = 2 ³ - 1 | | Given $a = 3$, $m = 2^5 - 1$ What is the period of this generator? | |
| What is the period of this generator? | ag ac inf-bailanya/ | | 220 |

Let us say that let us have the modulus as 2 to the power 3 - 1 and the multiplier a as 2. So, what would be the period of this generator. If you want to work this through then please pause the video otherwise let me present the solution let us let us choose the seed as 1 and then we find that the period is 1 well the period is 3 that the sequence is 1, 2, 4 and then it starts repeating itself. If you choose a different seed then you hit another sequence and obviously there is no intersection between the two sequences.

If on the other hand we choose the multiplier as 3 then you will find that this becomes indeed becomes a full period generator. So, the choice of the multiplier is actually very important. Let us look at the next Marseille prime to the power 5 - 1 as the modulus and if a is 2 then what would be the period of this generator if you again want to work this out then please pause the video otherwise it turns out that depending on the seed you get different sequences and none of them is a full period generator the period is 5 and there are 6 such sequences.

On the other hand if a is chosen to be 3 then this generated indeed becomes a full period

generator.