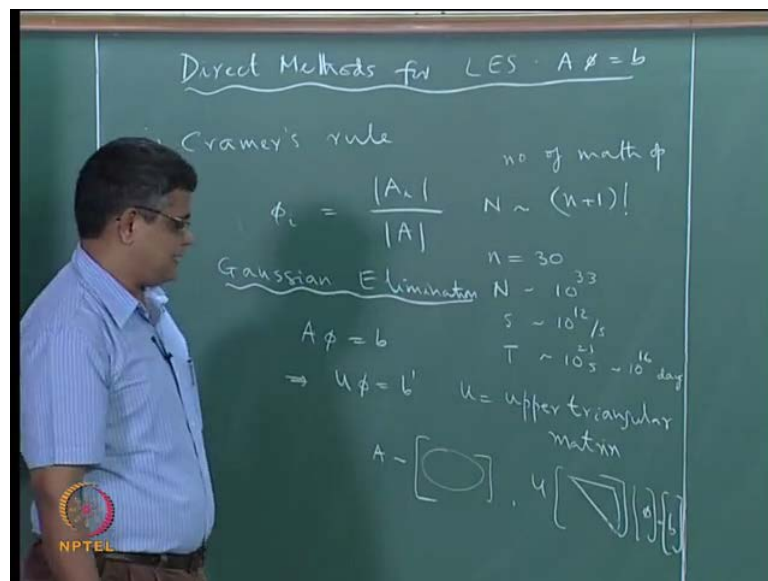


Computational Fluid Dynamics
Prof. SreenivasJayanti
Department of Chemical Engineering
Indian Institute of Technology, Madras

Module No.# 05
Solution of Linear Algebraic Equations
Lecture No. #23
Direct Methods for Linear Algebraic Equations
Gaussian Elimination Method

Let's look at some direct methods for the solution of $A\phi = b$.

(Refer Slide Time: 00:17)



So, we will call this as direct methods for linear equations. So, these are linear equation solvers LES, $A\phi = b$ and as we have said earlier, all the coefficients A are known and all the coefficients b are known therefore, this constitutes a linear equation a set of linear equations. And we also mentioned that a direct method is a method in which after a finite number of steps, you get the overall solution.

We mentioned several methods, the first method that we need to look at, we need to mention for the sake of completion is the Cramer's rule in which ϕ_i is given by $\phi_i = \frac{|A_i|}{|A|}$

determinant of A_i by A , where A_i is the coefficient matrix A here in which the i th column, where i represents the variable that we are looking for, the i th column is replaced by the column vector b . This way you have to replace the i th column by column vector b and find the determinant of the resultant matrix and divide that by the determinant of the coefficient matrix. So, this method is taught in all high school syllabus courses and the advantage of this method, that this will work for any nonsingular matrix A is outweighed by the difficulty of this for large matrices.

The number of mathematical manipulations mathematical operations for this method is $N!$ plus 1 factorial, where N is the number of variables that we have. So, if N is 30, the number of mathematical operations $N!$ is of the order of 10 to the power of 33 . So, if you have a speed of mathematical operations of 10 to the power of 12 per second, this would take the time taken for the solution of 30 equations using this method would be 10 to the power 21 seconds.

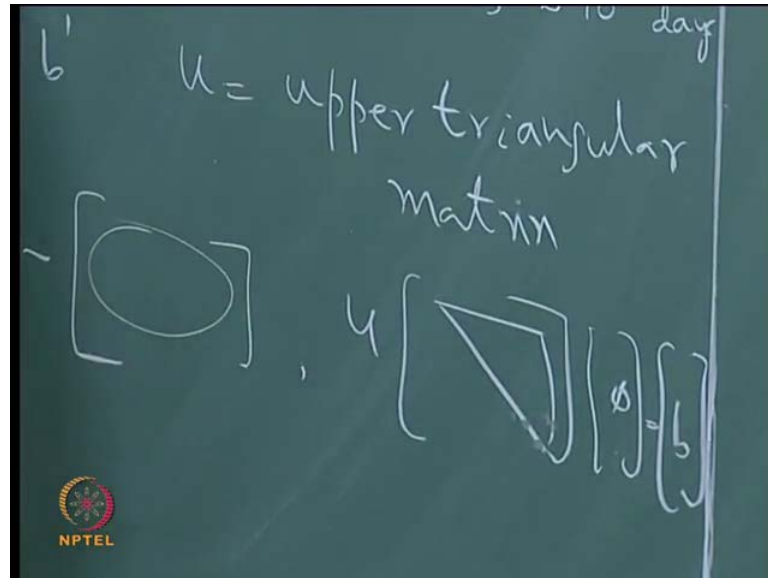
So, that is a very huge number that is 10 to the power, let us say 15,16 days. So, this is why this method we mention it for the sake of completion and especially in cases, where n is very small, may be this method would can still be used; if you have n equal to 4, may be this method is much easier and n equal to 3, we can use these kind of methods.

The most general method that one can recommend is a Gaussian elimination method; this is the probably the most useful and efficient methods method for the solution of $A\phi = b$, in case where A does not have any special features like bandedness or symmetry or sparseness. So, in such cases when A is almost fully populated without any specific structure, diagonal structure or other structure symmetry in such a case this can be taken to be the most efficient method and how do we do this? So, if you have $A\phi = b$, the objective is to convert this into a matrix like $u\phi = b'$, where u is an upper triangular matrix.

And upper triangular matrix is of, if A is like this, then with all these coefficients here, u will have coefficients only in this. So, if you now put this $u\phi = b'$, you will have n equations here and as you go down in this way, you will have fewer and fewer variables right at the bottom, you have a single variable here. So, that equation can be solved directly to give you ϕ_n . Now, you come to the upper level, there you have $n-1$ and n as the two variables and n is already known. So, you can solve this

directly for ϕ_{n-1} and then, if you come to the next row here, next higher row you have which will be $n-2$, $n-1$, n and by now you have computed $n-1$, you can directly get $n-2$.

(Refer Slide Time: 06:30)



So, the solution of $u \phi$ equal to b , where u is an upper triangular matrix is something that can be done easily by back substitution. So, the idea of the Gaussian elimination is to successively eliminate these variables; you have ϕ_1 , ϕ_2 , ϕ_3 , ϕ_4 like that. So, you eliminate ϕ_1 from all these equations, ϕ_2 from all the lower equations, ϕ_3 from all the lower equations and soon, until you get a single variable left in the last equation and thereby **in an** in the first stage of elimination, you convert A from a matrix which has coefficients all over the place into one which has **coefficients** only nonzero coefficients only in the upper triangular part of the original matrix A and the solution of this can be done efficiently.

(Refer Slide Time: 07:42)

$$\begin{aligned} a_{11} \phi_1 + a_{12} \phi_2 + a_{13} \phi_3 + a_{14} \phi_4 + \dots + a_{1n} \phi_n &= b_1 \\ a_{21} \phi_1 + a_{22} \phi_2 + a_{23} \phi_3 + a_{24} \phi_4 + \dots + a_{2n} \phi_n &= b_2 \\ a_{31} \phi_1 + a_{32} \phi_2 + a_{33} \phi_3 + a_{34} \phi_4 + \dots + a_{3n} \phi_n &= b_3 \\ &\vdots \\ a_{n1} \phi_1 + a_{n2} \phi_2 + a_{n3} \phi_3 + a_{n4} \phi_4 + \dots + a_{nn} \phi_n &= b_n \end{aligned}$$

Now, how do we do this elimination? So, let us take a specific case $a_{11} \phi_1 + a_{12} \phi_2 + a_{13} \phi_3 + a_{14} \phi_4 + \dots + a_{1n} \phi_n = b_1$. The next equation will be $a_{21} \phi_1 + a_{22} \phi_2 + a_{23} \phi_3 + a_{24} \phi_4 + \dots + a_{2n} \phi_n = b_2$ and we can write one more, $a_{31} \phi_1 + a_{32} \phi_2 + a_{33} \phi_3 + a_{34} \phi_4 + \dots + a_{3n} \phi_n = b_3$ and we have n such equations and the last one will be $a_{n1} \phi_1 + a_{n2} \phi_2 + a_{n3} \phi_3 + a_{n4} \phi_4 + \dots + a_{nn} \phi_n = b_n$.

All these coefficients b are in this side and all the coefficients a_{11} and a_{12} and all these things, they form the matrix the coefficient matrix a and ϕ that is $\phi_1 \phi_2 \phi_3 \phi_4 \phi_5$, all those things all the way up to ϕ_n are the variables.

So, in this, the idea of the elimination process is to make use of this equation which is known as the pivot equation and eliminate ϕ_1 from all these equations and how do we do that? If you say that, if you multiply this equation by a_{21}/a_{11} and subtract from this equation, then you get this term will be $a_{21} - (a_{21}/a_{11}) \times a_{11}$. So, this a_{11} will cancel out and what you have is $a_{21} - (a_{21}/a_{11}) \times a_{11}$, when you subtract from this, ϕ_1 gets cancelled out and this term here will become...

(Refer Slide Time: 10:19)

$$\begin{aligned}
 & a_{11} \phi_1 + a_{12} \phi_2 + a_{13} \phi_3 + a_{14} \phi_4 + \dots + a_{1n} \phi_n = b_1 \\
 & a_{21} \phi_1 + a_{22} \phi_2 + a_{23} \phi_3 + a_{24} \phi_4 + \dots + a_{2n} \phi_n = b_2 \\
 & a_{31} \phi_1 + a_{32} \phi_2 + a_{33} \phi_3 + a_{34} \phi_4 + \dots + a_{3n} \phi_n = b_3 \\
 & \vdots \\
 & a_{m1} \phi_1 + a_{m2} \phi_2 + a_{m3} \phi_3 + a_{m4} \phi_4 + \dots + a_{mn} \phi_n = b_m
 \end{aligned}$$

$$\begin{aligned}
 & a_{11} \phi_1 + a_{12} \phi_2 + a_{13} \phi_3 + \dots + a_{1n} \phi_n = b_1 \\
 & \left(a_{22} - \frac{a_{21} a_{12}}{a_{11}} \right) \phi_2 + \left(a_{23} - \frac{a_{21} a_{13}}{a_{11}} \right) \phi_3 + \dots + \left(a_{2n} - \frac{a_{21} a_{1n}}{a_{11}} \right) \phi_n = \left(b_2 - b_1 \frac{a_{21}}{a_{11}} \right)
 \end{aligned}$$

So, you have the first equation as it is, $a_{11} \phi_1 + a_{12} \phi_2 + a_{1n} \phi_n = b_1$ and here this becomes 0, because we have cancelled this out and this equation here will become... your subtracting this from this; this multiplied by a_{21} by a_{11} . So, this will now become $a_{22} - a_{21} \frac{a_{12}}{a_{11}}$ by a_{11} times $a_{12} \phi_2$ and similarly, $a_{23} - a_{21} \frac{a_{13}}{a_{11}}$ times $a_{13} \phi_3$ plus soon plus $a_{2n} - a_{21} \frac{a_{1n}}{a_{11}}$ by a_{11} times ϕ_n equal to $b_2 - b_1 \frac{a_{21}}{a_{11}}$.

So, we have modified the second equation by multiplying the first equation - the pivot equation - by this value $\frac{a_{21}}{a_{11}}$ and then, subtracting that resulting equation from this, so that the second equation gets modified like this and it does not have ϕ_1 .

(Refer Slide Time: 15:00)

$$\begin{aligned}
 & a_{21}\phi_1 + a_{22}\phi_2 + a_{23}\phi_3 + a_{24}\phi_4 + \dots + a_{2n}\phi_n = b_2 \\
 & a_{31}\phi_1 + a_{32}\phi_2 + a_{33}\phi_3 + a_{34}\phi_4 + \dots + a_{3n}\phi_n = b_3 \\
 & \dots \\
 & a_{n1}\phi_1 + a_{n2}\phi_2 + a_{n3}\phi_3 + a_{n4}\phi_4 + \dots + a_{nn}\phi_n = b_n
 \end{aligned}$$

$$\begin{aligned}
 & a_{11}\phi_1 + a_{12}\phi_2 + a_{13}\phi_3 + \dots + a_{1n}\phi_n = b_1 \\
 & \left(a_{32} - \frac{a_{31}a_{12}}{a_{11}} \right) \phi_2 + \left(a_{33} - \frac{a_{31}a_{13}}{a_{11}} \right) \phi_3 + \dots + \left(a_{3n} - \frac{a_{31}a_{1n}}{a_{11}} \right) \phi_n = \left(b_3 - b_1 \frac{a_{31}}{a_{11}} \right) \\
 & \left(a_{52} - \frac{a_{51}a_{12}}{a_{11}} \right) \phi_2 + \left(a_{53} - \frac{a_{51}a_{13}}{a_{11}} \right) \phi_3 + \dots + \left(a_{5n} - \frac{a_{51}a_{1n}}{a_{11}} \right) \phi_n = \left(b_5 - b_1 \frac{a_{51}}{a_{11}} \right) \\
 & \vdots \\
 & \left(a_{n2} - \frac{a_{n1}a_{12}}{a_{11}} \right) \phi_2 + \left(a_{n3} - \frac{a_{n1}a_{13}}{a_{11}} \right) \phi_3 + \dots + \left(a_{nn} - \frac{a_{n1}a_{1n}}{a_{11}} \right) \phi_n = \left(b_n - b_1 \frac{a_{n1}}{a_{11}} \right)
 \end{aligned}$$

NPTEL

Similarly, we take the third equation and then, we multiply this by a $\frac{a_{31}}{a_{11}}$. So, a $\frac{a_{31}}{a_{11}}$ times the pivot equation. So, this much is subtracted from this; so when we do that, a_{31} and this will cancel out; you get a $\frac{a_{32}}{a_{11}}\phi_1$, when you subtract from this, that cancels out and this term here becomes a $\frac{a_{32}}{a_{11}}\phi_2$ plus a $\frac{a_{33}}{a_{11}}\phi_3$ plus a $\frac{a_{3n}}{a_{11}}\phi_n$ equal to $b_3 - b_1 \frac{a_{31}}{a_{11}}$.

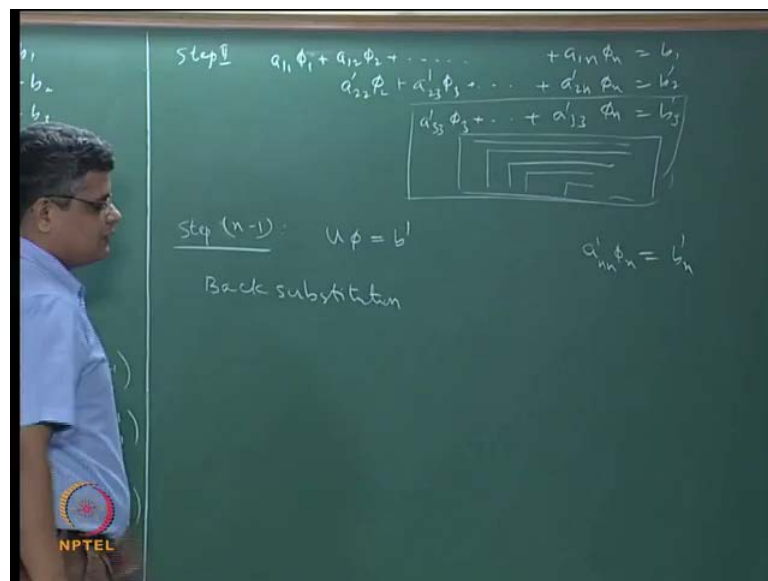
So, we have eliminated from the third equation also ϕ_1 by multiplying the pivot equation by this product, the product of the coefficient of **the first equation** the first term divided by the coefficient of the first term of the pivot equation. So, by doing that, we are able to eliminate ϕ_1 from the third equation also and then, we go on and then do it like this for all the $n-1$ equations and then finally, we have... So, this is zero plus zero plus zero plus you will have a $\frac{a_{n2}}{a_{11}}\phi_2$ plus a $\frac{a_{n3}}{a_{11}}\phi_3$ plus a $\frac{a_{nn}}{a_{11}}\phi_n$ equal to $b_n - b_1 \frac{a_{n1}}{a_{11}}$.

So, what we have now is **is** a set of equations in which a_{11} appears only in the first equation and in all the rest of the equations we do not have a_{11} and this becomes a smaller set with $n-1$ variables.

So, now we keep this like this; we take this **this** set of equations as something like a phi equal to b and then, **this is written as we eliminate there is something here yeah we eliminate**, we retain this and then, we eliminate phi 2 from this, phi 2 from the next equation, phi 2 from the next equation and soon and then, we eliminate phi 2 from **all the succeeding equations** all the n minus 2 succeeding equations; at the end of that, we will have **1 1 only in the first equation**, phi 1 only in the first equation, phi 2 only in the second equation and then, you will have a subset of n minus 3 equations in which phi 3, phi 4 and all those things appear and then, **you can eliminate** you can retain the third one and eliminate phi 3 from all the rest of the equations and then, we can go on doing this until we eliminate all, but phi n in the last equation. So, **this is if** this is original one.

(Refer Slide Time: 16:40) Step 1 will give us this.

(Refer Slide Time: 16:42)

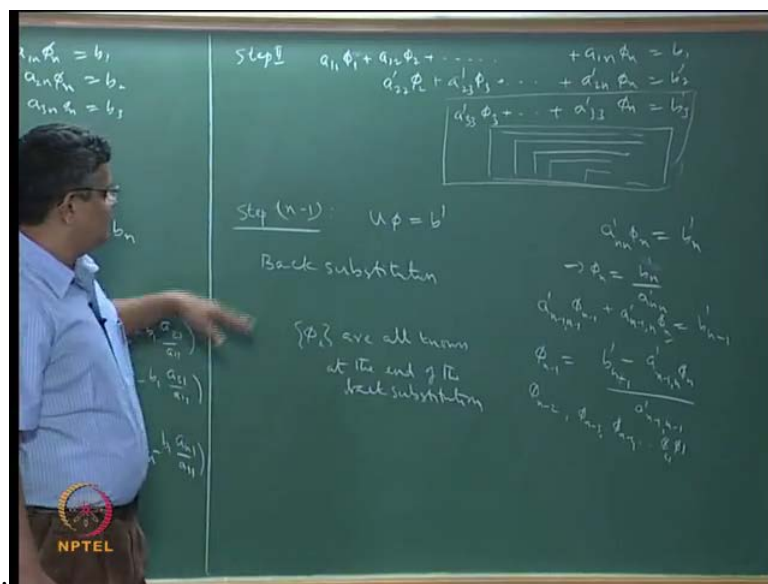


And step 2 will give us this, $a_{11}\phi_1 + a_{12}\phi_2 + \dots + a_{1n}\phi_n = b_n$ and then, $a'_{22}\phi_2 + \dots + a'_{2n}\phi_n = b'_2$; I am putting prime here, because that is modified; $a'_{23}\phi_3 + \dots + a'_{2n}\phi_n = b'_2$ plus $a'_{2n}\phi_n = b'_2$ and so on.

And. So, now, if we eliminate phi 2 from all these things, we will have $a'_{33}\phi_3 + \dots + a'_{3n}\phi_n = b'_3$ plus $a'_{3n}\phi_n = b'_3$. So, we will have a subset of these equations n minus 2 equations in which you have this. So, we **we** keep this and then from this subset, we eliminate phi 3. So, that we have this like this and after this, we retain the top one and from this subset we

eliminate ϕ_4 and again we eliminate ϕ_5 like this and then, we go on like this until we are left with an upper triangular matrix which will have finally, like this and then at the end of step $n-1$, we will have all the first equations and then we will have a $n \times n$ prime ϕ_n equal to b_n prime, where this prime of course has been multiplied several times. So, this is the final form and we have an upper triangular matrix and this is the form we have as **u prime equal to** ϕ equal to b prime. So, **and** this we solve by back substitution.

(Refer Slide Time: 19:16)



We will solve this for ϕ_n equal to b_n prime by a_{nn} prime and then above this we have a $n-1$ prime $n-1$ ϕ_{n-1} plus $a_{n-1,n}$ prime ϕ_n equal to b_{n-1} prime. So, this is the resulting equation here and in this we already know this value from here and then we can evaluate this. So, this is ϕ_{n-1} will be equal to b_{n-1} prime minus $a_{n-1,n}$ prime ϕ_n divided by $a_{n-1,n-1}$ prime. So, like this then we go to ϕ_{n-2} and then ϕ_{n-3} ϕ_{n-4} up to 1 up to ϕ_2 and then ϕ_1 .

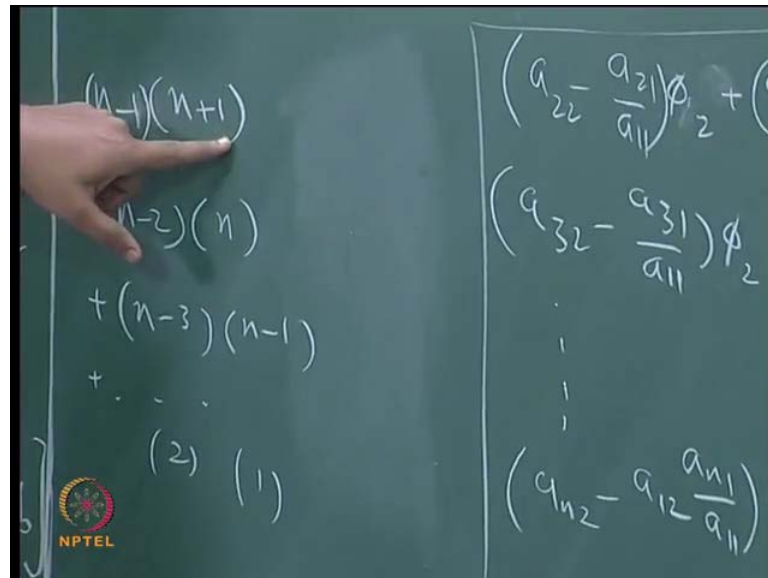
So, **we have** in the Gaussian elimination, we have a forward elimination stage in which we operate on $n-1$ equations to eliminate ϕ_1 and then, we operate on $n-2$ equations to eliminate ϕ_2 and then $n-3$ equations to eliminate ϕ_3 like that and then we ultimately end up with an upper triangular matrix and that triangular upper triangular matrix is solved by back substitution from the bottom row up to the first row

in a series of substitutions of values or variables which are already determined. So, at the end of the back substitution, we get all the ϕ_i values are all known at the end of the back substitution.

So, the Gaussian elimination is a good illustration of a direct method. You have a series of orchestrated steps by which you successively eliminate $\phi_1, \phi_2, \phi_3, \phi_4$ in a time consuming tedious calculation procedure to convert a into an upper triangular a ϕ equal to b into $u \phi$ equal to b' and this is then solved efficiently by a back substitution process and what is the number of mathematical operations? We have said that for the Cramer's rule it is $n!$, this looks like a very tedious thing you have to do so many multiplications and soon. We can make an estimate of $n!$, for example, to do the first step we have this; we have to multiply we have to determine this coefficient here with which we are multiplying each term here and then subtracting it. So, we need one multiplication here and then, we want to know that this is going to be 0, because we have done this. So, from all these $n-1$ terms that are here plus the right hand side term, we have to multiply by this. So, that is $n-1$ multiplications to do this multiplication here up to this and then n th multiplication. So, that gives you a total of $n-1$ of this plus 1 plus 1. So, that is n plus 1 multiplications are required to eliminate ϕ_1 from 1 equation and we have to do there are $n-1$ equations from that are from which we have to eliminate ϕ_{n-1} .

(Refer Slide Time: 24:23) So, the first step here takes n by n times n plus 1 number of operations sorry $n-1$ number of equations are there in which to eliminate ϕ_1 and for each elimination, we require n plus 1 terms. So, this is for the first step and for the second step of course, we have a smaller set. So, we have less number of equations to deal with. So, the second step will require $n-2$ elimination row elimination each requiring n multiplications and then the third one will require $n-3$ times $n-1$ plus soon all the way up to when we have only one thing here and then we have two multiplications for the last one.

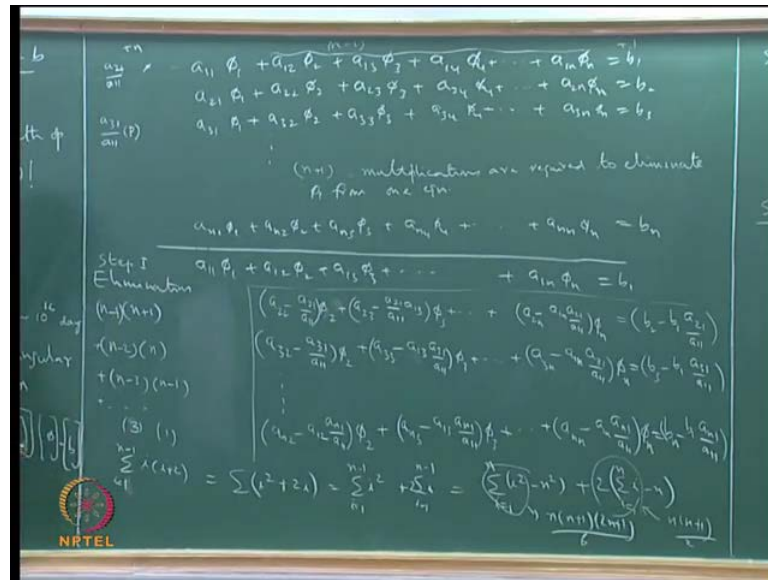
(Refer Slide Time: 25:33)



So, **the total number of** the total number of multiplications and of course, divisions that are required to do; step 1 elimination is n minus 1 times n plus 1; step 2 is n minus 2 times n ; step 3 is n minus 3 times n minus 1 and soon like this.

So, the total number required is the sum of all these things. So, we can write this as sum of i times i plus 2 this where i goes from 1 to n minus 1. If you want to get exactly this, then this **this** thing here is probably wrong, because this is 1 and 3. So, this must be 3 and 1 in which case this will be correct, because the difference between these two is 2. So, this will be 1 and 3 obviously, because in the last equation, we have this term and then the ϕ_1 coefficient to be determined and the b coefficient to be determined. So, it is 3 times one here. So, this formula works. So, the total number of mathematical multiplications and divisions required for the elimination is this.

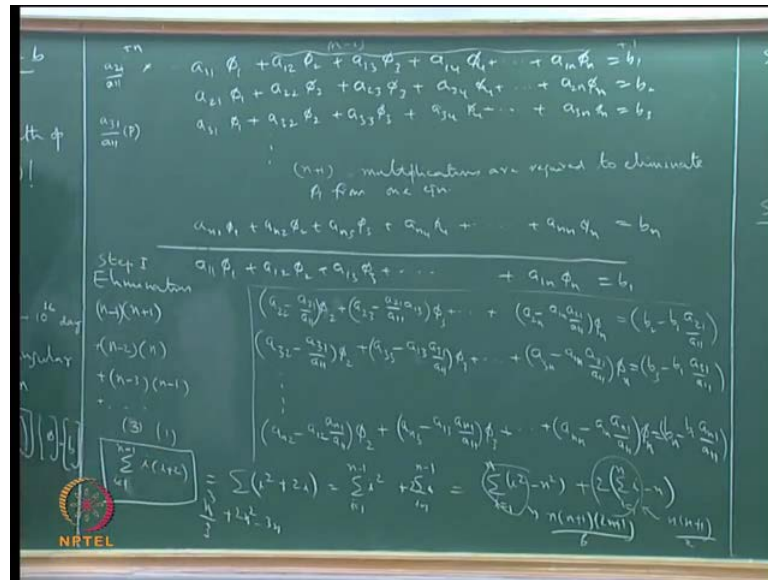
(Refer Slide Time: 26:45)



So, this is $i^2 + 2i$. So, that is sum of i^2 plus sum of $2i$, where i is from 1 to $n-1$ and this also i equal to $n-1$ and we have a formula for this. So, we can write this as $i^2 - n^2$; now i going from 1 to n .

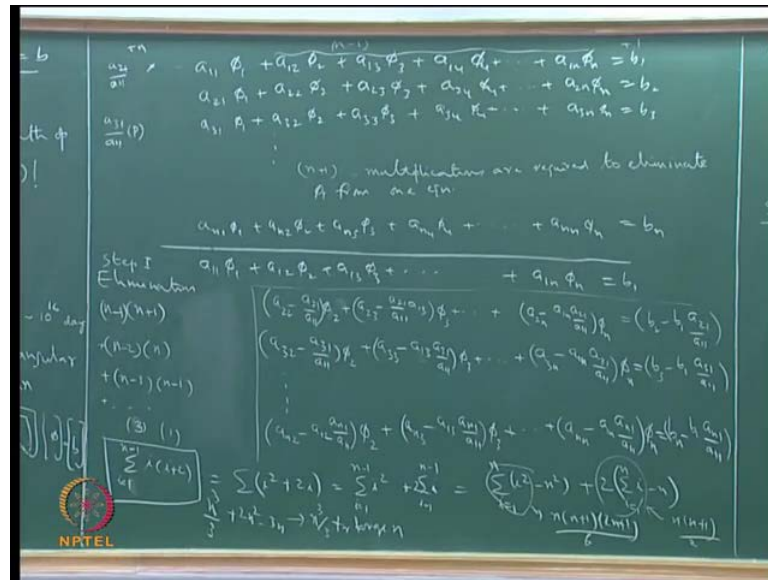
So, the last term will be $n^2 - n^2$, so we are cancelling out plus 2 times i equal to 1 to $n-1$. So, we have readymade formulas for this; when i goes from 1 to n , then **this is** this particular thing is n times $n+1$ times 2 plus 1 times 6 and this particular thing is n times $n+1$ by two. So, we can substitute these things and then finally, show that the number of operations required in the elimination process is something like $\frac{1}{3}n^3 + 2n^2$ or $\frac{1}{3}n^3$ **some** something like this, some $3n$ or something like that.

(Refer Slide Time: 28:40)



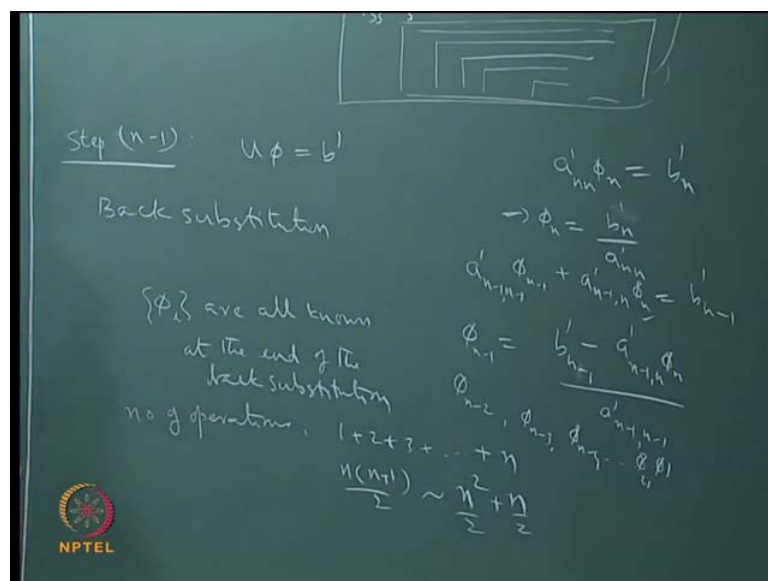
What is important is that, when n tends to large numbers like 1000, then this will be 10 to the power 9 and this will be 10 square. So, **that is only this is** this is a billion and this will be a million. So, billion plus 2 million or billion plus 1 million is nothing much. So, this tends to n cube by 3 for large n .

(Refer Slide Time: 29:14)



So, the idea is that in the Gaussian elimination process method, the elimination stage requires about $n^3/3$ number of steps and whatever back substitution. So, back substitution one can similarly make an operation count, for example, **for the first** for the bottom most row, we have taken one calculation, for the second bottom most row we require one calculation here and one more division here. So, that is 2 and for third one we need 3.

(Refer Slide Time: 30:00)



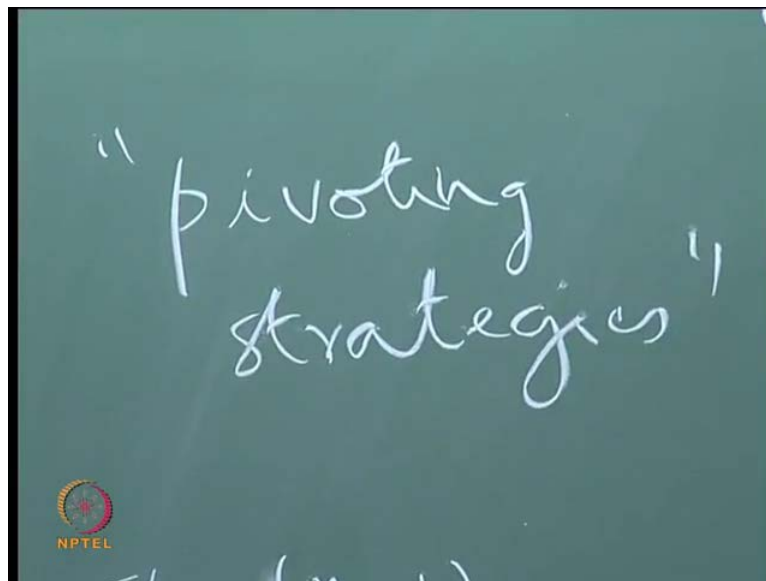
So, it becomes the number of operations is $1 + 2 + 3 + \dots + n$. So, this is $\frac{n(n+1)}{2}$. So, this is $\frac{n^2}{2} + \frac{n}{2}$. So, for large value of n , the back substitution takes only n^2 and the forward substitution takes n^3 . So, the point that we are trying to make is that for the entire solution for the Gaussian elimination solution, the number of multiplications or divisions required is of the order of n^3 and what is also important to note in this is, that the solution ϕ_i is obtained only after we do the elimination process and then the back substitution. And the back substitution constitutes only a small fraction of the overall computation required, because the elimination process requires n^3 number of operations and back substitution requires only n^2 number of operations. So, a small fraction if n is 1000, then only 0.1 percent of the time is required for back substitution and 99.9 percent effort is required for the forward elimination process and when n is even larger, that fraction correspondingly decreases.

So, what this means is that which is characteristic of a direct method and especially when the number of operations goes as n^3 or n^2 and all that is, that the solution from a direct method is known only after we do a large number of computation without even getting a hint of what the solution is. Only after we do the 99.9 percent of the overall solution effort do we get any idea of what the solution is? There is no intermediate guess/estimate of what the solution is and that is the disadvantage of a direct method when we want to incorporate this in an overall iterative procedure.

Typically, we have seen that in the method that we had been using for the simultaneous solution of the coupled equations of the Navier-Stokes equation, we make ϕ equal to b by making some estimates, for example, v and w and p when we want to solve the u momentum equation. So, that means, that this solution ϕ equal to b has coefficients which are determined, but which are not fully correct. So, even though we know that these coefficients are not fully correct, these are approximate we cannot get an approximate answer from the direct method, until we do 99.9 percent of the overall computation effort then it is too late. So, that is the disadvantage that a direct method typically has. Even though that is a case the direct method has the assurance that no matter what the structure of A is as long as there is a unique solution that is as long as the determinant of A is not zero, we can be confident that this method will give us a solution. So, we have the assurance of the widest possible applicability of the Gaussian elimination for the solution of $A\phi = b$; you do not have to worry whether

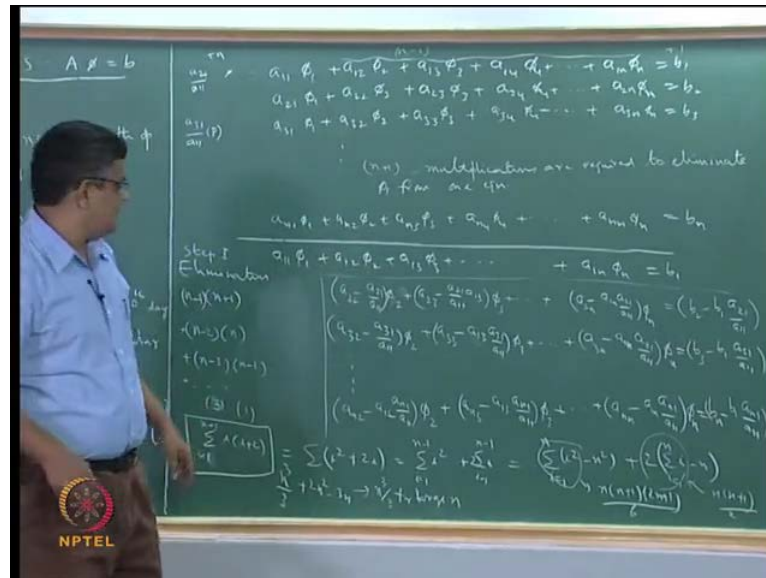
diagonal dominance is there or if it is not there. So, that kind of condition is not... Another feature of the Gaussian elimination method is that, we can see that in step 1 when you want to eliminate ϕ_1 , all the coefficients of the lower things are getting multiplied by a_{11} and then divided by a_{11} like this and then you are doing subtractions and then once you come to the third next step, even this a_{33} , which is originally modified to eliminate ϕ_1 is now again modified by division and multiplication by some other numbers to eliminate ϕ_2 and when you go to the next step, the ϕ_4 coefficient here which is already modified for ϕ_2 and ϕ_1 is again modified by further multiplications. So, as you go further and further down, each coefficient is modified several times by multiplication, division like that. So, unless you are careful, you can the round of error in representing these numbers will become intolerable when you have finite precision of arithmetic and when you have large number of equations.

(Refer Slide Time: 35:20)



So, that is why in order to reduce this error, strategies like what is known as the pivoting various pivoting strategies are incorporated in trying to find out the best pivot equation in order to do the elimination process.

(Refer Slide Time: 34:23)



So, at each step, for example, you take that coefficient which has the highest value, because **then if** since you are dividing by this a_{11} , all these things if a_{11} happens to be very small, then this coefficient here gets unnecessarily magnified and then, when you are trying to subtract from this, then because this is magnified here, **the** this subtraction may lose some significant digits.

So, you would like to **you'd like to minimize the** those kind of errors by choosing that coefficient among these things which has the highest value as the pivot. So, you have a **row wise exchange of** exchange of rows or exchange of columns is usually done to identify the right kind of pivot to do the elimination process.

So, usually when you are looking at a $Ax = b$ type of matrix solution, where A is full that is more or less with many coefficients which are nonzero, in such a case you would have to have some sort of pivoting strategy **that has to be...** and you have to do more much more than doing writing a program here, you have to do more logic type of programming steps to introduce a pivoting strategy. So, that successive elimination and modification of the coefficient does not lead to run off type of situations.

So, including all those pivoting strategies, Gaussian elimination is considered as the most efficient method for the solution of $Ax = b$, when A is full and does not have any structure. So, from that point of view this is a very good method.

When we come to c f d type of problems, where a is mostly sparse, we know that even if we have thousand equations, that is thousand variables the value of one variable is influenced only by only the 4 neighboring points or six neighboring points or 2 neighboring points. So, that means, that each row will have only a few nonzero coefficients. So, in such a case the buildup of arithmetic error because of successive multiplications and subtractions and all that is going to be less when you form a matrix obtained from a c f d solution.

Obtained for a c f d solution, because the discretization will make sure that only a few of the entire row have nonzero coefficient. So, in such a case the round of error and pivoting strategy is not. So, important for a c f d type of a ϕ equal to b, but for a general case it is very much true.

So, the point that we want to take from this is that, this is a general method which can be used for any nonsingular $A \phi$ equal to b type of equation and it requires n^3 number of mathematical operations, which is a great improvement on n plus 1 factorial, for example, when you say n equal to 30, then this only 30 cubed. So, that is about thirty thousand operations, where as this is giving to 10 to the power of 3 number of operations. So, this is very good for typically and less than thousand, we can safely use this kind of thing and there are certain applications in which they would prefer to use these type of elimination type of methods, because for a very coarse grid, may be you can you can use this type of situation.